Date

02.06.12

Cosmos

① Cormen
② Google

Algorithms

Syllabus

① Analysis
② Divide & Conquer
③ Greedy Technique
④ Dynamic Programming
⑤ Graph, Tree
⑥ Hashing
⑦ P, NP, NPC, NPH

Algorithm

Definition: It is a combination of sequence of finite steps to solve a particular problem.

Properties of Algorithm:-

① It should take finite time to produce o/p.
② It should produce the correct o/p.
③ Every step in the algo should be unambiguous (or) deterministic.
④ Every step in the algorithm should perform some task
⑤ Every algo should produce atlest 1 output.
⑥ Every algo should take zero or more i/p.
⑦ Algorithm should be independent of programming language.

Non-Deterministic Algo:- Random no generator

Analysis Of Algorithm

Note:- How to check in the available algorithms which is the best one?

① Time ② Space

How to find time complexity of the algo?

T(A) :- time complexity of the algorithm A

T(A) : C(A) + R(A)

C(A) — compile time of A
- It depends on the compiler)
- S/W

R(A) — Run-time of A
- (It depends on the processor)
- H/W

★ [n] is considered to be very large

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Analysis

Apostiary	Apriori
① It is programming of a compiler & h/w dependent analysis. ← depends upon compiler & processor.	① It is independent of programming language & h/w.
② It is dependent, the complexity/ answer changes from comp. to computer.	② The answer will not change. (i.e. the complexity doesn't change from computer to computer.)
③ Exact answer.	③ Approximated answer

Apriori Analysis:-

is the determination of order of magnitude of a statement

```
Main()
{ int x,y,z; ——→ order of magnitude of this statement is 1.
  x=y+z;       (which means that this statement is executed once when the program executes.)
}
```

```
Main()
{ int x,y,z,i,n;      → order of magnitude=1
  x=y+z;              → order of magnitude=1
  for(i=1;i<=n;i++)   → order of magnitude= n
    {x=y+z;
    }
}
```

No of statements & steps = 3

$n+2 = O(n)$

```
Main()
{ int x,y,z,i,j,n;    → order of magnitude=1
  x=y+z;              → " " " =1
  for(i=1;i<=n;i++)   →
    x=y+z;            → " " " = n
  for(j=1;j<=n;j++)
  {for(j=1;j<=n;j++)  } → order of magnitude = n²
    {x=y+z;
    }
  }
}
```

$n+2+n^2 = O(n^2)$

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④ Main()
{ int x,y,z,i,j,n; → order of magnitude = 1
x = y+z; → " " " = 1
for(i=1 to n) → " " " = n
x=y+z;
for(i=1 to n) → " " " = n
for(j=1 to n/2) → " " " = n/2 } → $n^2/2$
x=y+z;
}

$\frac{n^2}{2} + n \Rightarrow \frac{1}{2}n^2 = O(n^2)$ → as n is c
therefore $\frac{n^2}{2}$

★

⑤ Main()
{ int x,y,z,n; → order of magnitude = 1
x=y+z; → " " " = 1
while(n>1)
{ x=y+z
n=n/2 }
}

if n=16	if n=8	if n=32
then loop will be executed 4 times.	then loop will be executed 3 times.	then loop will be executed 5 times.

$\log_2 n$

∴ 2($\log_2 n$) statements will be executed.
& $2\log_2 n + 2$
but we can neglect +2 & multiple 2
∴ complexity is $\log_2 n$

but if n=n/3	but if n=n-1
then $\log_3 n$	then O(n)
if n=n/10	if n=n-2
then $\log_{10} n$	then $\frac{n}{2}$, & O(n)

⑥ Main()
{ int x,y,z,i,j,k,n;
for(i=1 to n) → order of magnitude = n
{ for(j=1 to i) → 1, 2, 3, . . . , n
{ for(k=1 to 135) → order of magnitude = 135
{ x=y+z;
}
}
}

$135 \times n\left[\frac{n(n+1)}{2}\right] = \frac{135 \times n^2(n+1)}{2}$

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i = 1	i = 2	i = 3	. . .	i = n
j = 1 to 1	j = 1 to 2	j = 1, 2, 3		j = 1, 2, 3, . . . , n
⇓	⇓	⇓		⇓
135	135 135	135 135 135		135 135 135 135

$T = 1\times135 + 2\times135 + \ldots + n\times135$

$= 135[1+2+\ldots+n]$

$= 135\,\frac{n(n+1)}{2} \Rightarrow O(n^2)$

```
Main()
{ int x,y,z,i,j,k,n;
  for(i= 1 to n)
   { for j=1 to i^2)
      { for (k=1 to n/2)
        { x=y+z;
        }
      }
   }
}
```

i = 1	i = 2	i = 3
j = 1 to 1	j = 1 to 4	j = 1 to 9
k = 1 to n/2	k = 1 to n/2	k = 1 to n/2

$T = 1\times\frac{n}{2} + 4\times\frac{n}{2} + 9\times\frac{n}{2} + \ldots + n^2\times\frac{n}{2}$

$= \frac{n}{2}[1 + 4 + 9 + \ldots + n^2] = \frac{n}{2}\left[\frac{n(n+1)(2n+1)}{6}\right]$

$\Rightarrow O(n^4)$

```
Main()
{ n= 2^2^k;
  for (i=1; i<=n; i++)
   { j= 2;
     while (j<=n)
      { j= j^2;
      }
   }
}
```

i = 1	i = 2	. . .	i = n
j = 2	j = 2		j = 2
$2 \le n$	$\log_2 n$		$\log_2 n$
$\log_2 n$			

$n = 2^{2^k}$

$\log_2 n = 2^k$

$\log_2(\log_2 n) = k$

$n \times (k+1)$

$n \times [\log_2 \log_2 n] + n$

$\Rightarrow O(n \log_2 \log_2 n)$

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k=1	k=2
n=4	n=16
2 ≤ 4	2 ≤ 16
4 ≤ 4	4 ≤ 16
16 ≤ 4	16 ≤ 16
2	256 ≤ 16
	3

k+1

O(1):- it will take same time always.

① A(n)
{ if (n≤1) return;
else
return(A(√n));
}

↳ O(log log₂ n)

A(1000)
↳ A(33) → 4 times
↳ A(6)
↳ A(3)
↳ A(1)

n=1	n=2	n=3
1	A(√2)	A(√3)

{a=√n;
b=A(a);
return b;
}

(a) O(1)
(b) O(log n)
(c) O(log log n)
(d) none

$T(n) = \begin{cases} O(1), & n = 2^{2^k} = k \\ T(n^{1/2}) + 2O(1) & \text{otherwise} \end{cases}$

put $n = 2^k$

$T(2^k) = T(2^{k/2}) + O(1)$

$S(k) = S(k/2) + O(1)$

$k^{\log_b a} = k^{\log_2 1} = k^0 = O(1)$

$\therefore O(k^{\log_b a} \log k)$

$= O(\log k) = O(\log \log n)$

② i/p: An array of n-elements in the range [1 ..... n]
o/p: print repeated elements.
What is the time complexity?
→ The best algo as:-

① flag [1 2 3 4 5 6 7] ⇒ O(1)

② for (i=1 to n)
flag[i] = 0 ⟶ n

③ for (i=1 to n)
{ if (flag[a[i]] != 0)
flag[a[i]] = 1
else
printf [a[i]];
}

⟶ n

2n ⇒ O(n)

Space complexity increases in this case

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Multiplication

$A = \begin{bmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \end{bmatrix}_{m \times n}$ $B = \begin{bmatrix} 1 & 4 \\ 2 & 5 \\ 3 & 6 \end{bmatrix}_{n \times p}$

$C = A \times B = \begin{bmatrix} \bullet & \bullet \\ \bullet & \bullet \end{bmatrix}_{m \times p}$ each element will take n multiplications to compute.

∴ there are $m \times p$ elements

∴ total no. of multiplications = $m \times n \times p$

∴ $O(mnp)$

) Addition of 2 $n \times n$ matrices require order of $O(n^2)$.
) Multiplication of 2 $n \times n$ matrices require $O(n^3)$.

Asymptotic Notation

- Big-oh-Notation (O)
- Omega-Notation (Ω)
- Theta-Notation (θ)

Big-oh-Notation

Let $f(n)$ & $g(n)$ be 2 +ve functions.

final value of the answer will be +ve

$n^2 + 10 \rightarrow$ +ve fn

$n^2 - 10 \rightarrow$ +ve fn (because n is very large value)

$10 - n^2 \rightarrow$ -ve fn

(n) = O(g(n))

(n) is order of g(n) iff

$f(n) \leq c\, g(n)$ $\forall n, n \geq n_0$

where c & n_0 are two constants & $c > 0$ & $n_0 \geq 1$

(n) = n

(n) = n^2

$f(n) \leq c\, g(n)$ $\forall n, n \geq n_0$

$n \leq n^2$, $n \geq 1$

(n) = $n^2 + 10$

g(n) = n^2

(n) = O(g(n))

$n^2 + 10 \leq c \cdot n^2$, $n \geq n_0$

1 , 1

2 , 4

C :- Speed of the processor.

n :- no. of inputs as no. of I/Ps can't be fractions, ∴ $n_0 \geq 1$

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- $f(n) = n^2$
- $g(n) = n$

$f(n) \leq C \cdot g(n)$

. there is no C for these functions

$\therefore f(n) \neq O(g(n))$

☆ $g(n)$ is greater by factor 'c' than $f(n)$.

Omega - Notation (Ω)

$f(n) = \Omega(g(n))$

or

$f(n)$ is omega of $g(n)$ iff

$f(n) \geq C g(n) \quad \forall n, n \geq n_0$

where c & n_0 are 2 constants, $C > 0$ & $n_0 \geq 1$

- $f(n) = n^2 - 10$
  $g(n) = n^2$
  $f(n) = \Omega(g(n))$
  $n^2 - 10 \geq C \cdot n^2$ ($n_0 \geq 5$)
  $C \leq 1/2$

$g(n)$ is smaller than $f(n)$ by 'c' times.

- $f(n) = n^2$
  $g(n) = n$
  $n^2 \geq C \cdot n$, $n_0 \geq 1$.
  $C \leq 1$

- $f(n) = n$
  $g(n) = n^2$
  $n \geq C \cdot n^2$
  there is no such c.
  $\therefore f(n) \neq \Omega(n^2)$

Theta - Notation

$f(n) = \theta(g(n))$

or

$f(n)$ is theta of $g(n)$ iff

① $f(n)$ is Order of $g(n)$

$f(n) \leq C_1 g(n)$ &

② $f(n) = \Omega(g(n))$

ie $f(n) \geq C_2 g(n)$

such that there exist 2 constants $C_1 C_2$ & n_0 where

e.g.

$f(n) = n^2 + 10$

$g(n) = n^2$

① $n^2 + 10 \leq C_1 n^2$

$C_1 = 2$

② $n^2 + 10 \geq C_2 n^2$

$C_2 = 1$

$n^2 + 10 = \theta(n^2)$

※ [the value of n_0 sho
he same for both O
$\Omega(n)$?

$f(n) = n^2$
$g(n) = n^2$
$n^2 = O(n^2)$
$n^2 = \Omega(n^2)$
$n^2 = \theta(n^2)$

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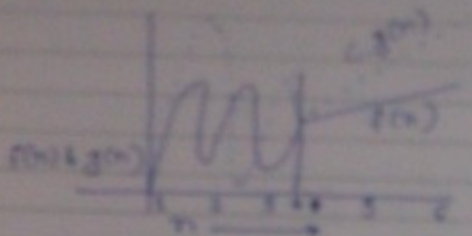
Big-Oh Notation



Worst Case

If $T(n) = O(n^2)$

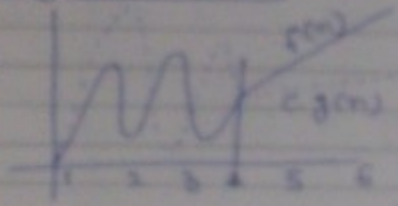
this means that whatever time is given to algo A it will not take more than n^2 time.

$\therefore$ $n = O(n^2)$
$n^2 = O(n^2)$
$n^3 = O(n^2)$

$f(n) = O(g(n))$
$f(n) \leq c \cdot g(n)$ $\forall n, n \geq n_0$

here, $g(n)$ is upperbound on all the values of $f(n)$.

Omega Notation

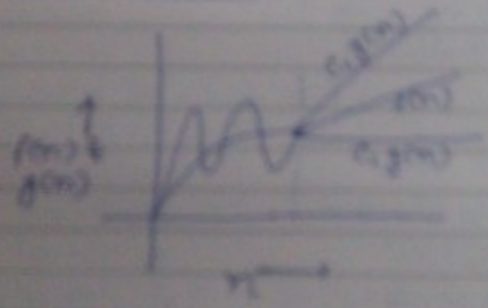


$f(n) = \Omega(g(n))$
$f(n) \geq c \cdot g(n)$

Ex.
$g(n)$ is smaller than $f(n)$ by c times

Best Case:
$T(n) = \Omega(n^2)$
this means that whatever time is given to algo A it will not take more than n^2 time.
$\therefore$ $n^4 = \Omega(n^2)$
$n^3 = \Omega(n^2)$

Theta Notation



$f(n) = \Theta(g(n))$
$f(n) \leq c_1 g(n)$ → [Big-Oh Notation]
$f(n) \geq c_2 g(n)$ → [Omega Notation]

c_1 & c_2 may be different or may be same.

We can use any symbol for average time complexity.

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*If there are 2 algo A & B.
$T(A) = O(T(B)) \longrightarrow$ A is better
$T(A) = \Omega(T(B)) \longrightarrow$ B is better
$T(A) = \theta(T(B)) \longrightarrow$ Both are equivalent.

Conclusion :-

Big-oh-Notation (O)
$n^2 = O(n^2) \longrightarrow$ Tightest upper bound [the upper bound which is exactly equal to the f(n).
$n^2 = O(n^3)$
$n^2 \neq O(n)$
$n^2 = O(n^{100})$
$A = O(B)$ here B is upper bound, where B may or may not be tight upper bound.

Small-oh-Notation :-

$n^2 = o(n^3)$
$n^2 = o(n^{100})$
$n^2 \neq o(n^2)$

$f(n) = o(g(n))$
$f(n) < cg(n), \forall n, n \geq n_0$
such that $c > 0$ & $n_0 \geq 1$.

$A = o(B)$, here B is upper bound, where B is not tightest upper bound.

Big-Omega Notation (Ω) :-

$n^2 = \Omega(n^2) \longrightarrow$ tightest lower bound.
$n^2 = \Omega(n)$
$n^2 = \Omega(1)$
$A = \Omega(B)$, here B is lower bound which may or may not tightest lower bound.

Small Omega Notation (ω) :-

$n^2 \neq \omega(n^2)$
$n^2 = \omega(n)$
$n^2 = \omega(1)$

$f(n) = \omega(g(n))$
$f(n) > cg(n)$
such that $c > 0$ & $n \geq n_0$
& $\boxed{n_0 \geq 1}$

$A = \omega(B)$, here B is lower bound which is not tightest lower bound.

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Theta-Notation (Θ)

$n^2 = \Theta(n^2)$ ⟶ Both Tightest Upper bound & tightest lower bound.
$n^2 \neq \Theta(n)$
$n^2 \neq \Theta(n^3)$
$A = \Theta(B)$ → B is both Tightest Upper bound & tightest lower bound.

Types Of Time Complexity :-

(1) Constant time complexity ⇒ $O(1)$
(2) Logarithmic time complexity ⇒ $O(\lg n)$ — $O(\lg\lg n)$ [in b/w] — $O(\ln n)$ [in b/w]
(3) Linear " " ⇒ $O(n)$ — $O(n\lg n)$
(4) Quadratic " " ⇒ $O(n^2)$ — $O(n^2\lg n)$
(5) Cubic " " ⇒ $O(n^3)$ — $O(n^{3+\varepsilon})$
(6) Polynomial " " ⇒ $O(n^k)$ $k \geq 1$
(7) Exponential " " ⇒ $O(k^n)$ $k \geq 2$

$n > (\log n)^2$ $\quad 2^n < n^n$
$n > (\log n)^{100}$ — value of n is so large that constant powers don't overpower them
$n < (\log n)^n$

[$n > \log n$ & $n > (\log n)^c$ where c is a constant, by $n < (\log n)^n$

---

① $T = \sum_{i=1}^{n} i = 1+2+3+4+\ldots+n = \frac{n(n+1)}{2} = O(n^2) = \Omega(n^2)$

② $T = \sum_{i=1}^{n} i^2 = 1^2+2^2+3^2+\ldots+n^2 = \frac{n(n+1)(2n+1)}{6} = \boxed{O(n^3)} = \Omega(n^3)$

★ this means that the sum value can't cross n^3.

→ To add the squares of n natural nos, it will take $O(n)$, i.e. only one for loop.

③ $T = n \times n \times n \times \ldots \times n$ [n times]
$T = O(n^n)$ → this means T value can't exceed n^n.
It will take $O(n)$ to compute T.

(iii) $T = n!$

$= n \times (n-1) \times (n-2) \times \dots \times 3 \times 2 \times 1$

$\leq n \times n \times n \times \dots \times n \times n \times n$

$= O(n^n) \rightarrow n!$ can't exceed n^n.

→ It will take $O(n)$ time to solve it, i.e. to calculate factorial because it will take only one for loop.

(iv) $T = \sum_{i=1}^{n} x^i$

$= x + x^2 + x^3 + \dots + x^n$

almost equal $= \frac{x(x^n - 1)}{(x-1)} = O(x^n)$ → it will can't exceed x^n with appropriate value of c.

① $f(n) = n^2$
$g(n) = 2^n$
$n^2 \leq c\,2^n$
$\therefore f(n) = O(g(n))$
$c = 1$ & $n_0 = 4$

② $f(n) = 2^n$
$g(n) = n^n$

n	2^n	n^n
1	2	1
2	4	4
3	8	27
4	16	256

$\therefore f(n) = O(g(n))$
$c = 1$ & $n_0 = 2$
or $2^n = O(n^n)$
or $n^n = \Omega(2^n)$

③ $f(n) = n^2 \log n$
$g(n) = n(\log n)^{10}$

$n^2 \log n \quad n(\log n)$
$n \cdot n \log n \quad n \log n$
$n > (\log n)$
$\therefore f(n) > g(n)$
$\therefore f(n) = \Omega(g(n$

for very large values of n
e.g. $n = 2^{1000}$
$(\log n)^4 = (100$
$2^{1000} >>>> (1$

(1) Which statement is false?

(a) $n^2 2^{3\log_2 n} = \Theta(n^5)$

(b) $\frac{4^n}{2^n} = \Theta(2^n)$

(c) $2^{\log_2 n} = \Theta(n^2)$

(d) None

(a) $n^2 2^{3\log_2 n}$
$n^2 2^{\log_2 n^3}$
n^5

n^5
n^5

$n^{\log_b a} = a^{\log_b n}$

(2) Which statement is T/F?

(a) $100\, n \log n = O\left(\frac{n \log n}{100}\right)$ → T

★ in such problems remove the constants.

$100 \cdot n \log n \leq c \cdot \frac{n \log n}{100}$

$(n+k)^m \neq n^m + k^m$

(b) $\sqrt{\log n} = O(\log\log n) \rightarrow F$ $\sqrt{\log n} \not\leq c \log\log n$

(c) $0 < x < y$ then $n^x = O(n^y) \rightarrow T$ $(\log n)^{1/2} \not\leq c \log\log n$

$\frac{1}{2}\log\log n \not\leq c \log\log n$

(d) $2^n \neq O(n^k)$ where k is constant $\rightarrow T$

$2^n \not\leq c \cdot n^k$

★ exponential is always faster than the polynomial.

④ T/F = ?

(a) $(n+k)^m = O(n^m)$ where k is constant $\rightarrow T$

(b) $2^{n+1} = O(2^n) \rightarrow T$

(c) $2^{2n} = O(2^n) \rightarrow F$

(b) $2^{n+1} \leq c \cdot 2^n$

(c) $2^{2n} \leq c\, 2^n$

$2 \cdot 2^n \leq 3 \cdot 2^n$

(a) $(n+k)^m \leq c \cdot n^m$

$n^m + {}^mC_1 n^{m-1}k + {}^mC_2 n^{m-2} + \ldots + k^m$

$\leq n^m + k^m$

$\leq n^m + n^m$

$= O(n^m)$

★ if K is not constant, then $O(n^m)$ & $O(k^m)$ can both be the answer.

★ constant is always smaller than the function.

⑤ Consider the following two functions :-

$f(n) = n^3 \quad 0 \leq n < 10,000$

$= n^2 \quad n \geq 10,000$

$g(n) = n \quad 0 \leq n < 100$

$n^3 \quad n \geq 100$

• for n b/w 1 & 100

$g(n) = O(f(n))$ as $f(n) = \Omega(g(n))$

• for n b/w 101 & 10,000

• for $n \geq 10,000$

$f(n) = O(g(n))$

$(n+k)^m \neq n^m + k^m$

(b) $\sqrt{\log n} = O(\log\log n) \rightarrow F$ $\sqrt{\log n} \leq c \log\log n$
$(\log n)^{1/2} \leq c \log\log n$
$\frac{1}{2}\log\log n \leq c \log\log n$

(c) $0 < x < y$ then
$n^x = O(n^y) \rightarrow T$

(d) $2^n \neq O(n^k)$ where k is constant $\rightarrow T$

$2^n \leq c \cdot n^k$

★ exponential is always faster than the polynomial.

⑥ T/F = ?

(a) $(n+k)^m = O(n^m)$ where k is constant $\rightarrow T$
(b) $2^{n+1} = O(2^n) \rightarrow T$
(c) $2^{2n} = O(2^n) \rightarrow F$

(b) $2^{n+1} \leq c \cdot 2^n$
(c) $2^{2n} \leq c \cdot 2^n$
$2 \cdot 2^n \leq 3 \cdot 2^n$

(a) $(n+k)^m \leq c \cdot n^m$
$n^m + {}^mC_1 n^{m-1}k + {}^mC_2 n^{m-2} + \ldots + k^m$
$\leq n^m + k^m$
$\leq n^m + n^m$
$= O(n^m)$

★ if k is not constant, then $O(n^m)$ & $O(k^m)$ can both be the answer.
★ constant is always smaller than the function.

Q Consider the following two functions:-

$f(n) = n^3 \quad 0 \leq n < 10,000$
$= n^2 \quad n \geq 10,000$

$g(n) = n \quad 0 \leq n < 100$
$= n^3 \quad n \geq 100$

• for n b/w 1 & 100
$n^3 \quad n$
$g(n) = O(f(n))$ or $f(n) = \Omega(g(n))$

• for $n \geq 10,000$
$n^2 \quad n^3$
$f(n) = O(g(n))$

• for n b/w 101 & 10,000
$n^3 \quad n^3$

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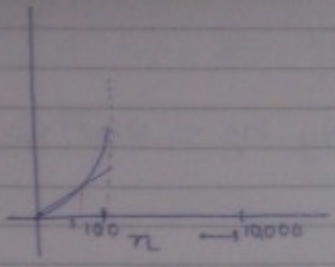
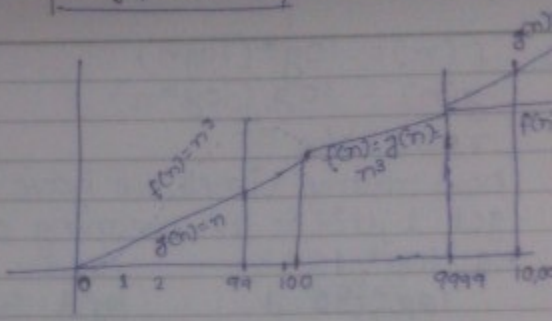
$n! = O(n^n)$
$n! \neq \omega(n^n)$

Graph





Q. Consider the following fn. :-

$f_1 = 2^n$
$f_2 = n^{3/2}$
$f_3 = n \log n$
$f_4 = n^{\log n}$

arrange above fn. in increasing order.

Ans $n \log n < n^{3/2} < n^{\log n} < 2^n$

$n^{\log n} > n^{3/2}$ because fn. is always greater than constant.

2^n $n \log 2$

$n^{\log n}$ $\log n \times \log n$

~~$n \log n$~~ ~~$n^{1/2}$~~

$\sqrt{n} > \log n$

$\frac{1}{2} \log n > \log \log n$

Q $f_1 = (\log n)!$
$f_2 = \log(n^n)$
$f_3 = (\log n)^{\log n}$

$f_1 = x!$
$f_2 = \log(n!) \leq \log(n^n) = n \log n = nx$
$f_3 = x^x$

as $nx > x!$
$\therefore f_2 < nx$
$\therefore nx > f_2 < x! > f$

$(\log n)! < \log(n!)$

$(\log n)! \quad \log$

$(\log n)^{\log n} \quad \log(n!)$
$\log n \log(\log n) \quad \log(n!)$

$f_2 < f_1 < f_3$

Cosmos

Q. $f(n) = \log^*(\log n)$
$g(n) = \log(\log^* n)$
rela blw ?
how many times we have to apply log so that we can get 1, this is the meaning of *.
e.g $\log^*_2 16 = 3$
$\log^*_2 2^{16} = 4$

$f(n) = \log^*(\log 65,536) = 4$
$g(n) = \log(\log^* 65,536) = \log_2 5 = 2$
$f(n) > g(n)$

Properties Of Asymptotic Notation :-

① Reflexive
Reflexive property
(i) $f(n) = O(f(n))$
(ii) $f(n) = \Omega(f(n))$
(iii) $f(n) = \Theta(f(n))$

② Symmetric Property
(i) If $f(n) = O(g(n))$
then $g(n) \neq O(f(n))$
(ii) If $f(n) = \Omega(g(n))$
then $g(n) \neq \Omega(f(n))$
(iii) If $f(n) = \Theta(g(n))$
then $g(n) = \Theta(f(n))$

③ Transitive Property
(i) If $f(n) = O(g(n))$ & $g(n) = O(h(n))$
then $f(n) = O(h(n))$
(ii) If $f(n) = \Omega(g(n))$ & $g(n) = \Omega(h(n))$
then $f(n) = \Omega(h(n))$
(iii) If $f(n) = \Theta(g(n))$ & $g(n) = \Theta(h(n))$ then $f(n) = \Theta(h(n))$

④ If $f(n) = O(g(n))$ then $2^{f(n)} \neq O(2^{g(n)})$
because If $f(n) = 2n$ & $g(n) = n$
$2n = O(n)$
but $2^{2n} \neq O(2^n)$

⑤ $f(n) \neq O(f(\frac{n}{2}))$
because take $f(n) = 2^{2n}$

Cosmos

⑥ $f(n) \neq O((f(n))^2)$

because take $f(n) = \frac{1}{n}$

$\frac{1}{n} > \frac{1}{n^2}$

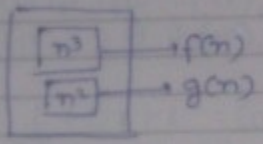
⑦ If $f(n) = O(g(n))$

$d(n) = O(h(n))$

then

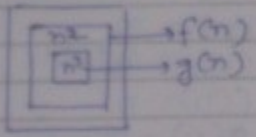
(i) $f(n) + d(n) = \max\{g(n), h(n)\}$ → $O(g(n)+h(n))$

(ii) $f(n) \cdot d(n) = O(g(n) \cdot h(n))$



both are executed individually

nested for loops



⑧ If $f(n) = O(g(n))$

then

$h(n) \cdot f(n) = O(h(n) \cdot g(n))$

e.g.

$3 < 5$

$6 \times 3 < 6 \times 5$

★ h(n) should be +ve function.

Q Consider following 4 fn :-

$f(n) = O(g(n))$ $g(n) = O(h(n))$ & $g(n) \neq O(f(n))$ & $h(n) = O(g(n))$

True?

(a) $f(n) + g(n) = O(g(n))$ → T

(b) $f(n) = O(h(n))$ → T (by transitivity)

(c) $h(n) = O(f(n))$ → F

(d) $h(n) \cdot f(n) = O(h(n) \cdot h(n))$ → T

(d) $f(n) = O(h(n))$

$h(n) \cdot f(n) = O(h(n) \cdot h(n))$

(a) By ~~OR~~ $f(n) = O(g(n))$

$\therefore f(n) \leq g(n)$

$\therefore f(n) + g(n) = \max$

but $g(n) \neq O(f($

$\therefore g(n) \not\leq f($

$\therefore f(n) < g(n)$

$\therefore f(n) + g(n) = O($

Cosmos

$f(n) = n^3$

Q. Suppose $T_1(n) = O(f(n))$ & $T_2(n) = O(f(n))$ then

$T_1(n) \le C_1 \cdot f(n)$
$T_2(n) \le C_2 \cdot f(n)$

(a) $T_1(n) + T_2(n) = O(f(n)) \longrightarrow T$
(b) $T_1(n) = O(T_2(n)) \longrightarrow F$
(c) $T_1(n) = \Omega(T_2(n)) \longrightarrow F$
(d) $T_1(n) = \Theta(T_2(n)) \longrightarrow F$

→ Nothing can be said, as we dont know the rem b/w $T_1(n)$ & $T_2(n)$
$T_1(n)$ can be n^5, $T_2(n)$ can be n^4
$f(n)$ can be n^{10} & $T_1(n)$ can also be n^4

Q.(i) $f(n) = 1 + 10 + \frac{1}{n^2} = 1 + 10 + 0 = 11 \Rightarrow O(1)$ [for large nos. of n : $\frac{1}{n^2} \to 0$]

(ii) $6n^3 2^n + n + \log n \Rightarrow O(n^3 2^n)$

(iii) $f(n) = n$
$g(n) = n^{1+\sin n}$
→ These two fn are uncomparable.

$f(n) = n$
$g(n) = n \cdot n^{\sin n}$

n | $n^{\sin n}$
1 | $n^{\sin n}$

n	$f(n)$	$g(n)$
0	0	0
90	90	90^2
180	180	180
270	270	1
360	360	360

This repeats for every 360° cycle.

(iv) $f(n) = n^{\sin n}$
$g(n) = n^{\cos n}$
→ These two fn. are uncomparable.

n	$f(n)$	$g(n)$
0	0	0
30	$30^{1/2}$	$30^{\sqrt{3}/2}$
45	$45^{1/\sqrt{2}}$	$45^{1/\sqrt{2}}$
60	$60^{\sqrt{3}/2}$	$60^{1/2}$
90	90	1

Date 09.06.12

Cosmos

★ Problem is small when no. of element is only 1.

Binary Search

```
a[i] = 10  20  30  40  50  60  70
i    =  1   2   3   4   5   6   7
```

array → no. to search

BS(a, i, j, x)

first element → last element

```
if (i==j)
 { if (a[i]==x)  ——→ O(1)
      return i;
   else
      return (-1);
 }
else
 { mid = (i+j)/2;  ——→ O(1)
   if (a[mid]==x) return(mid); ——→ O(1)
   else
    if (a[mid]<x)  ——→ O(1)
      BS(a, mid+1, j, x);
    else
      BS(a, i, mid-1, x);
 }
}
```

$$T(n) = \begin{cases} O(1), & \text{if } n=1 \\ O(1)+O(1)+O(1)+T\left(\frac{n}{2}\right) & \text{otherwise} \\ = T\left(\frac{n}{2}\right) + c \end{cases}$$

$$T(n) = T\left(\frac{n}{2}\right) + c$$
$$= T\left(\frac{n}{4}\right) + c + c$$
$$= T\left(\frac{n}{8}\right) + c + c + c$$
$$= T\left(\frac{n}{2^k}\right) + kO(1)$$

$n = 2^k$

$\log_2 n = k$.

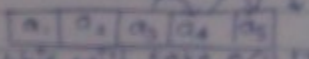
Cosmos

Q1 I/p: A sorted array of n elements
o/p: Find 2 elements a & b such that a+b>1000

func(a,i,j)
{

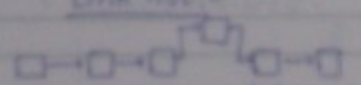
a[i] →	1000	2000	3000	4000	5000	6000	7000	8000	9000
i →	1	2	3	4	5	6	7	8	9

Array:-



- this will take O(1) to search the element after which we want to add
- this will take O(n) to add the element after the given element.

Link list:-



- this will take O(n) to search the address of the node of it on which we want to add the element
- this will take O(1) to add the element

& → reference operator
* → dereference operator

Solution:- take the last two elements & add them, if their sum > 1000, then it is the answer, if they don't give, then no other combination will give as it is the biggest combination.

algo:- ① goto last 2 element (a,b) → O(1) [as it an array]
② if (a+b>1000) return (a,b); → O(1) [one comparison to check greater than 1000]
else
return (-1);

∴ only one comparison
∴ complexity = O(1) + O(1) = O(1).

Cosmos

Q2. similar to Q.1.
but only a+b==1000.

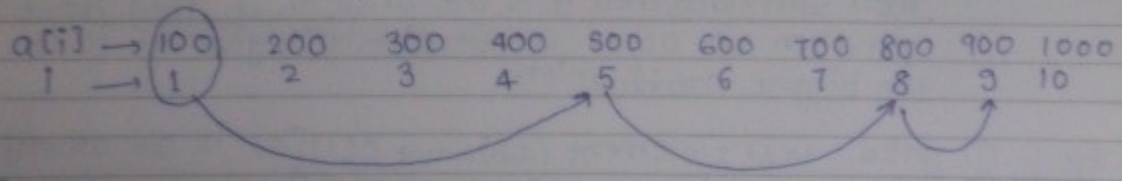
```
for(i=1; i<=n; i++)
  { for(j=i+1; j<=n; j++)
      { if(a[i]+a[j]==1000])
          return(a[i], a[j]);
        else if(a[i]+a[j]>1000)
          return(-1);
      }
  }
return(-1);
```

Time Complexity:-
$O(n^2)$

Algo:-
→ ① For the element a apply L.S. to find another element b such that a+b==1000. → O(n)
② continue for all elements until a+b==1000. → $O(n^2)$

$O(n^2)$

Improvement using Binary Search:-



① take a
② search b for ~~ea~~ every a which will take O(lgn)
there are n a's.
∴ Complexity = O $n \times \log_2 n$ → $O(n \log_2 n)$

Cosmos

3rd Improvement:-

a+b = 1300

a[i] →	100	200	300	400	500	600	700	800
i →	1	2	3	4	5	6	7	8
	a	a						b

a[1] + a[8] = 900 < 1300
increase 1 to 2
∴ a[2] + a[8] = 1000 < 1300
increase 2 to 3
a[3] + a[8] = 1100 < 1300
increase 3 to 4
a[4] + a[8] = 1200 < 1300
increase 4 to 5
a[5] + a[8] = 1300 ⟶ a = 500, b = 800
∴ it will take only one scan,
∴ complexity (time) = O(n)

→ Greedy Technique.

Algo:-

① Take 2-variables A & B.
A → 1st element
B → last element.

② If (A+B == C)
return (A,B);
else
if (A+B > C)
B = B-1;
else
A = A+1;

③ continue 2nd step until A+B == C.
But if there are no elements with A+B == C, then stop when i == j, (i.e. when the position of A & B are equal.

Cosmos

Q. i/p:- A sorted array of n-distinct elements both +ve & -ve.

o/p:- find an element b such that A[i]==i.

Algo:-

My practice:-

① Go to the middle element & check whether it is +ve or -ve.

② If

```
func(a,i,j){
for(i=1; i<=n; i++)
  { mid = (i+j)/2 ; if(a[mid]==mid) return(mid);
if(a[mid]>0)
   i = mid+1;
else
   j = mid-1;
func(a,i,j);
}
```

Sir's Soln.:-

① Using Linear search.

```
check
for(i=1; i<=n; i++)
  {if(a[i]==i)
      return(i);
  }
```

⟶ O(n)

② Using Binary search.

Cosmos

```
mid = [i+j]/2
if (a[mid] == mid)
    return (mid);
if (a[mid-1]
```

a[i] →	-100	-50	-35	4	10	12	13	14	17	110	125	135	140
i →	1	2	3	4	5	6	7	8	9	10	11	12	13

example:-

	mid
a[i] →	14
i →	8

algo:-
if (position > value)
    go right
else
    go left

O(log n)

as we can see
8 < 14
∴ the 7th element can be 7,
∴ go left.
we can't go right because
9th element will be > 14
[∵ since distinct elements,
in worst case we will have
9th element as 15
10th " " 16
11th " " 17
∴ position can never be equal
to value.]

Q. i/p :- A sorted array of n-elements contain only 0 & 1
o/p :- which is greater, i.e. whether no. of zeroes are greater or no. of 1's are greater.

Cosmos

Algo:-

```
n0=0; n1=0;
for (i=1; i≤n; i++)
  { if (a[i]==0)
      n0++;
    else
      n1++;
  }
if (n0>n1)
  return 0;
else
  return 1;
```

→ O(n)

a[i] →	0	0	0	0	0	0	1	1	1	1
i →	1	2	3	4	5	6	7	8	9	10

① go to mid element
② check mid element, if it is 0, then check

- if n is even.

if $(\frac{n}{2})$ is 0 && $(\frac{n}{2}+1)$ is 0, then no. of 0's greater.
else if $(\frac{n}{2})$ is 0 && $(\frac{n}{2}+1)$ is 1 then no. of 0's = no. of 1's
else if $(\frac{n}{2})$ is 1 && $(\frac{n}{2}-1)$ is 1 then no. of 1's greater.

} → O(1)

- if n is odd.

if $\frac{n+1}{2}$ is 0, then no. of 0's are greater.
else no. of 1's are greater.

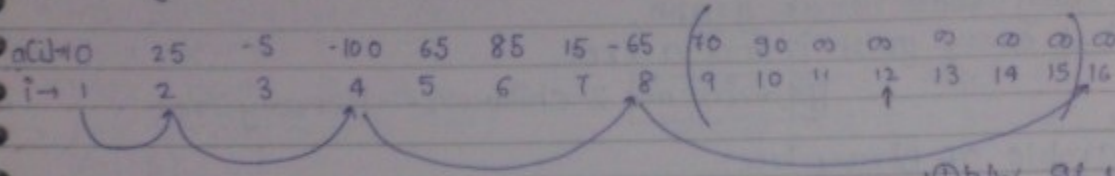
} → O(1)

Q. i/p :- An array of n elements unknown size, which contain both +ve & -ve no. until some place & afterwards all are ∞.
o/p :- Find the 1st infinite place in the array.

Cosmos

a[i] →	0	-15	25	65	-5	100	∞	∞	∞	∞
i →	1	2	3	4	5	6	7	8	9	10

① Using Linear Search → O(n).



Algo:-

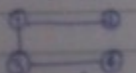
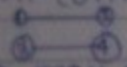
① Start with 1st element, if a[1] == ∞ then stop. else multiply the position by 2 continue till a[i] == ∞. position ie

② After reaching the ∞, then check whether the position before that ∞ is • also ∞, then apply binary search b/w $\frac{i}{2}+1$ & i-1; & check whether middle element is ∞ if it is go left otherwise go right

① b/w 9 & 15 check 12th p
if is ∞
② check 11th is then apply b/w 9 & 11.
③ check 10th is its not, then e.g. b/w 11 &

* ④ we find is ∞, which the answe

Complexity:- $O(\log_2 n)$

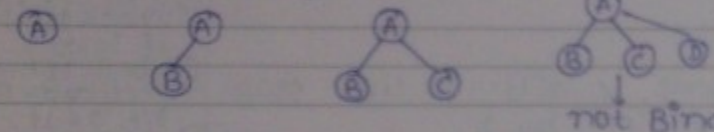
Graph	Tree
① No root element.	① root element is present.
② If there is a path b/w any two vertices, then it is a connected graph. 	② Tree is always connected.
If there is no path b/w any two vertices, then it is non-connected graph.  May or may not be connected	
③ May or may not be a directed graph.	③ It is always directed. (By default the directions are

May be cyclic or acyclic.	④ It is always acyclic.

- Every tree is a graph, but every graph is not a tree.

Maximum 2 children → Binary Tree

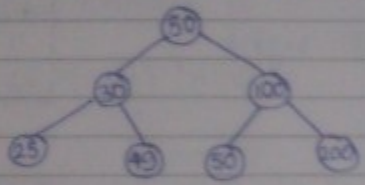
Example of Binary Tree:-



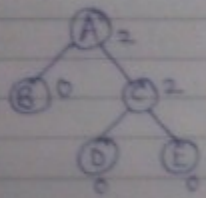
not Binary Tree.

Binary Search Tree:- In a given binary tree, at every node, root is ≥ all nodes present in left subtree & ≤ all nodes present in right subtree.

e.g.



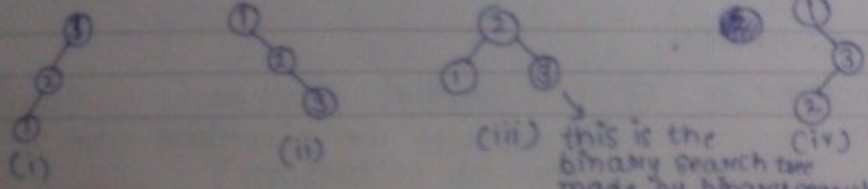
Strict Binary Tree:-



Every node can have 0 children or 2 children. [not a single child]

how many Binary Search Tree possible with 3 nodes.

1, 2, 3



Cosmos

★ No. of levels created by binary search algorithm = $\log_2 n$ [n = no. of elements]

★ height of tree = no. of levels − 1

- For Binary Tree :- (n = 5)
  → Max. levels = 5
  → Max. height = 5 − 1 = 4
- For Binary Search Tree :- (n = 5)
  → no. of levels (max.) = 5
  → Max. height = 4

- For Balanced Binary Search Tree (n = 5)
  → no. of levels = $\lceil \log_2 5 \rceil = 3$
  → height = 3 − 1 = 2.

★ No. of binary search trees possible with n-elements :-

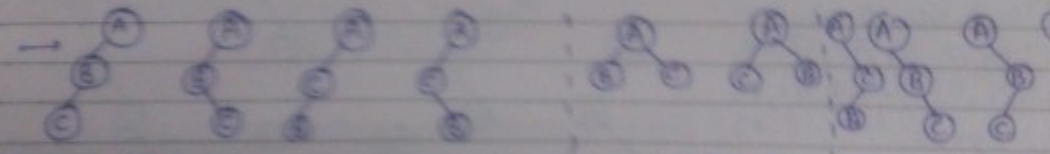
$^{2n}C_n$ / $(n+1)$

→ if n = 3

$^6C_3 = \frac{6!}{3!\,3!} = \frac{6 \times 5 \times 4}{3 \times 2} = 5 \times 4$

$\frac{^6C_3}{4} = \frac{5 \times 4}{4} = 5$

★ No. of Binary Trees with 3 nodes :-



3×3+2 3×4=12 2×3=6 3×4=12

12+12+6 = 30

★ ∴ no of binary trees with n-nodes :-

$n! \times \left(\frac{^{2n}C_n}{n+1}\right)$

Cosmos

★ No. of unlabeled Binary Trees = no. of binary search trees = $\frac{^{2n}C_n}{n+1}$.

★ No. of labeled Binary Trees = no. of binary trees × $n! \times \frac{^{2n}C_n}{n+1}$.

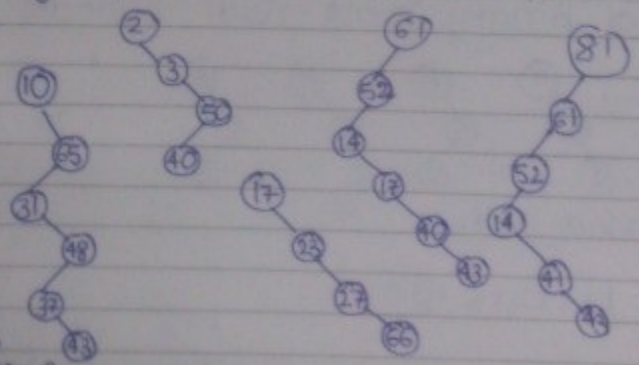
Q. A set of n distinct elements & an unlabeled binary tree with n-nodes.
In how many ways we can populate the tree with the given set ~~set~~ of elements so that it becomes BST.

(A) 1
(B) 0
(C) $n!$
(D) $\frac{^{2n}C_n}{(n+1)}$

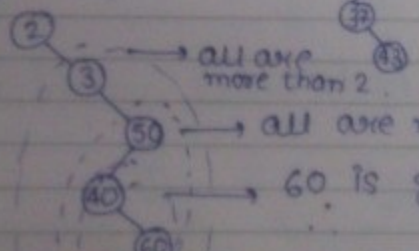
→ My answer:- because only 1 tree is given, ∴ with 1 tree, we can arrange the nodes in only 1 way so that it becomes a BST.

Q. A binary search tree is used to locate no. 43, then which of the following sequence is not possible?

(A) 61, 52, 14, 17, 40, 43
(B) 2, 3, 50, 40, 60, 43
(C) 10, 65, 31, 48, 37, 43
(D) 81, 61, 52, 14, 41, 43
(E) 17, 23, 27, 66, 18, 43



To check:-
(B)

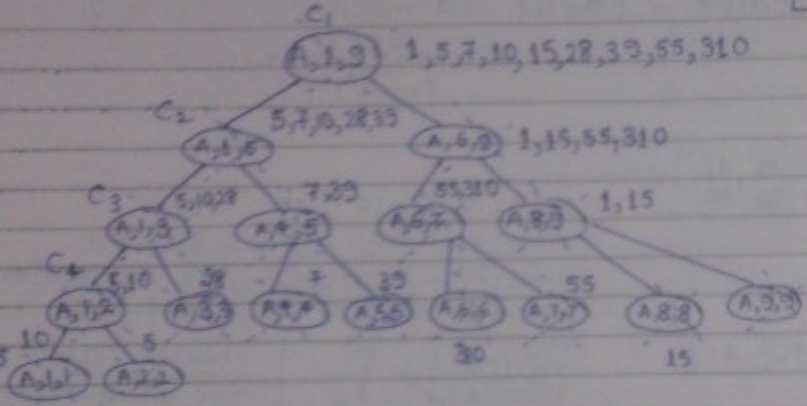


Merge Sort

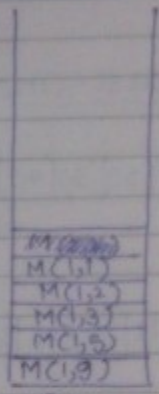
Note: Merging 2 sorted sub arrays.

ex.:-

A [10 5 28 7 39 310 55 15 1]

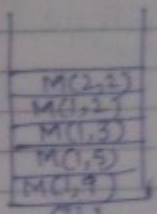


★ Max. stack size = no. of levels in the tree.
So the levels & height are calculated to measure the stack size.



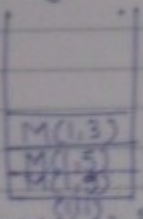
(i)

• M(1,1) is successfully completed.



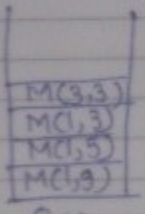
(ii)

Now M(2,2) is successfully completed.

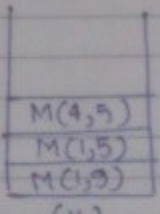


(iii)

M(1,3) is completed because its both childrens are completed.



(iv)



(v)

M(3,3) is complet
& ∴ M(1,3) is com
because its both
children are su
completed.

★ Space Complexity: i/p size + stack size. → i/p size + 5n (n - no. of each space)

★ Time complexity = sum of all function calls.

Time complexity = 17 * (function call).

↳ no. of function calls.

Merge procedure

• sorting with using the same array:- inplace sorting.

Cosmos

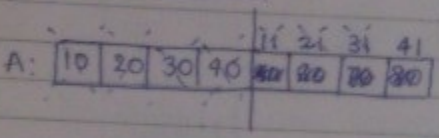
i/p to merge-procedure → M(a, i → mid, mid+1 → j)

- sorted array from i to mid
- sorted array from mid+1 to j.

o/p of merge-procedure → Sorted (a, i → j)

- sorted array from i to j.

Worst-Case :-

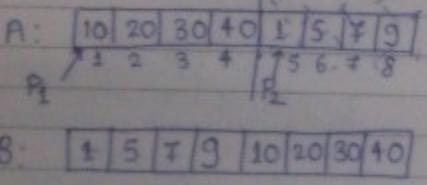


min(10,11) = 10
min(20,11) = 11
min(20,21) = 20
min(30,21) = 21
min(30,31) = 30
min(40,31) = 31
min(40,41) = 40
41

→ Total no. of comparisons 4+4-1 = 7

Worst case :-
Time complexity :-
m+n-1 → m: size of array
n: size of array
O(m+n)
When both the array size is equal :-
n/2 + n/2 - 1 ⇒ n-1 → O(n)

Best Case :-



min.(10,1) = 1
min(10,5) = 5
min(10,7) = 7
min(10,9) = 9
10
20
30
40

→ Best Case
Time complexity :-
no. of comparisons =
if 1st part of array has size m
& 2nd part of array has size n
∴ Best case complexity O[min(m,n)]
If both are of n/2 size then :- n/2 + n/2
Ω(n) → n/2 ≥ cn
it can't go n/2 ≤ cn less than this (as we have to atleast read all the n-elements :)

If we use don't use the second array B, then we have to re-sort the two array parts of A which increases the time complexity.

Cosmos

Merging Algorithm:-

```
Merge(a,i,mid, mid+1,j)
{  k= mid+1; P=i
   while(i<= mid && k<=j)
   {  if(a[i]<a[k])
      {  b[P]=a[i];
         i++;
         P++;
      }
      else
      {  b[P]=a[k];
         k++;
         P++;
      }
   }
   for(;i<=mid;i++)
   {  b[P]=a[i];
      P++;
   }
   for(;k<=j;k++)
   {  b[P]=a[k];
      P++;
   }
   for(k=1;k<=j;k++)
   {  a[k]=b[k];
   }
}
```

Merge-Sort

```
MergeSort(a,i,j)
{  if(i==j)
      return(a[i]);   ——> O(1)
   else
      Mergesort(i,(i+j)/2); ——> O(1)
```

Worst Case:-

$T(n) = n-1 = O(n)$

Best Case:-

$T(n) = \frac{n}{2} = \Omega(n)$

Average Case:-

$\theta(n)$

$$T(n) = \begin{cases} O(1), & \text{if } n=1 \\ T(\frac{n}{2}) + T(\frac{n}{2}) + \theta(n) & \text{o.w} \\ = 2T(\frac{n}{2}) + \theta(n) \\ = 2T(\frac{n}{2}) + n \end{cases}$$

Cosmos

$$T(n) = 2T\left(\frac{n}{2}\right) + n$$

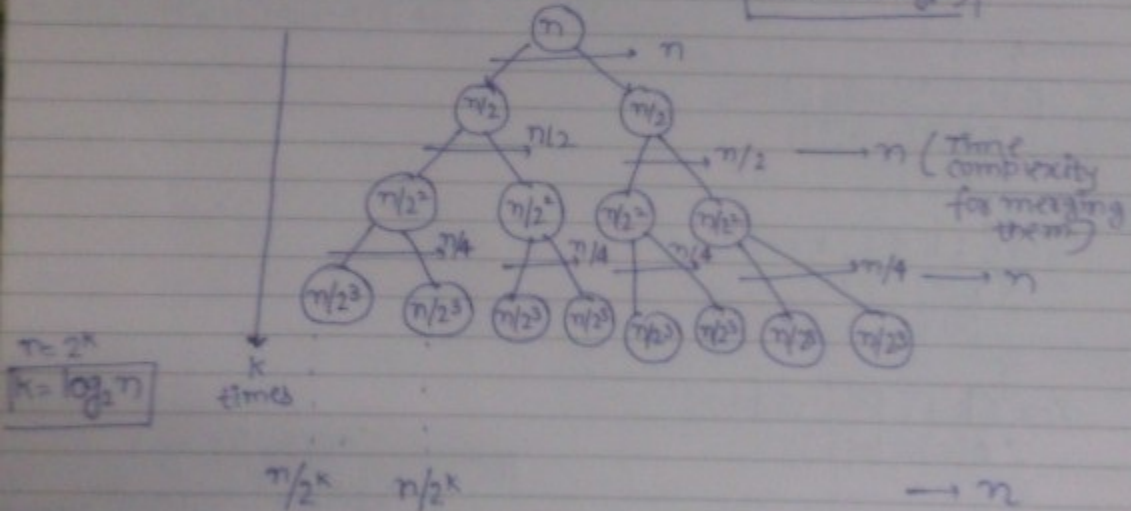
$$= 2\left[2T\left(\frac{n}{4}\right) + \frac{n}{2}\right] + n$$

$$= 2\left[2\left[2T\left(\frac{n}{8}\right) + \frac{n}{4}\right] + \frac{n}{2}\right] + n$$

$$= 2\left[2^2 T\left(\frac{n}{8}\right) + \frac{n}{2} + \frac{n}{2}\right] + n$$

$$= 2^3 T\left(\frac{n}{8}\right) + 2n + n \quad \cdots \quad 2^k T\left(\frac{n}{2^k}\right) + kn$$

$\therefore$ $T(n) = \theta(n \log_2 n)$ put $k = \log_2 n$



$\therefore$ no. of levels $= k = \log_2 n$

& at each level the time complexity $= n$.

$\therefore$ Time Complexity $= n \log_2 n$

Cosmos

Merge-Sort is not inplace sorting technique because in the merging procedure we are taking an extra array of size n.

Space complexity = $2n + 1C \log_2 n! + 2n \Rightarrow O(n)$

- $2n$: array size of merge (A) [n-elements, each element of 2B, ∴ total size = 2n B]
- $1C$: no. of variable size in each function
- $\log_2 n!$: no. of function calls from the tree of prev. page.
- $2n$: stack size array size of B.

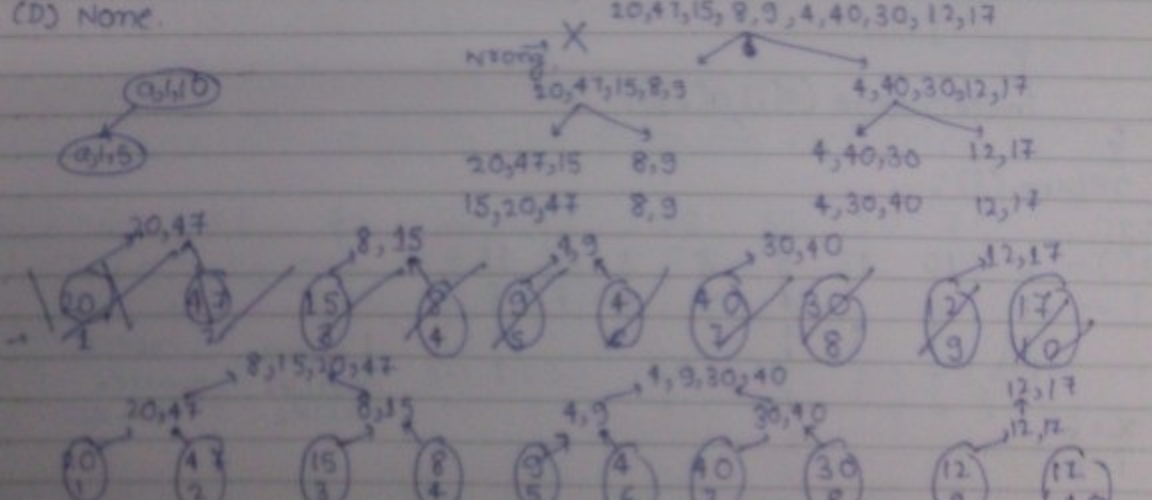
Note:- If the array size is very small, merge sort is not advisable (insertion sort is advisable), so for larger size arrays merge sort is advisable.

Q If one uses straight 2-way merge sort algo to sort following elements.

20, 47, 15, 8, 9, 4, 40, 30, 12, 17.

then the order of these elements after 2nd pass of the algo. is

(A) 8, 9, 15, 20, 47, 4, 12, 17, 30, 40
(B) 8, 15, 20, 47, 4, 9, 30, 40, 12, 17
(C) 15, 20, 47, 4, 8, 9, 12, 30, 40, 17
(D) None.



* Quick sort is used to find out the n-th smallest element in the array.

* In straight 2-way merge-sort
  - We combine & sort 1st & 2nd element, then 3rd & 4th, then 5th & 6th, & so on. [in 1st pass.]
  - In 2nd pass we merge 1st 2nd & 3rd 4th element & sort them, then 5th 6th 7th 8th, & so on.

→ In above example, no. of levels required to sort is 4.
→ ∴ no. of levels = $\log_2 n$
& each level take n time.
∴ complexity is $n\log_2 n$.

Quicksort :- [Tony Hoare]

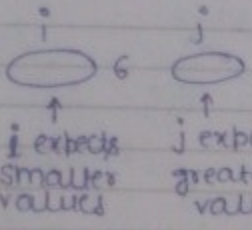
Randomized Quicksort (when pivot element is choosen randomly).
Normal Quicksort (when pivot element is choosen from same position.)

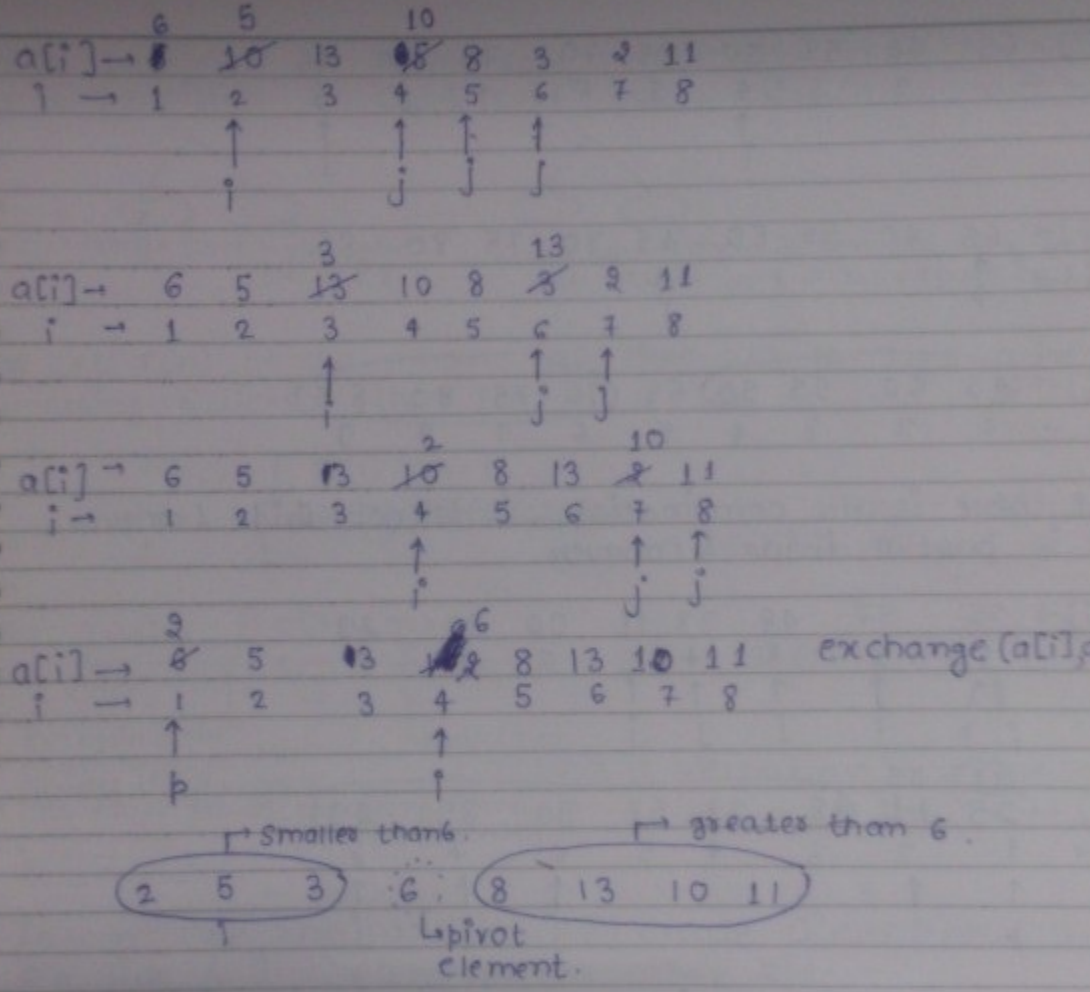
① It uses Divide & Conquer Algorithm.
② In-place sorting algo.
③ Not stable.

```
Partition(a,p,q)
{ x = a[p]; i=p;
  for(j=p+1; j<=q; j++)
    { if(a[j]<=x)
        {i=i+1;
         interchange(a[i],a[j]);
        }
    }
    Interchange(a[i],a[p]);
}
return(i);
}
```

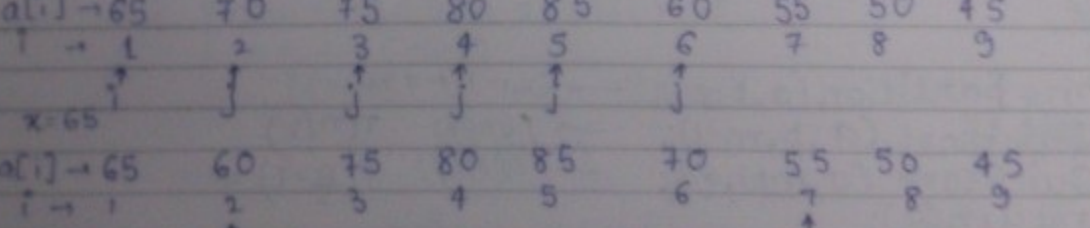

a[] →	6	10	13	5	8	3	2	11
→	1	2	3	4	5	6	7	8
	p	i	j	j	j			1
								2

x = 6
i = p;





•Apply partition algo in the following elements:-



Cosmos

a[i] → 65 60 55 50 85 70 75 80 45
i → 1 2 3 4 5 6 7 8 9
↑ i
↑ j

a[i] → 65 60 55 50 45 70 75 80 85
i → ↑ p
↑ j

a[i] → (45 60 55 50) 65 (70 75 80 85)
i → 1 2 3 4 5 6 7 8 9

★ As there is no combination, ∴ it is not divide & conquer, it is partial divide & conquer.

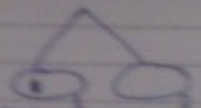
a[i] → 25 57 48 38 11 90 89 29
i → 1 2 3 4 5 6 7 8
↑ i ↑ p ↑ j ↑ j ↑ j ↑ j

11 25
a[i] → ~~25~~ ~~11~~ (48 38 57 90 89 29)
i → 1 2 3 4 5 6 7 8
↑ p ↑ i ↑ j ↑ j ↑ j

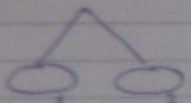
Quicksort(a,p,q) (40 70 10 20 60 50 30)
{ if (p==q)
return(a[p]);
else
{
m = partition(a,p,q); → O(n) → n
Quicksort(a,p,m-1); → T(n/2) T(n/2)
Quicksort(a,m+1,q); → T(n/2) T(n/2-1)
return(a);
}

Cosmos

* To partition in Merge sort → mid = $\lfloor \frac{l+r}{2} \rfloor$ → O(1)
* To partition in Quick sort → m = partition(a, p, q) → O(n)



Merge sort → no rel^n b/w 1 & 2



Quick sort → there is rel^n. b/w 1 & 2
all values in 2 > all values in 1.

* Let T(n) be the amount of time req. to sort n elements using Quicksort, then T(n)

$$T(n) = \begin{cases} O(1), & n=1 \\ n + T(\frac{n}{2}) + T(\frac{n}{2}-1) & \text{if } n>1. \end{cases}$$

$$T(n) = \begin{cases} O(1), & n=1 \\ 2T(\frac{n}{2}) + n, & n>1 \end{cases} \longrightarrow \Omega(n \log_2 n)$$

(Best Case) occurs when everytime the partition algo divides the elements into equal groups (containing same no of elements)

↳ partition time (this is merge time in merge sort)

Worst Case :-

In worst case the partition will create 0 elements on one side & (n-1) elements on other side.

$$T(n) = \begin{cases} O(1), & n=1 \\ n + T(0) + T(n-1) & \text{if } n>1 \\ = n + T(n-1) & \\ = T(n-1) + n & \longrightarrow O(n^2) \end{cases}$$

$n + n-1 + n-2 + \ldots$

$\leq n + n + \ldots$

$\leq n \cdot n$

$\leq n^2$

$O(n^2)$

★ When array is not at all sorted, then Quicksort is the best sort.

$T(n) = T(\frac{n}{2}) + T(\frac{n}{2}-1) + 2n$

$= 2T(\frac{n}{2}) + n$

$= n\log_2 n$

Average Case:-

$T(n) = 2T(\frac{n}{2}) + n \rightarrow$ Lucky Case

$T(n) = T(n-1) + n \rightarrow$ Unlucky Case

$T(n) = 2T(\frac{n}{2}) + n$

$= 2\left[T(\frac{n}{2}-1) + \frac{n}{2}\right] + n$ [Lucky followed by unlucky]

$= 2T(\frac{n}{2}-1) + 2n$

$\approx 2T(\frac{n}{2}) + n \rightarrow \Theta(n\log_2 n)$

↳ which is similar to Best Case

$T(n) = T(n-1) + n$

$= 2T(\frac{n-1}{2}) + n - 1 + n$

$= 2T(\frac{n-1}{2}) + 2n - 1$ → constants can't change the complexity.

$\approx 2T(\frac{n}{2}) + n$

) Quicksort is ran on 2 i/ps shown below to sort in ascending order

i) $1, 2, 3, 4, \ldots, n$

ii) $n, n-1, n-2, \ldots, 1$

et C_1 & C_2 be the no of comparisons made for p's (i) & (ii) respec., then

) $C_1 < C_2$ (B) $C_1 > C_2$ (C) $C_1 = C_2$ (D) Not comparable

★ If array is almost/already sorted then insertion sort is best sort [O(n)].

(i) 10 20 30 40 50 60 ⟹ 6 comparisons (actually 5, but we take

1 2 3 4 5 6

i j j j j j

(ii) 10 20 30 40 50 60 ⟹ 5 comparisons (actually 4, but we take 5).

i j j j j

6+5+4+3+2+1

T(6) = T(5) + 6 → 6 comparisons in 1st partition

= T(4) + 6 + 5

= T(3) + 6 + 5 + 4

= T(2) + 6 + 5 + 4 + 3

= T(1) + 6 + 5 + 4 + 3 + 2

= 1 + 2 + 3 + 4 + 5 + 6

→ depth of stack is 'n'

→ worst case

n / n-1 / n-2

2nd 60 50 40 30 20 10

i j j

10 50 40 30 20 60 → n swaps in worst case (1 swap in best case)

i j

10 50 20 30 40 60

i i

10 20 50 30 40 60

★ 1st time partition → 6 comparisons

2nd " " → 5 "

3rd " " → 4 "

1 "

∴ T(6) = T(5) + 6 = T(4) + 6 + 5 = $O(n^2)$

★ Quicksort gives worst case when array is sorted (whether in descending or ascending).

★ In descending order, the partitions will contain (n-1) & 0 elements which given $O(n^2)$ time complexity

Cosmos

★ Quicksort is best sort even though its worst case is $O(n^2)$ because the worst case i.e. (when the array is already sorted) never happens since we never sort an already sorted array.

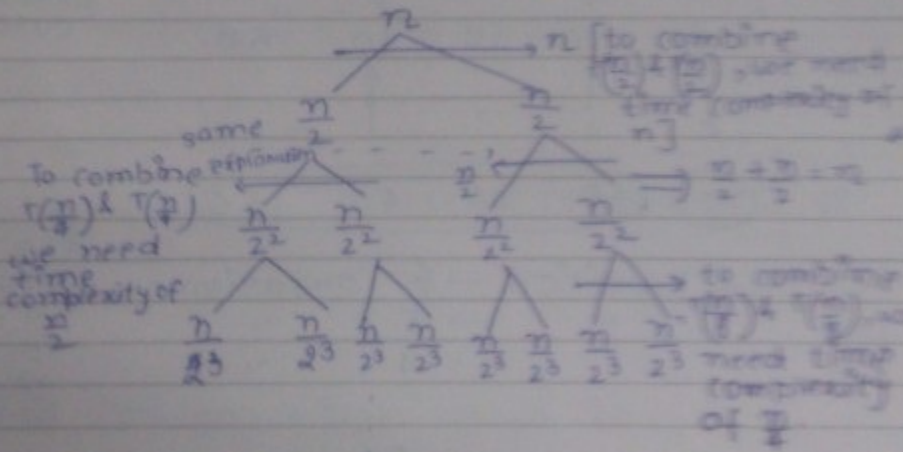
★ After applying few passes of quicksort ~~quick sort~~ algo on the given array we got following o/p

1 10 5 8 25 44 55 30 70

then how many pivot elements are there in above o/p

Ans. (5)

maximum 3 times the partition algo was applied it may be 1, 2 or 0 because [0 → when array is like that only.



$\frac{n}{2^k}$ $\frac{n}{2^k}$

Date

Cosmos

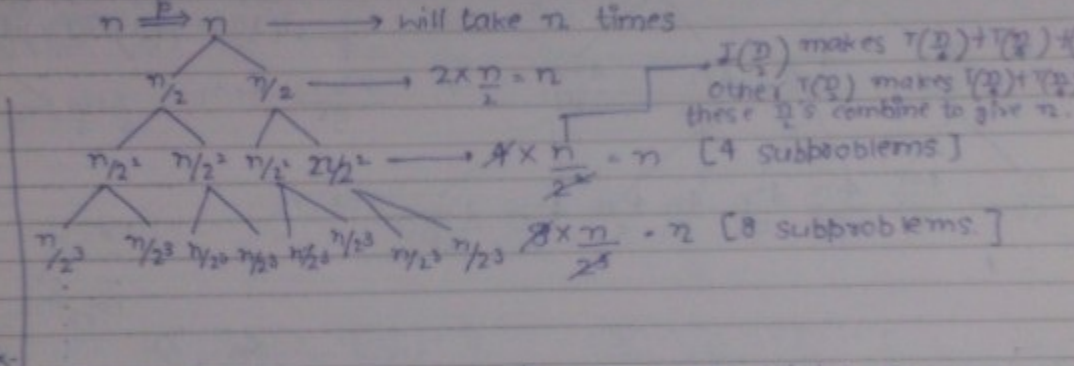
10.06.12

Recursive Tree Method :-

$T(n) = 2T(\frac{n}{2}) + n$ → combination

$T(\frac{n}{2}) + T(\frac{n}{2}) + n$

[method is useful when there are 3 elements]



$2^k \rightarrow$ n should be equal to 2^k

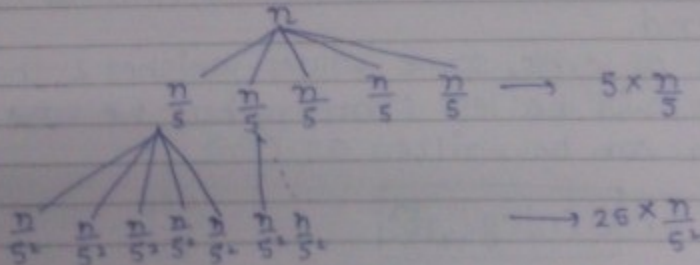
$n = 2^k \rightarrow k = \log_2 n$

[because in the end each subproblem should contain only one element]

$n \times k = n \times \log_2 n$

$O(n \log_2 n)$

Q. $T(n) = 5T(\frac{n}{5}) + n$



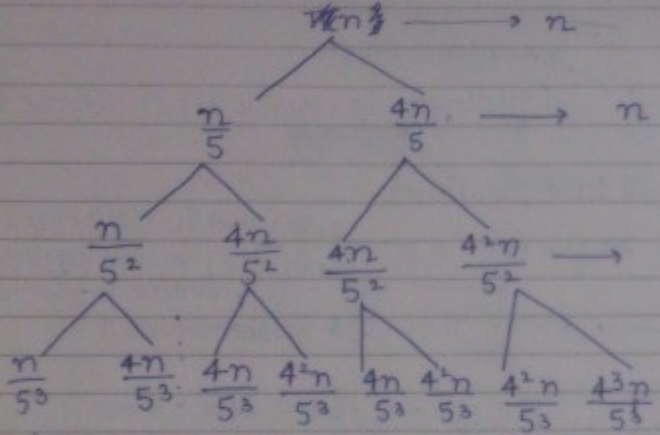
$5^k = n$

$k = \log_5 n$

$T(n) = O(n \log_5 n) + O(n)$

Cosmos

Q. $T(n) = T(\frac{n}{5}) + T(\frac{4n}{5}) + n$



$T(\frac{n}{5}) = T(\frac{n}{5^2}) + T(\frac{4n}{5^2}) + \frac{n}{5}$

$T(\frac{4n}{5}) = T(\frac{4n}{5^2}) + T(\frac{4^2n}{5^2}) + \frac{4n}{5}$

$\frac{25n}{5^2} \rightarrow n$

$\frac{n}{5} + \frac{4n}{5} = n$

$n = 5^3$
$n = 125$
$n = 5^k$
$\log_5 n = k$

$\frac{4^3n}{5^3} = \frac{n}{(5/4)^3}$

Dividing by 5/4 will create more levels, so $n/(5/4)^k$ creates most no. of levels, ∴ we for max. no. of levels, we equate n with $(\frac{5}{4})^k$

$4^k n = 5^k$

$n = (\frac{5}{4})^k$

$(\log_{5/4} n) = k$

My answer: $O(n \log_{5/4} n)$

★ $n/5^3$ will stop somewhere in the middle, whereas $n/(5/4)^k$ will go till the end.

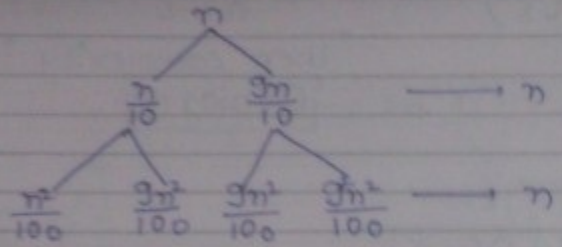
★ after some no. of levels, some values vanishes & the partition cost will be less than n, but we take $O(n)$, as less than n can be written as $O(n)$.

Sir's answer → $O(n \log_{5/4} n)$

Q. $T(n) = T(\frac{n}{10}) + T(\frac{9n}{10}) + n$

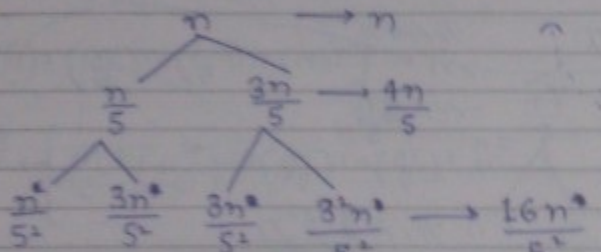
Cosmos

Similar to previous one



Complexity :- $\boxed{O(n \log_{(10/9)} n)}$

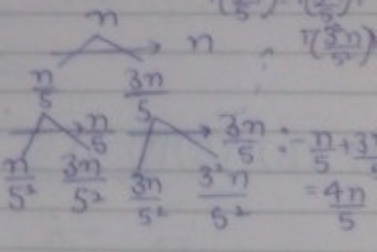
Q. $T(n) = T(\frac{n}{5}) + T(\frac{3n}{5}) + n$



$\frac{4^3 n}{5^3}$

$T(\frac{n}{5}) = T(\frac{n}{5^2}) + T(\frac{3n}{5^2}) + \frac{n}{5}$

$T(\frac{3n}{5}) = T(\frac{3n}{5^2}) + T(\frac{3^2 n}{5^2}) + \frac{3n}{5}$



$n + \frac{4n}{5} + \frac{4^2 n}{5^2} + \ldots + \frac{4^k}{5^k} n$ $\qquad \log_{5/3} n = k$

$\frac{n[1 - (\frac{4}{5})^{\log_{5/3} n}]}{\frac{4}{5} - 1} = 5n\left[1 - (\frac{4}{5})^{\log_{5/3} n}\right]$

↳ this is less than 1

∴ Complexity :- $5n\left[1 - (\frac{4}{5})^{\log_{5/3} n}\right]$

Cosmos

Note:-

If $T(n) = T\left(\frac{n}{a}\right) + T\left(\frac{n}{b}\right) + n$ if $\frac{n}{a} + \frac{n}{b} = n$

then complexity is $\boxed{O(n \log n)}$

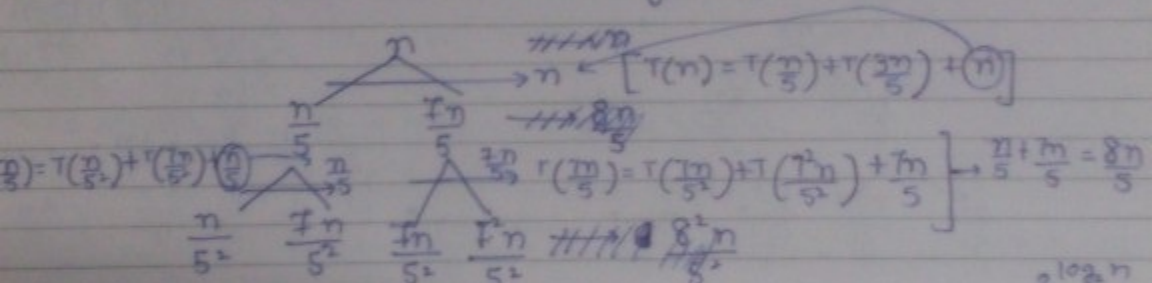
but if $\frac{n}{a} + \frac{n}{b} < n$

then complexity is $\boxed{O(n)}$

Q. $T(n) = 8T\left(\frac{n}{2}\right) + n^2 \longrightarrow O(n^2 \log n)$

Q $T(n) = T\left(\frac{n}{5}\right) + T\left(\frac{7n}{5}\right) + n$

this means problem is increasing stepwise



$n + \frac{8n}{5} + \frac{8^2 n}{5^2} + \ldots + \frac{8^k n}{5^k}$

$\frac{5^k}{8^k} = n$

$k = \log_{5/8} n$

$= \frac{n\left[\left(\frac{8}{5}\right)^{\log_{\frac{8}{5}} n} - 1\right]}{\frac{3}{5}} = \frac{5n}{3}\left[\left(\frac{8}{5}\right)^{\log_{5/8} n}\right]$

$= O\left(\left(\frac{8}{5}\right)^{\log_{8/5} n}\right)$

$\approx n \times \left(\frac{8}{5}\right)^{\log_{5/7} n} \approx n \times n^{\log_{5/7} 8/5}$

$\therefore$ complexity is $\boxed{O(n \times n^{\log_{5/7} 8/5})}$

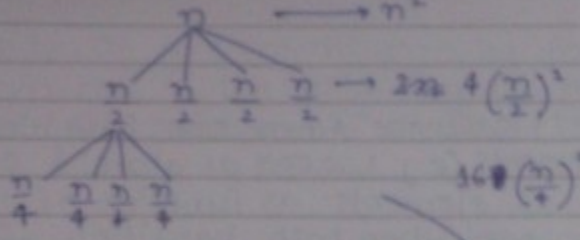
$2^{\log_2 n}$

$n^{\log_2 2}$

$= n$

Cosmos

Q. $T(n) = 4T(\frac{n}{2}) + n^2$

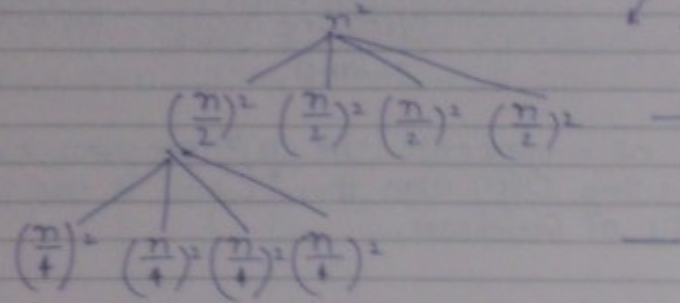


$\rightarrow 2nd\ 4\left(\frac{n}{2}\right)^2 = \frac{4n^2}{4} = n^2$ [each partition is taking $\left(\frac{n}{2}\right)^2$ time]

$\rightarrow 16\left(\frac{n}{4}\right)^2 = n^2$

no. of levels $\rightarrow n = 2^k \rightarrow k = \log_2 n$

$\therefore$ Complexity $\rightarrow O(n^2 \log_2 n)$



$\rightarrow n^2$

$\rightarrow 4\left(\frac{n}{2}\right)^2 = n^2$

$\rightarrow 16\left(\frac{n}{4}\right)^2 = n^2$

Quicksort :-

★ Quicksort will give the best case even when a single element is on one side of partition & even copies of element on other side.

$T(n) = T(\alpha n) + T(n - \alpha n) + n$ $\quad 0 < \alpha < 1$ [not equal, because if $\alpha = 0$ or 1 then it will give the worst case]

[Best Case]

★ $T(n) = T(10) + T(n-10) + n$

$\downarrow$

$O(n^2)$ → worst case [as n is very large & hence $(n-10)$ won't make any effect]

$\downarrow$

$T(n) = T(1,00,000) + T(n - 1,00,000) + n$

$O(n^2)$

☆ Worst-Case:-
Quicksort will give worstcase when one partition has constant no. of elements (e.g. 1,00,000 as above) & other partition has no. of elements as function of n.

☆ Best Case will happen when both partition will have no. of elements as function of n.

Worst-Case :- $T(n) = T(c) + T(n-(c+1)) + n$

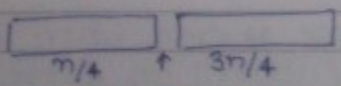
c :- constant
c+1 :- 1 because of pivot element [it won't be taken in the partition]

but as +1 in (c+1) won't affect the answer, ∴ we can write it as (n-c).

Q. In Quicksort, sorting of n-elements, the $(\frac{n}{4})^{th}$ smallest element is selected as pivot using $O(n)$ algo, then what is the worst case time complexity of Quicksort.

(a) $O(n^2)$
(b) $O(n \log n)$
(c) $O(n)$
(d) $O(n^3)$

Ans.

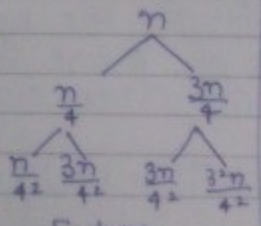


$O(n \log_{4/3} n)$

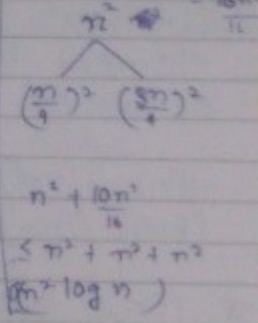
$T(n) = T(\frac{n}{4}) + T(\frac{3n}{4}) + O(n) + O(n)$

$O(n)$: place pivot at correct position

$O(n)$: selecting pivot acc. to the ques.



[when



$T(n) = T(\frac{n}{4}) + T(\frac{3n}{4}) + 2n$

Cosmos

If we have choosen n^{th} smallest element :-
then it is worst case.

Q. The median of n-elements can be found using O(n) algo. Which one of the following is correct about complexity of quicksort if median is selected as pivot
(A) $O(n^2)$
(B) $O(n \log n)$
(C) $O(n)$
(D) $O(n^3)$

Median :- The middle element of a sorted array.
this means $(\frac{n}{2})^{th}$ smallest element.

Stable Sorting Technique :-

The relative order of repeated elements not changed after sorting, then sorting technique is called Stable Sorting Techni

eg.

a[i] → 10_a, 11, 20, 10_b, 40, 50, 60
i → 1 2 3 4 5 6 7
i j j j

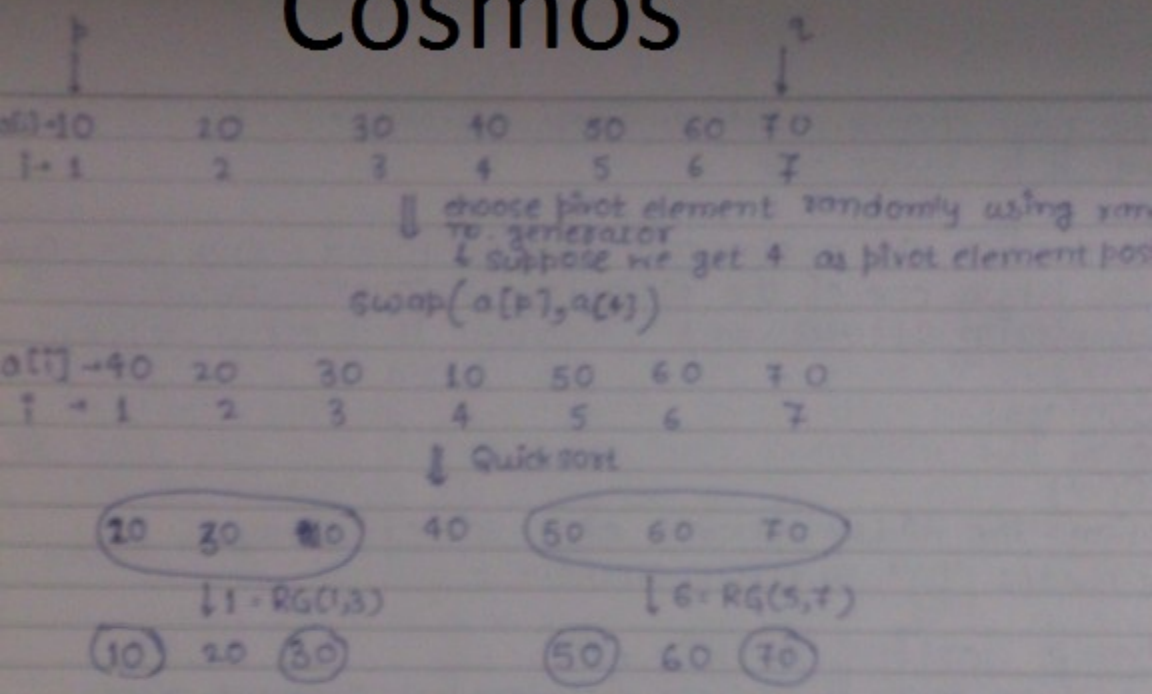
a[i] → 10_a 10_b 20 11 40 50 60 now, interchange a[i] & a[P]
i → 1 2 3 4 5 6 7
P i j

a[i] → 10_b 10_a 20 11 40 50 60
i → 1 2 3 4 5 6 7

★ Quicksort is not stable, but Mergesort is stable.

Randomized Quicksort :-

Cosmos



① Choosing pivot element randomly is called randomly quicksort

② No specific I/P will generate worst case performance for both sorted & unsorted assay will have $O(n \log n)$ complexity.

③ It is independent of I/P order.

④

Randomized Quicksort :-

```
RandomizedQS(a,p,r)
{ if(p<r)
    return(a[p]);
  else
  { i=RG(p,r);
    swap(a[p],a[i]);
    m=partition(a,p,r);
    RandomizedQS(a,p,m-1);
```

Quicksort	Randomized Quicksort
Best: $n \log n$	$n \log n$
Avg.: $n \log n$	$n \log n$
Worst: n^2	n^2
(almost sorted)	(when all the elements a
99% its not going to happen	same) 99.99999% its not going t happen

Selection Procedure:-

i/p:- An array of n-element & integer k

o/p:- Selecting k^{th} smallest element.

e.g.

40	70	20	60	10	50	30	k=4
1	2	3	4	5	6	7	

Algo-1

① Find 1st element & delete.
② " 2nd " " "
③ " 3rd " " "
④ " : k^{th} " & return it

⇒ $O(kn)$ when k=1 O(n) → best
when k=n $O(n^2)$ → worst case

Algo-2

① Sort the array ⟶ $n \log n$
② Return a[k] ⟶ $O(1)$

Best & Worst case → $O(n \log n)$

Algo-3

40	70	20	60	10	50	30
1	2	3	4	5	6	7

Selectionprocedure (a, i, j, k) → k^{th} smallest element

```
{ if (i==j)
    {if (i==k)
       return (a[i]);
    else
       return (-1);
    }
  else
    { if (m<k)
        partition
        Selectionprocedure(a,
      else
```

Cosmos

Let T(n) be the amount of time req., then

$T(n) = \begin{cases} \Omega(n) \rightarrow \text{best case} \\ O(n \log n) \rightarrow \text{worst case} \end{cases}$ → My answer

Sir's answer:-

$$T(n) = \begin{cases} O(1), & n = 1 \\ n + 2 + T\left(\frac{n}{2}\right) & \text{if } n > 1 \end{cases}$$

Best case (when the partition algo creates "almost" equal no. of elements on both sides.)

time for partition

$= T\left(\frac{n}{2}\right) + n$

$= \Omega(n)$

Worst case:-

$$T(n) = \begin{cases} O(1) & \text{if } n = 1 \\ n + 2 + T(n-1) & \text{if } n > 1 \end{cases}$$

→ Worst case (when the partition algo creates partition when one partition contains "constant" no. elements.)

time for partition

$= T(n-1) + n$

$= O(n^2)$

Strassen's Matrix Multiplication :-

① Without Divide & Conquer :-

$A_{n \times n}$, $B_{n \times n}$

1. $C_1 = A + B \longrightarrow O(n^2)$
2. $C_2 = AB \longrightarrow O(n^3)$

C_1 :- To select one element it will take $O(n^2)$
& to add two elements it will take $O(1)$
∴ complexity :- $O(n^2)$

C_2 :- To get the element in final matrix → $O(mnp)$
↓
$m \times n$ $n \times p$

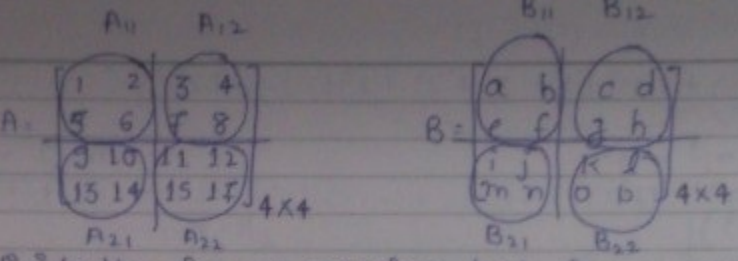
② Using Divide & Conquer :-

(i) If the given two matrices are $\leq 2 \times 2$, then the problem is small.

(ii) Only for square matrices.

(iii) The size of square matrices should be powers of 2.

endian



① divide the size of matrices by 2 till we reach the matrix size of 2×2 (e.g. the above is divided till we get 2×2)

e.g. $T(16) = T(8)$
$= T(4)$
$= T(2)$

$C = \begin{bmatrix} C_{11} & C_{12} \\ C_{21} & C_{22} \end{bmatrix}_{4\times4}$

$C_{11} = A_{11} * B_{11} + A_{12} * B_{21}$
$C_{12} = A_{11} * B_{12} + A_{12} * B_{22}$
$C_{21} = A_{21} * B_{11} + A_{22} * B_{21}$
$C_{22} = A_{21} * B_{12} + A_{22} * B_{22}$

8 multiplications of matrices 2×2

$T(4) = 8\,T\left(\frac{4}{2}\right) + 4 \times \left(4/2\right)^2$

4 additions of matrices of sizes 2×2

$T(8) = 8\,T\left(\frac{8}{2}\right) + 4 \times \left(\frac{8}{2}\right)^2$

$T(n) = \begin{cases} O(1), & \text{if } n \le 2\times2 \\ 8\,T\left(\frac{n}{2}\right) + 4 \times \left(\frac{n}{2}\right)^2 \end{cases}$

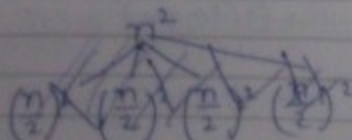
$T(16) = 8\,T\left(\frac{16}{2}\right) + 4 \times \left(\frac{16}{2}\right)^2$

8 matrix multiplications of size 8×8

4 matrix additions of size 8×8

each addition takes $O(n^2)$ complexity
∴ a 8×8 matrix will take $(8)^2$

$T(n) = 8\,T\left(\frac{n}{2}\right) + n^2$



$T(n) = 8\,T\left(\frac{n}{2}\right) + n^2 = 64\,T\left(\frac{n}{4}\right) + 8\frac{n^2}{4} + n^2$

Cosmos

$= 2^k T\left(\frac{n}{2^k}\right) + n^2\left[1 + 2^1 + 2^2 + 2^3 + \ldots + 2^{k-1}\right]$

$\frac{n^2[2^k - 1]}{1}$

$n_2 + n^3 - n^2$

$\leq n^3 + n^3 - n^3$

$O(n^3)$

$2^k = n$

Note:-

① Using divide & conquer & without using it, it will take same time $O(n^3)$ or $\theta(n^3)$.

② Without using divide & conquer is better because of stack space

★ Strassen Method:-

$$T(n) = \begin{cases} O(1), & \text{if } n \leq 2 \times 2 \\ 7T\left(\frac{n}{2}\right) + 16 \times \left(\frac{n}{2}\right)^2 & \end{cases}$$

$aT\left(\frac{n}{b}\right) + f(n)$

$n^{\log_b a} = n^{\log_2 7}$

$f(n) = n^2$ (as constants dont affect complexity)

$n^{\log_2 7} > n^2$ ∴ complexity:- $O(n^{\log_2 7})$

7 multiplications of size $\frac{n}{2} \times \frac{n}{2}$

16 additions of size $\frac{n}{2} \times \frac{n}{2}$ [each will take $a\left(\left(\frac{n}{2}\right)^2\right)$ time $[O(n^2)]$ for $n \times n$]

$O(n^{2.81})$ → Strassen method

But the best case is → $O(n^{2.32})$

Conclusion

① Max.-min. → $2T\left(\frac{n}{2}\right) + 2 \longrightarrow O(n)$

② Power of an element → $T\left(\frac{n}{2}\right) + 1 \longrightarrow O(\log n)$

③ Binary Search → $T(n/2) + C \longrightarrow O(\log n)$, $\Omega(1)$

④ Merge Sort → $2T\left(\frac{n}{2}\right) + n \longrightarrow \theta(n \log n)$

⑤ Quick Sort → $2T\left(\frac{n}{2}\right) + n \longrightarrow \Omega(n \log n)$

(e) Selection procedure $\begin{cases} T\left(\frac{n}{2}\right)+n \longrightarrow \Omega(n) & \text{Avg} \\ T(n-1)+n \longrightarrow O(n^2) \end{cases}$

(f) Matrix Multiplication $\longrightarrow T(n) = 8T\left(\frac{n}{2}\right)+n^2 \longrightarrow O(n^3)$

$\hookrightarrow T(n) = 7T\left(\frac{n}{2}\right)+n^2 \longrightarrow O(n^{2.81})$

Q.

```
A(n)
{ if(n≤1) return;
  else
    return(A(√n));
}
```

(A) O(n)
(B) O(log n)
(C) O(log log n)
(D) $O(n^2)$

rec. reln: $T(n) = \begin{cases} O(1), & n \le 1 \\ T(\sqrt{n}) + O(1), & \text{otherwise} \end{cases}$ ← to calculate $\sqrt{n}$

$T(n) = T(\sqrt{n}) + O(1)$

$= T(n^{1/2}) + O(1)$

$= T(n^{1/4}) + O(1)$

$= T(n^{1/8}) + O(1)$

$= T(n^{1/2^k}) + O(1)$

$= T(1) + kc$

$= O(1) + c \log\log n$

$= O(\log\log n)$

$n^{1/2^k} = 2$

$\frac{1}{2^k}\log_2 n = \log_2 2$

$\frac{1}{2^k}\log_2 n = C_1$

$\frac{1}{C_1}\log_2 n = 2^k$

$\boxed{\log_2\log_2 n = k}$

Q.

```
DAA(n)
{ if(n≤1) return;
  else
  {
  return(DAA(n/2)+DAA(n/2)+100);
  }
}
```

$T(n) = \begin{cases} O(1), & n \le 1 \\ 2T\left(\frac{n}{2}\right)+c, & \text{otherwise} \end{cases}$

$= 2\left[2T\left(\frac{n}{4}\right)+c\right]+c$

$= 2^2 T\left(\frac{n}{4}\right) + 2c + c$

$= 2^2\left[2T\left(\frac{n}{8}\right)+c\right]+2c+c$

$T(n) = nT(1) + C_1 \quad [n = 2^k]$

$= n + C_1$

$= \boxed{O(n)}$

```
A (n)
{ if (n ≤ 1)
    return;
  else
    return (A(n/2) + A(n/2) + n);
}
```

Similar to:
```
b = A(n/2);
c = A(n/2);
d = b + c + n;
return(d);
```

This is just addition & it will take O(1) time for addition. [not a function call.]

$$T(n) = \begin{cases} O(1), & n \le 1 \\ T(\frac{n}{2}) + T(\frac{n}{2}) + C, & \text{otherwise} \end{cases}$$

$= 2T(\frac{n}{2}) + C$

$= 2\left[2T(\frac{n}{4}) + \frac{n}{2}\right] + C$

$= 2^2 T(\frac{n}{2^2}) + n + n$

$= 2^2\left[2T(\frac{n}{2^3}) + \frac{n}{4}\right] + n + n$

$= 2^3 T(\frac{n}{2^3}) + n + n + n$

$= 2^k T(\frac{n}{2^k}) + kn$

put $2^k = n$

$\therefore nT(1) + n\log_2 n\, C$

$= \boxed{O(n)}$

My mistake. [taken n as function call.]

```
return (A(n/2) * A(n/2));
```

$T(n) = 2T(\frac{n}{2}) + C$ → for multiplication.

* Master theorem is applicable when $f(n)$ is a [illegible] polynomial function.

Master Theorem :-

It is applicable only in the form of divide & conquer, e.g.

$T(n) = aT\left(\frac{n}{b}\right) + f(n)$ $\quad a \geq 1$, $b > 1$

Case 1: $f(n) = O(n^{\log_b a - \epsilon})$ where $\epsilon > 0$

then

$$\boxed{T(n) = \theta(n^{\log_b a})}$$

Case 2: $f(n) = \Omega(n^{\log_b a + \epsilon})$ where $\epsilon >$

then $\boxed{T(n) = \theta(f(n))}$

Case 3: $f(n) = \theta(n^{\log_b a})$

then $\boxed{T(n) = \theta(n^{\log_b a} * \log n)}$

e.g. ① $T(n) = 8T\left(\frac{n}{2}\right) + n^2 \rightarrow a=8, b=2 \quad n^{\log_b a} = n^3$

$n^2 = O(n^3)$

$f(n) < n^{\log_b a}$ $\quad \boxed{\theta(n^3)} \rightarrow$ Case 1

② $T(n) = 2T\left(\frac{n}{2}\right) + n \longrightarrow a=2, b=2 \quad n^{\log_b a} = n$

$f(n) = n^{\log_b a}$

$\therefore \boxed{\theta(n \log n)} \longrightarrow$ Case 3

③ $T(n) = 2T\left(\frac{n}{2}\right) + 1$

$n^{\log_b a} = n \quad , f(n) = 1 \quad f(n) < n^{\log_b a} \longrightarrow \boxed{\theta(n)} \rightarrow$ Cas

④ $T(n) = T\left(\frac{n}{2}\right) + c$

$n^{\log_2 1} = n^0 = 1 \quad f(n) = c \quad n^{\log_b a} = f(n) \quad \therefore \theta(n^0 \log n) = \theta(\log n$

⑤ $T(n) = 4T\left(\frac{n}{2}\right) + n^3$ $\quad$ Case 3 :- $\boxed{\theta(n^3)}$

$n^{\log_2 4} = n^2$

Imp.

⑥ $T(n) = 8T\left(\frac{n}{2}\right) + n^2 \quad \theta(n^3)$

$n^{\log_2 8} = n^3$

$n^2 = O(n^3)$

★ ϵ tells that n^3 is how many times greater

* Master theorem is applicable when $f(n)$ is a func polynomi function.

Master Theorem :-

It is applicable only in the form of divide & conquer, e.g.

$T(n) = aT\left(\frac{n}{b}\right) + f(n)$ $\quad a \geq 1,\ b > 1$

Case 1:- $f(n) = O(n^{\log_b a - \epsilon})$ where $\epsilon > 0$

then $\boxed{T(n) = \theta(n^{\log_b a})}$

Case 2:- $f(n) = \Omega(n^{\log_b a + \epsilon})$ where $\epsilon > 0$

then $\boxed{T(n) = \theta(f(n))}$

Case 3:- $f(n) = \theta(n^{\log_b a})$

then $\boxed{T(n) = \theta(n^{\log_b a} * \log n)}$

e.g. ① $T(n) = 8T\left(\frac{n}{2}\right) + n^2 \rightarrow a=8, b=2 \quad n^{\log_b a} = n^3$

$n^2 = O(n^3)$

$f(n) < n^{\log_b a}$ $\quad \boxed{\theta(n^3)} \rightarrow$ Case 1

② $T(n) = 2T\left(\frac{n}{2}\right) + n \longrightarrow a=2, b=2 \quad n^{\log_b a} = n$

$f(n) = n^{\log_b a}$

$\boxed{\theta(n \log n)} \longrightarrow$ Case 3

③ $T(n) = 2T\left(\frac{n}{2}\right) + 1$

$n^{\log_b a} = n \quad , f(n) = 1 \quad f(n) < n^{\log_b a} \longrightarrow \boxed{\theta(n)} \rightarrow$ Case

④ $T(n) = T\left(\frac{n}{2}\right) + C$

$n^{\log_2 1} = n^0 = 1 \quad f(n) = C \quad n^{\log_b a} = f(n) \quad \therefore \theta(n^0 \log n) = \theta(\log n)$

⑤ $T(n) = 4T\left(\frac{n}{2}\right) + n^3$ Case 3 :- $\boxed{\theta(n^3)}$

$n^{\log_2 4} = n^2$

⑥ $T(n) = 8T\left(\frac{n}{2}\right) + n^2$ $\quad \theta(n^3)$

$n^{\log_2 8} = n^3 \quad n^2 = O(n^3)$

Imp.

★ ϵ tells that n^3 is how many times greater than

Cosmos

Cosmos

$n^{\log_b a}$ must be atleast polynomial times greater than $f(n)$.

Case 1:- $n^{\log_b a}$ is polynomial times greater than $f(n)$.
Case 2:- $n^{\log_b a}$ is polynomial times smaller than $f(n)$.
Case 3:- $n^{\log_b a}$ is asymptotically equal.

$T(n) = 2T(\frac{n}{2}) + n\log n$

$n^{\log_2 2} = n$

$f(n) > n^{\log_b a}$ ⟶ Case 2

$\frac{f(n)}{n^{\log_b a}} = \log n$:- but we can't write $\log n$ in terms of n^ϵ

$n^{\log_b a}$ is logarithmic time smaller than $f(n)$, so master theorem is not applicable. So, $n\log n > n^1$

$T(n) = T(\sqrt{n}) + c$ ⟶ ?

$n^{\log_b a} = 1$ $[a=1, b=1]$ but b should be greater than 1.
$\Theta(\log n)$

The above method is not in divide & conquer format.
Now to use master method, we will convert it into divide & conquer format

Step 1:- Assume $n = 2^k$

$T(2^k) = T(2^{k/2}) + c$

Step 2:- assume $T(2^k) = S(k)$

$S(k) = S(\frac{k}{2}) + c$

$n^{\log_2 1} = n^0 = k$ $\quad k^{\log_2 1} = k^0 = 1$

$n\neq 1$, $\Theta(n^0 \log k)$ Case 3:- $\Theta(k^0 \log k)$

$= \boxed{\Theta(\log \log n)}$

$T(n) = 2T(\sqrt{n}) + \log n$

put $n = 2^k$

$T(2^k) = 2T(2^{k/2}) + \log_2 2^k$ ⟶ $k^{\log_2 2} = k^1 = k$

$T(2^k) = 2T(2^{k/2}) + k$

$k = k$

$\Theta(k)$ ⟶ $\boxed{\Theta(\log_2 n)}$

put $S(k) = 2S(\frac{k}{2}) + k$

Insertion Sort:-

Note:- Insert & Sort.

example:- A[10 20 30 25 40 50 15 45]
1 2 3 4 5 6 7 8

★ It is inplace sorting technique.

1st pass:-

1. [10]
2. [10|20]
3. [10|20|30]
4. [10|20|30|25]
5. [10|30|25|30|40]
6. [10|20|25|30|40|50]
7. [10|20|25|30|40|50|15]
8. [10|15|20|25|30|40|50|45] → [10|15|20|25|30|40|45|50]

Best Case:- 10,20,30,40,50

1. [10]
2. [10|20]
3. [10|20|30]
4. [10|20|30|40]
5. [10|20|30|40|50]

→ n-1 comparisons

∴ Ω(n) → Best case (when it is already sorted)

Note ①:- To sort n elements using insertion sort, it will take (n-1) comparisons in the best case.

② If array is almost sorted, then insertion sort is the best

③ If array size is small, then insertion sort is best sort (merge sort is worst in this case)

Worst Case:-

70, 60, 50, 40, 30, 20, 10

1. [70]
2. 60 [70|60]
3. 50 [60|70|50]
4. 40 [50|60|70|40]
5. 30 [40|50|60|70|30]
6. 20 [30|40|50|60|70|20]

[20|30|40|50|60|70]

Comparisons	Swap
0	0
1	1
2	2
3	3

$2n^2$ → O(n²) → worst case

Average case:-

Cosmos

While applying insertion sort, to insert a particular element to its correct place, for that we are using linear search, then what is the worst case time complexity of insertion sort if we replace by binary search?

) $O(n^2)$ (C) $O(\log n)$
) $O(n \log n)$ (D) $O(n)$

My answer :- (B)

Sir's Answer :-

Worst case

LS comparisons	LS swap	BS comp.	BS swap
•	•		
0	0	0	0
1	1	log1	1
2	2	log2	2
3	3	log3	3
n-1	n-1	log(n-1)	n-1
n^2	n^2	nlogn	n^2

→ swap operations do not decrease in binary search

$2n^2$

$O(n^2)$

$n \log n + n^2 \longrightarrow \boxed{O(n^2)}$ → ans.

How to decrease no. of swaps.

If I/p is present in Linked List.

In linked list we can only perform Linear Search.

Which will take $O(n^2)$ n^2 comparisons & 0 swaps.

∴ complexity :- $\boxed{O(n^2)}$

$0+1+2+3+\dots+n \le n \cdot n = n^2$

Selection Sort

	25	5	58	45	32	64	3	11
	1	2	3	4	5	6	7	8

S1 :- i = 1

min = 1 < 7 [compare (i) select position 1 to be the 1st minimum.

(minimum's (ii) compare the minimum with

Straight Selection Sort :-

(iii) swap(a[i], a[min])

eg. a[i] → 3 5 58 45 32 64 25 11
1 2 3 4 5 6 7 8

pass 2 :- now, i = 2
min = 2

pass 3 :- i = 3
min = ~~3 4 5 7 8~~

a[i] → 3 5 11 45 32 64 25 58
1 2 3 4 5 6 7 8

pass 4 :- i = 4
min = ~~4 8 7~~

a[i] → 3 5 11 25 32 64 45 58

pass 5 :- i = 5
min = 5

pass 6 :- i = 6
min = ~~6 7~~

a[i] → 3 5 11 25 45 64 58

pass 7 :- i = 7
min = ~~7 8~~

a[i] → 3 5 11 25 45 58 64

Complexity :-

1st pass :- n-1 comparisons
2nd pass :- n-2 comparisons
3rd " :- n-3 "
4th " :- n-4 "

∴ (n-1) + (n-2) + (n-3) + ... + 1 ≤ n + n + n + ... n times = n^2

Best Case :- $\Omega(n^2)$ $\boxed{\Theta(n^2)}$
Worst Case :- $O(n^2)$

Bubble Sort :- $\boxed{\Theta(n^2)}$

i/p :- 75 58 10 84 5
1 2 3 4 5

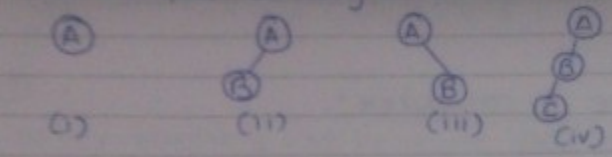
pass 1 : (75) (58) 10 84 5

pass 2 :- (58) (10) 75 5 84
10 (58) (75) 5 84
10 58 (75) (5) 84

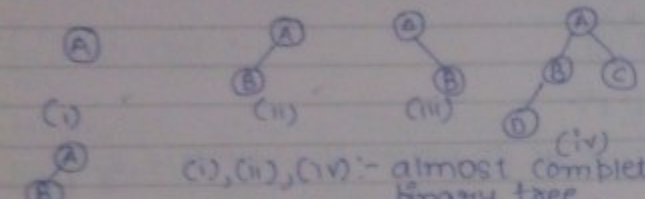
Cosmos

Heapsort

Almost Complete Binary Tree.



(i),(ii),(iii),(iv)
↓
Binary Tree



(i),(ii),(iv) :- almost complete binary tree.

(iii),(v) :- not an almost complete binary tree.

In almost Complete Binary Tree :-

(i) Without completing left, we can't go to the right.

(ii) Without completing a level, we can't go to the next level.

A binary tree is said to almost complete binary tree iff :-

① At every node, after completing left child, only then go to right child

② At every node, after completing the present level, only then go to next level.

In the given almost complete binary tree, if last level is also filled completely, then it is called complete binary tree.

Another name for complete binary tree :- maximal binary tree.

$1+2+4+\ldots+2^n$

$\frac{1(2^{n+1}-1)}{2-1}$

$= 2^{n+1}-1$

No. of levels = $n+1$

Note :- A complete binary tree with K-levels, the no of nodes = 2^k-1

A binary tree with k-levels, the max. no. of nodes = 2^k-1.

Note :- The no. of nodes in a complete binary tree at level i = 2^{i-1}

2. If there are n nodes in a binary tree :-

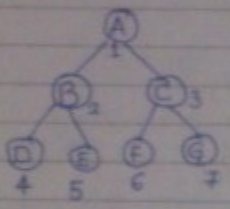
max. no. of levels = n

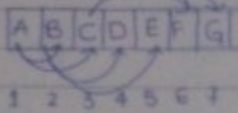
Cosmos

★ If there are binary tree & complete binary tree, then complete binary tree will take less time for any arbitary operation.

How to represent binary tree in computer:-

Properties:-





Parent of node at i^{th} place:- if i is even = $\frac{i}{2}^{th}$ po

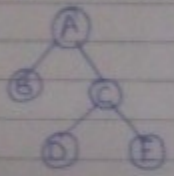
if i is odd = $\lfloor \frac{i}{2} \rfloor^{th}$ p

(floor v

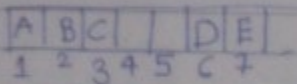
Child:- Left Child of node at i^{th} position = $2i^{th}$ position

Right Child of node at i^{th} position = $(2i+1)^{th}$ position

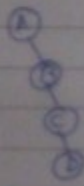
Q Store the following binary tree using array representation -

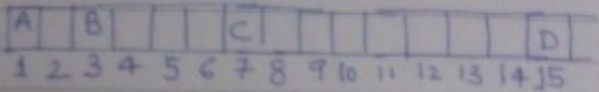


⇒



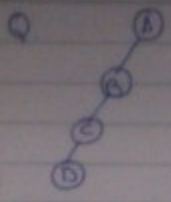
Q Store the following Binary Tree in form of array:-

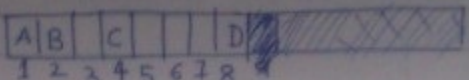




★ Array Representation of Binary Tree is favourable when tree contai is almost complete binary tree. Linked List Representation is favoured when tree is not almost complete (like the above question).

Cosmos





★ After reaching last node, we ~~stopp~~ stop.

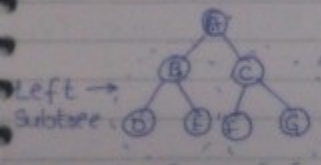
Heapsort :-

- Heapsort require Heap trees.
- Heap trees { Min heap Tree / Max Heap Tree

★ Heap Trees are almost complete binary trees.

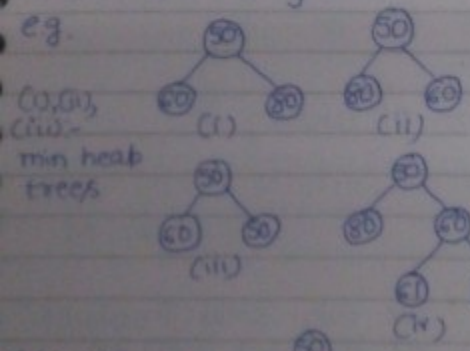
Min Heap Tree

- An almost comp
- In a given almost complete binary tree, at every node root is minimum, (or equal) in comparison to its children, then the tree is called min heap tree.



★ We dont have to compare Sub trees, but only children.

eg.

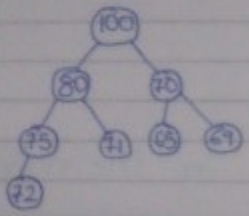


(i), (ii), (iii) - min heap trees

(iv) - not a min heap tree because it is not

Max Heap Tree

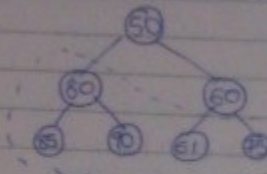
- In a given almost complete binary tree, at every node is maximum, (or equal) in comparison its children, then it is max heap tree.



★ If there are n no. of leaf nodes, then no. of internal nodes are n-1
So total no. of nodes = 2n-1

[Only for (complete) binary trees.]

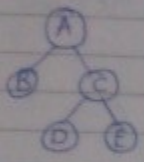
★ The root of min heap tree is the minimum element.
★ The 2nd min. is from 2nd level.
★ The 3rd min. can be :-



The 3rd min. can be in dotted places.

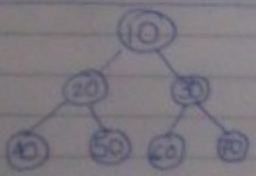
Strict Binary Tree :-

• Every node must have either 0 children or 2 children



★ The max. element in the min heap tree is at the last level

Insertion in Min heap tree :-



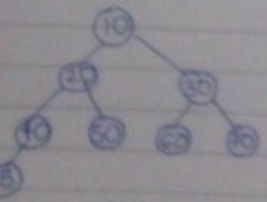
⇒

(i) Store tree in form of array :-

10	20	30	40	50	60	70
1	2	3	4	5	6	7

(ii) Now, insert '5'.

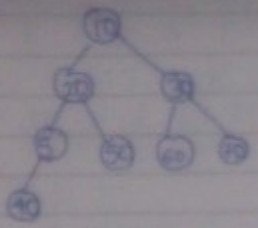
10	20	30	40	50	60	70	5
1	2	3	4	5	6	7	8



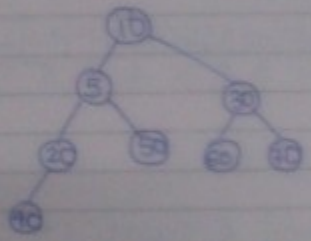
5 is the child of (4÷2)

We can only insert at inserted position because we have to make it almost complete binary tree.

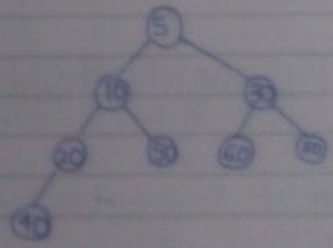
Now 5 will be compared with 40, where 5 is min, so it will be swapped



(iv) Now 5 will be compared with its parent (20), & hence swapped



(v) Now 5 will be compared with its parent (10), & hence swapped.

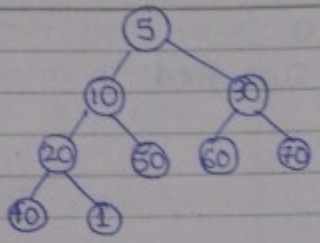


Now, in array:-

10					

Now, Insert '1'

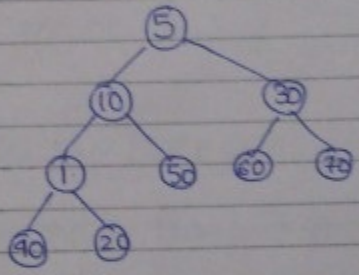
Cosmos



5	10	30	20	50	60	70	40	1
1	2	3	4	5	6	7	8	9

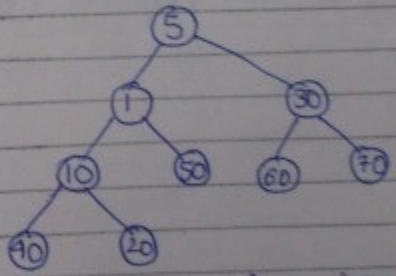
1 is right child of 20 which is at 4, ∴ 1 is at 4×2+1=9.

Now, compare '1' with its parent & swap it:-



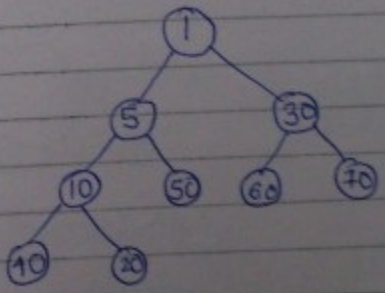
5	10	30	1	50	60	70	40	20
1	2	3	4	5	6	7	8	9

Now, compare '1' with its parent & swap it:-



5	1	30	10	50	60	70	40	20
1	2	3	4	5	6	7	8	9

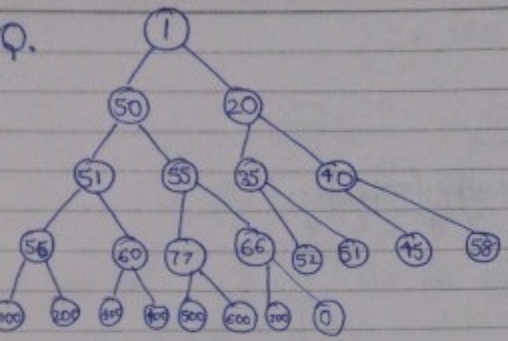
Now, compare '1' with its parent & swap it:-

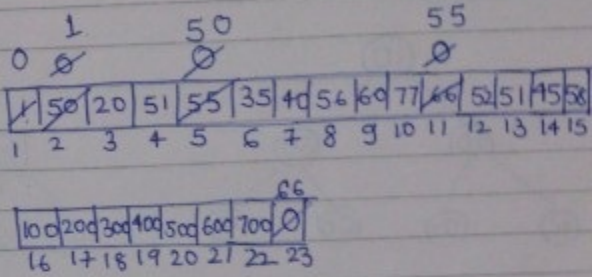


1	5	30	10	50	60	70	40	20
1	2	3	4	5	6	7	8	9

Cosmos

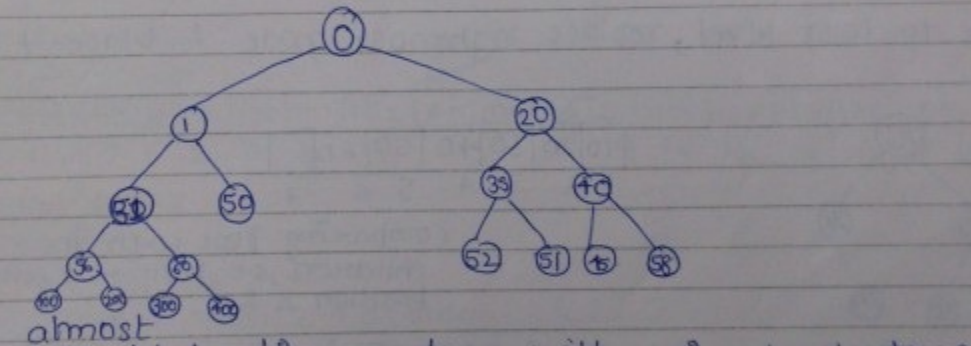
Q.





Now, insert '0'.

- Compare 0 with element at 11th position



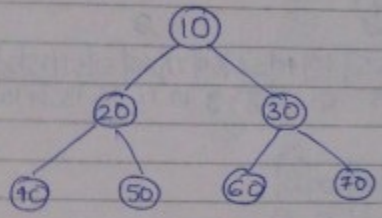
* In complete binary tree with min heap tree properties, there are $\lceil \log_2 n \rceil$ levels, so ~~this~~ at max. there will be $\lceil \log_2 n \rceil$ comparisons & $\log_2 n$ swaps. $\therefore \log_2 n + \log_2 n = 2\log_2 n$
- Worst Case :- $O(\log_2 n)$
- Best Case :- $\Omega(1)$
- Average Case :- $\Theta(\log_2 n)$

* In order to insert an element into min heap o max. heap which already contains n elements requ $O(\log_2 n)$ (Worst Case & Average Case).

Cosmos

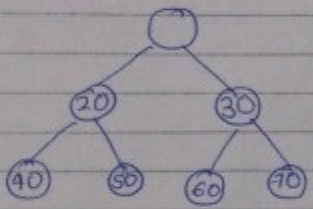
★ We always delete the root.

Deletion



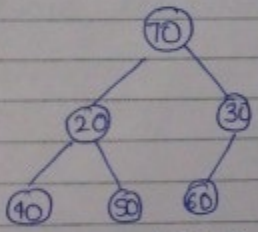
10	20	30	40	50	60	70
1	2	3	4	5	6	7

- Now, delete '10'.



	20	30	40	50	60	70
1	2	3	4	5	6	7

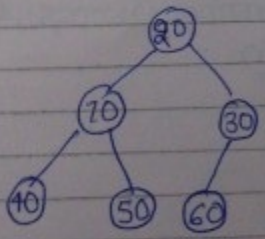
Now, go to last level, replace rightmost node & place it at root.



70	20	30	40	50	60	
1	2	3	4	5	6	7

Comparing root with its children, i.e. position 1 with position 2 & 3.

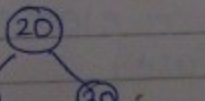
Now, compare 70 with 20 & 30, 20 is min., ∴ Swap replace 70 & 20



20	70	30	40	50	60	
1	2	3	4	5	6	7

Comparing position 2 with its childre (4 & 5)

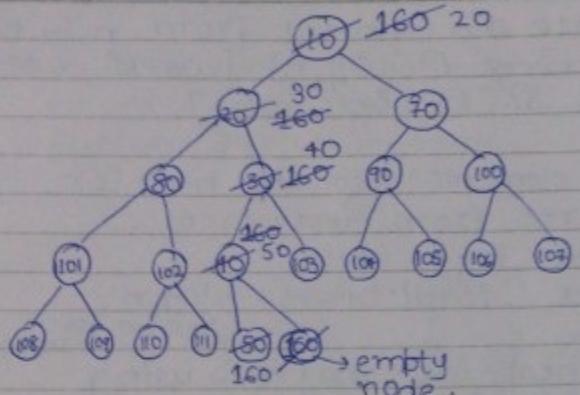
Now compare 70 with 40 & 50, 40 is min, ∴ Swap 70 & 40.



20	40	30	70	50	60	

Cosmos

Q.



Answer

1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20
20	30	70	80	40	90	100	101	102	50	103	104	105	106	107	108	109	110	111	160

1	2	3	4	5	6	7	8	9	10	11	12	13	14	15	16	17	18	19	20	21
10	20	30	80	40	90	100	101	102	50	103	104	105	106	107	108	109	110	111	160	

1. ~~Swap 1 & 2~~
1. Delete 1.
2. Place 21 at 1
3. Compare 1 with 2 & 3, swap 1 & 2.
4. Compare 2 with 4 & 5, swap 2 & 5
5. Compare 5 with 10 & 11, swap 5 & 10.
6. Compare 10 with 20, ~~& 21~~ swap 10 & 20.

In deletion:- (or level)
At every step, there are 2 comparisons & 1 swap,
so no of levels = $\log_2 n$.
∴ Time complexity = $3\log_2 n = \boxed{O(\log_2 n)}$ → Worst Case

Average Case:- $\boxed{\theta(\log_2 n)}$

Best Case:- $\boxed{\Omega(1)}$

Cosmos

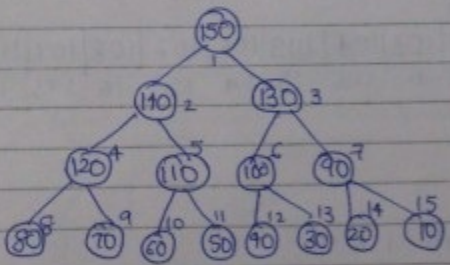
Note:- In order to delete an element from min. heap or max. heap requires $O(\log_2 n)$ [worst case & average case.] & $\Omega(1)$ [Best Case.]

* If we have to insert n elements into the heap (i.e. constructing min heap tree from empty tree):- each insertion will take $\log_2 n$ time, as there are n nodes, ∴ total time = $n \log_2 n$. (Brute Force Method)

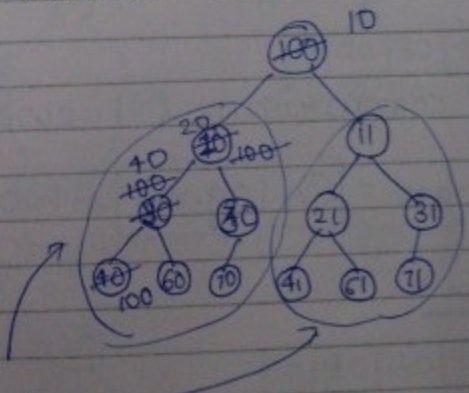
* But Constructing min heap (or) max heap using Build heap will take $O(n)$ time.

i/p:- 150, 140, 130, 120, 110, 100, 90, 80, 70, 60, 50, 40, 30, 20, 10
create min-heap tree.

Step 1:- Construct almost complete binary tree.



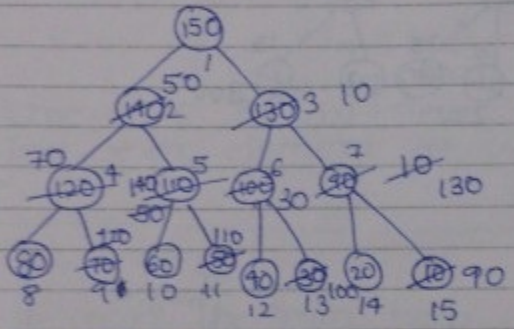
* We can use min heapify method only when left & right subtrees of root are themselves minheap tree.



Cosmos

No

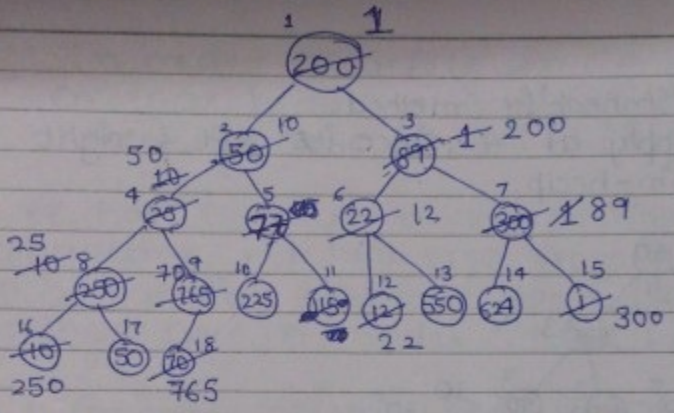
Step 2:- Now, apply Minheapify method we can only apply at 7, since its left & right subtree are min-heap.

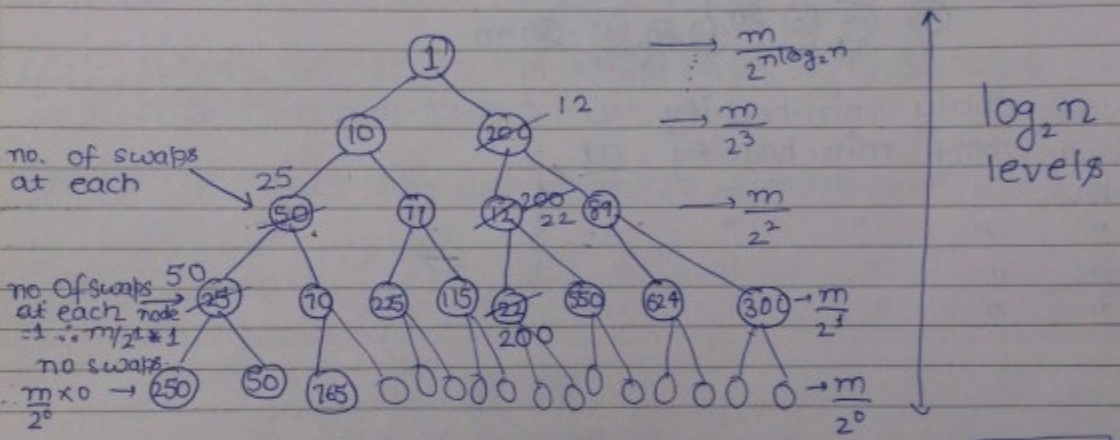


- Now, apply min-heapify at 6.
- Now, apply min-heapify at 5.
- Now, " " " " 4.
- " , " " " " 3.
- " , " " " " 2.
- " , " " " " 1

). Construct min heap tree for following methods:-
200, 50, 89, 25, 77, 42, 300, 250, 765, 225, 115, 12, 550, 624, 1, 10
70.

P.T.O.





m :- no. of nodes at last level
n :- total no. of nodes in tree.
] → $m = \frac{n}{2}$

Time complexity of build heap method =

$$\frac{m}{2^0} \times 0 + \frac{m}{2^1} \times 1 + \frac{m}{2^2} \times 2 + \dots + \frac{m}{2^{(\log n)-1}} \times [(\log n) - 1]$$

$$= m \underbrace{\left[\frac{1}{2^1} + \frac{2}{2^2} + \frac{3}{2^3} + \dots + \frac{[(\log_2 n) - 1]}{2^{(\log n)-1}} \right]}_{O(1)}$$

$$= m * O(1)$$

$$= \frac{n}{2} * O(1)$$

$$= \boxed{O(n)}$$

Cosmos

* Deleting
* Finding 1st min. element → $\log_2 n$
* Finding 2nd min. element → $2\log_2 n$
* Finding 3rd min. element → $3\log_2 n$ — V. Imp.
* Finding n^{th} min. element → $O(n)$

because n^{th} min. element = max. element.

& max. element is in last level,
& so to find max. element at last level, it will take $O(n)$ time.

Heapsort

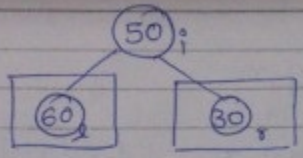
) $O(n)$ time to make the min. heap.
) Deleting the i^{th} min. element from min. heap :- $O(i \log_2 n)$, this i^{th} element is

eapsort:-
) Create min. heap Using Build heap :- $O(n)$
) Delete the elements one by one & store from right hand side of original array.
each deletion takes $O(\log_2 n)$ & storing at right hand side takes $O(1)$ time,
so to delete 'n' elements & storing them at right hand side :- $O(n \log_2 n)$

Complexity :- $O(n) + O(n \log_2 n)$
$= \boxed{O(n \log_2 n)}$

Min heap will give descending order. → In-place sorting.
Max heap " " ascending " .

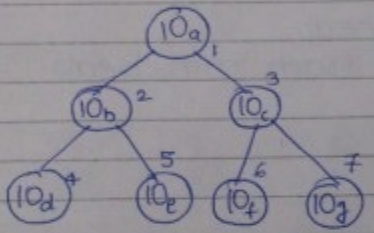
Cosmos




```
Min heapify(i)
{ min = i;
  l = 2*i;
  r = (2*i)+1;
  if (a[min] > a(l))
      min = l;
  if (a[min] > a(r))
      min = r;
  swap(a[i], a[min]);
}
```

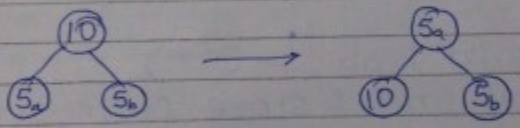
10_a 10_b 10_c 10_d 10_e 10_f 10_g

onstruct min-heap:-

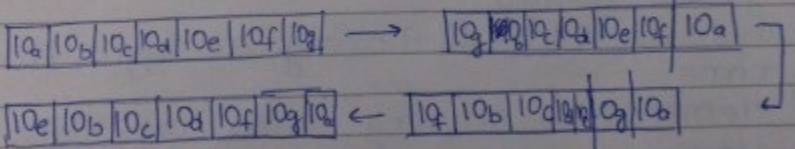


Using min-heapify:-
there will be no swaps.

&



[from min-heapify]



Heapsort is not stable sorting technique as the ordering of similar elements changes.
Heapsort & Quicksort are not stable, rest others are stable.
All sorting techniques are in-place except Merge-Sort.

Cosmos

Conclusion :-

① Min heapify :- $O(\log_2 n)$, $\Omega(1)$
② Build-heap :- $\Theta(n)$
③ Finding min. :- $O(1)$
④ Deleting min :- $O(\log_2 n)$
⑤ Deleting 5th min. :- $5\log_2 n$:- $O(\log_2 n)$
⑥ Finding 5th min. :- Deleting upto 4th min. :- $4\log_2 n$:- $O(\log$
⑦ Deleting nth min. :- Deleting max. element :- $O(n)$
max. element is in last level, & hence taking only last level & finding max. element from it & deleting it will take $\frac{n}{2}$ which is $O(n)$.

Greedy Technique (V.V.Imp.)

i/p :- Most of the problems in Greedy contain n-i/p's & our goal is finding subset which will optimize our objective.

Basics :-

① Solution Space :- all possible solutions. (For given n-i/p all possible soln (may or may not be correct) is called soln space.
② Feasible Solution.
③ Optimal Solution

- Feasible :- It is one of the soln from soln. space which satisfy our condition.
- Optimal :- It is one of the feasible soln. which optimizes our goal.

★ e.g. if we have a class of 200 students & we want to find out top 10 students.

- Solution Space :- $^{200}C_{10}$:- all combinations of 10 students.
- Feasible Solution :- all that combinations of 10 studen which satisfy the objective, i.e. whose avg. is greater than a specified value.

★ To find top 10 students from 200 students, then create min heap which will take $O(n)$ time, & delete first 10 max element which will take $10\log_2 n$ time. ∴ $O(n) + O(10\log_2 n) = \boxed{O(n)}$.

Control abstraction of Greedy

```
Greedy Technique(a,n)        // a → array of n-elements
{ solution = ∅               // no soln. in the beginning
  for(i=1; i≤n; i++)
    { x = select(n);         // finding the feasible solutions
      if (Feasible(x))       // checking for feasible soln.
        add(x, solution);
    }
}
```

Time complexity depends upon :-
- Select(n)
- Feasible(x)
- add(x, solution)

Best Case :- $\Omega(n)$ [When all three fn. take $O(1)$ tim
- example :- Select(n) $\longrightarrow n$
  feasible(x) $\longrightarrow n^2$
  add(x, solution) $\longrightarrow n^4$

  ∴ complexity :- $\boxed{O(n^5)}$ (because of for loop)

★★ Q. The time complexity of Greedy Technique :-
(a) $O(n)$ (b) $O(n^2)$ (c) $O(n^3)$ (d) $O(n^4)$

→ though we can't say anything about the complexity of the functions, so it may even be $O(n^{100})$ but from the available option (d) is the most appropriate option.

Applications Of Greedy :-
1. Job Sequencing with Deadline
2. Knapsack problem
3. Huffman Coding
4. Optimal Merge Pattern
5. Minimum Cost Spanning Tree
   (i) Prim's

Cosmos

⑥ Single Source Shortest Path
(i) Dijkstra's
(ii) Bellman-Ford

Job Sequencing with Deadline

① Single-CPU :- at a time only single job is executing (Round-Robin is not possible).
② Interleaving is not allowed
③ Arrival times of all jobs are same. (FCFS is not possible)
④ 1-unit of running time for every job. (SJF is not possible).

	J_1	J_2	J_3	J_4
Deadline	2	1	2	1
Profit	50	75	45	80

$(J_4 \rightarrow J_1)$ → Optimal soln.

Soln. Space :-

$J_1\ J_2$ J_2J_1 J_3J_1 J_4J_1 J_1 $\{\}$
$J_1\ J_3$ J_2J_3 J_3J_2 J_4J_2 J_2 ↑
J_1J_4 J_2J_4 J_3J_4 J_4J_3 J_3 no job at all.
J_4

$J_1J_2J_3$ 3 jobs → ${}^4P_3 = \frac{4!}{1!} = 4!$

$J_1J_2J_3J_4$ / J_2 4 jobs → ${}^4P_4 = \frac{4!}{0!} = 4!$

∴ Soln. space :- $17 + 2 \times 4!$
$= 17 + 48$
$= \boxed{65}$

$\{J_2, J_1\}$ → 125
$\{J_2, J_3\}$ → 120
$\{J_4, J_1\}$ → (130)
$\{J_4, J_3\}$ → 125
$\{J_1, J_3\}$ → 95
$\{J_3, J_1\}$ → 95
$\{J_1\}$ → 50
$\{J_2\}$ → 75
$\{J_3\}$ → 45
$\{J_4\}$ → 80
$\{\}$ → 0

→ Feasible Solution

Optimal Soln. :-

J_2, J_1
J_2, J_3
J_4, J_1
J_4, J_3
J_1, J_3
J_3, J_1

Objective :-

complete as many jobs as possible in the deadline.

Cosmos

To calculate Optimal Soln. directly :-

1	2
J_4	J_1

Step (i) :- jobs that can be done in 2nd month. (J_1 & J_2), we take J_1 for max. profit

Step (ii) :- jobs that can be done in 1st month (all jobs can be done in 1st month, but J_1 is done in slot 2). So, only J_2, J_3, J_4 are in competition, J_4 is giving max. profit

Jobs	J_1 ✓	J_2 ✓	J_3 ✓	J_4	J_5 ✓	J_6 ✓
Profits	24	18	22	10	12	30
Deadlines	5	3	3	2	4	2

$n = 6$

J_6	J_2	J_3	J_5	J_1
1	2	3	4	5

↑ start from this side.

$n = 5$

Jobs	J_1	J_2 ✓	J_3 ✓	J_4	J_5 ✓
Profits:	10	20	15	5	80
Deadlines:	3	3	3	4	4

J_1	J_2	J_3	J_5
1	2	3	4

↑ start from this side

Profit = 125

(i) :- Sort all the jobs in decreasing order of profit.

J_5	J_2	J_3	J_1	J_4
80	20	15	10	5
4	3	3	3	4

Stop here

Cosmos

Q. n=9

Jobs	J_1	J_2	J_3	J_4	J_5	J_6	J_7	J_8	J_9
Profits:	15	20	30	18	18	10	20 (23)	16	25
Deadlines	7	2	5	3	4	5	2	7	3

Sort acc. to profits

Jobs	J_3	J_9	J_7	J_2	J_4	J_5	J_8	J_1	J_6
Profits	30	25	23	20	18	18	16	15	10
Deadlines	5	3	2	2	3	4	7	7	5

J_2	J_7	J_9	J_5	J_3	J_1		J_8
1	2	3	4	5	6		7

Profit = 147

J_4 & J_6 are left o

Algorithm

① Sort all jobs in decreasing order of profits → O(nlo
② Find maximum deadline in the array of n-deadlin & take array of that size, & start from RHS. → O(n
③ For every slot no. 'i', apply linear search to find job which contain deadline ≥ 'i'. → $O(n^2)$

Knapsack Problem :-

problem: $\sum_{i=1}^{n} w_i > M$

M :- weight that knapsack can hold

w_i :- weight of i^{th} object.

Feasible Soln. :- $\sum_{i=1}^{n} x_i w_i \leq M$

Optimal Soln. :- $\sum_{i=1}^{n} x_i p_i$ is maximum.

★ (Fractions are allow in Greedy)
★ When fractions are n allowed, it is shifted Dynamic prog

now,
n = 3 M = 20
objects ob, ob, ob,

My answer

Optimal Soln. :- 18×25 + 2 ×

★ All 3 are feasible solutions.

① Greedy About Profit:-

$1 \times 25 + \frac{2}{15} \times 24 + 0 \times 10$

$= 25 + 3.2 = \boxed{28.2}$

② Greedy About Weight:-

Take less weight object 1st:-

$(0, \frac{10}{15}, 1)$ $\boxed{10}$ 10

- Take 3rd object 1st
- Now, out of 1st & 2nd object, 2nd object has less weight.

$\therefore$ Profit :- $0 \times 25 + \frac{10}{15} \times 24 + 1 \times 15$

$= 16 + 15 = \boxed{31}$

Greedy About Both Profit & weight :- (Sir's answer)

obj 1 :- 18 weight → 25 profit
1 " → $\frac{25}{18} = 1.3$

obj 2 :- 15 weight → 24 profit
1 " → $\frac{24}{15} = 1.6$

- take obj 2 first, then obj 3 & then obj 1.

obj 3 :- 10 weight → 15 profit
1 " → 1.5 "

$(-, \frac{1}{2}, \frac{5}{10})$

profit = $0 \times 25 + 24 \times 1 + 15 \times \frac{1}{2} = 24 + 7.5 = \boxed{31.5}$

Greedy Knapsack will give Optimal soln. always by

Q. n=5 M=12

Objects	Ob1	Ob2	Ob3	Ob4	Ob5
profits	5	2	2	4	5
weights	5	4	6	2	1

Ob1 :- 5 weight → 5 profit
1 " → 1 "

Ob2 :- 4 weight → 2 "
1 " → 0.5 "

Ob 3 :- 6 " → 2 "
1 " → 0.33 "

Ob 4 :- 2 " → 4 "
1 " → 2 "

Ob 5 :- 1 " → 5 "

profits

Ob5 → 5 — 11
Ob4 → 4 — 9
Ob1 → 5 — 4
Ob2 → 2 — 0
16

Objects → (1, 1, 0, 1, 1)
Profit = 16

Q. n=7 M=15

Objects	ob1	ob2	ob3	ob4	ob5	ob6	ob7
profits	10	5	15	7	6	18	3
weights	2	3	5	7	1	4	1
profit:weights	5	1.6	3	1	6	4.5	3

Objects = (1, 2/3, 1, 0, 1, 1, 1)

Profit = 10 + 2/3 × 5 + 15 + 6 + 18 + 3
= 52 + 3.2
= 55.2

Algorithm

for (i = 1 to n)

pi - profits
wi - weights

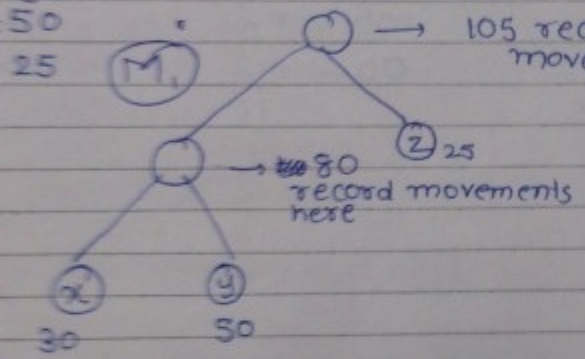
Cosmos

Take one by one object untill capacity of knapsack beco
es zero. ⟶ $O(n)$

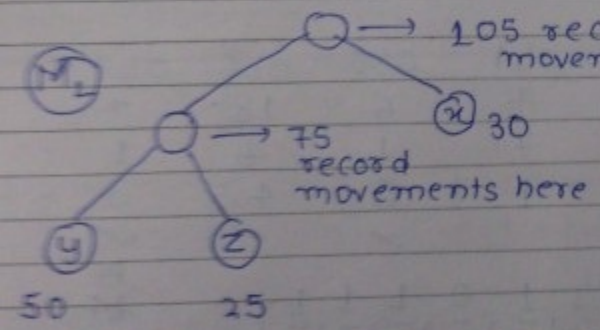
∴ Time complexity = $O(n) + O(n \log_2 n) + O(n)$
= $O(n \log_2 n)$

Optimal Merge Pattern

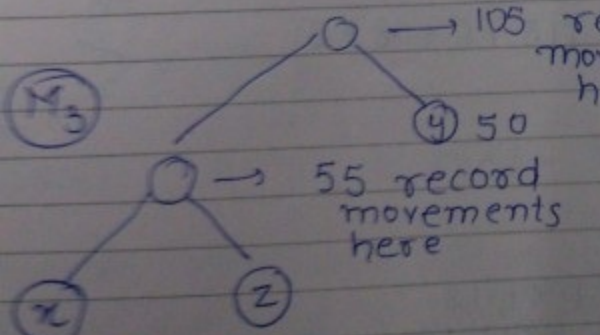
5 files
30
50
25



∴ total record movements =
80 + 105 = 185 record movemen



∴ total record movements = 75 + 105
= 180 record movem



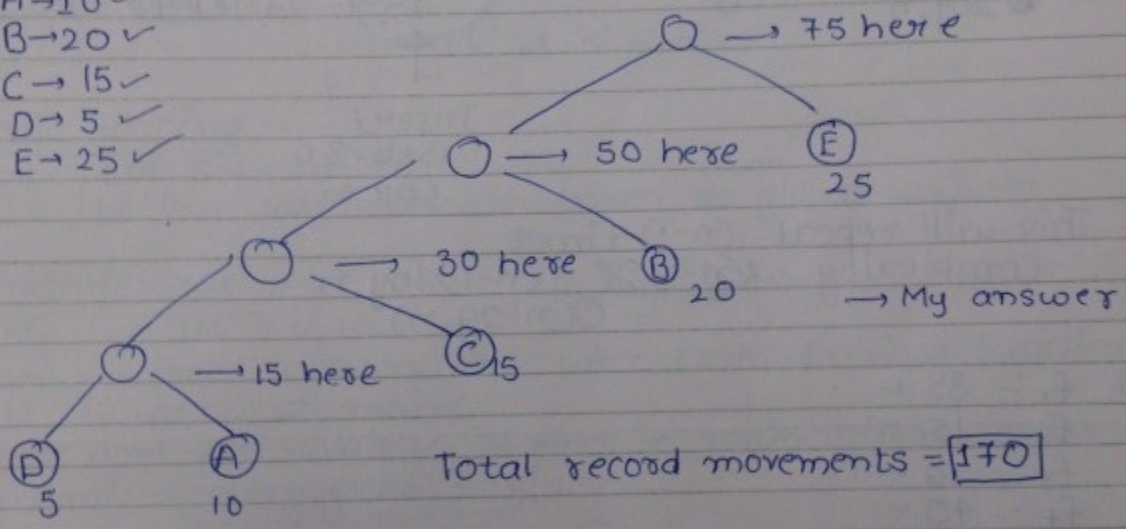
∴ total record movements = 55 + 1
= 160 record movem

Cosmos

For the given problem, there are 3 merge patterns, M_1, M_2, M_3, in all of them M_3 is optimal merge pattern because it is taking least record movements.

he Min. no. of record movements required to merge 5 files

A → 10 ✓
B → 20 ✓
C → 15 ✓
D → 5 ✓
E → 25 ✓

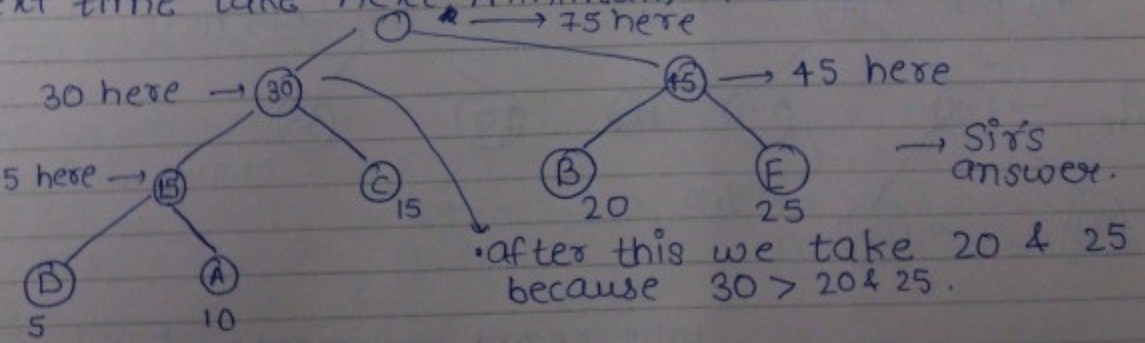


Total record movements = 170

★

1st

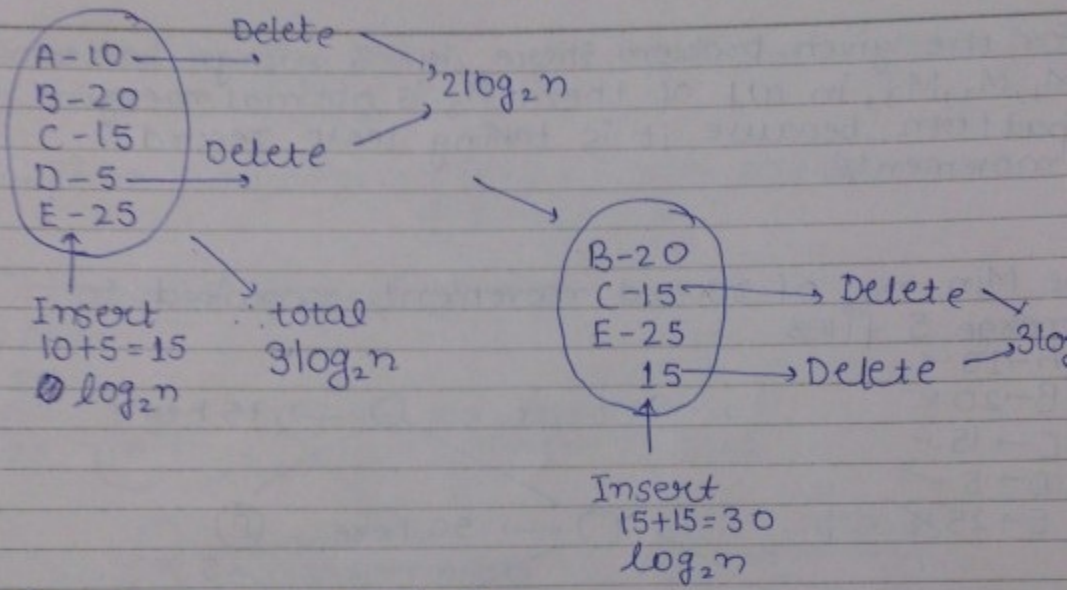
r these problems, take two minimum first time. ext time take next minimum & so on.



after this we take 20 & 25 because 30 > 20 & 25.

75 + 45 + 30 + 15 = 165

Cosmos



This will repeat $(n-1)$ times

Complexity :- $3(n-1)\log_2 n$

$= O(n\log_2 n)$

Q. f_1 :- 35 ✓

f_2 :- 15 ✓

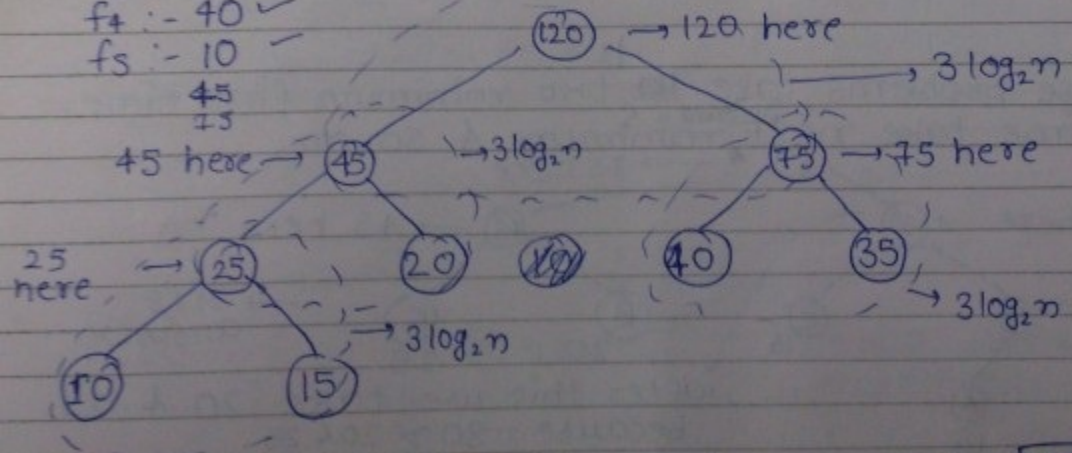
f_3 :- 20 ✓

f_4 :- 40 ✓

f_5 :- 10 ✓

45

15



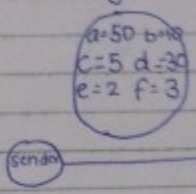
total record movements = 265

Cosmos

Huffman Coding (Huffman coding is best technique because it is application of greedy & hence it gives best optimal solution).

1) Data Encoding
2) Data Compression

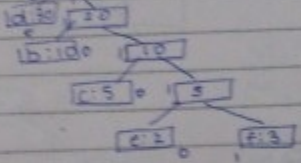
Message contains 100 characters



If each character is transferred through ASCII, then 7 bits req.
total bits to transfer 100 characters = 7×100 = 700 bits

compressing

We need to transfer 6 characters, then 3 bits req.,
∴ total bits to transfer 100 characters = 3×100 = 300 bits

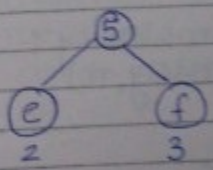


a → 0	1×50	→ 50
d → 10	2×30	→ 60
b → 110	3×10	→ 30
c → 1110	4×5	→ 20
e → 11110	5×2	→ 10
f → 11111	5×3	→ 15
		185

Huffman coding

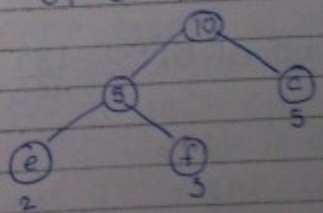
Note:-
More freq ⇒ less bits
Less freq ⇒ more bits

Step 1:-



a=50	b=10
c=5	d=30
e	5

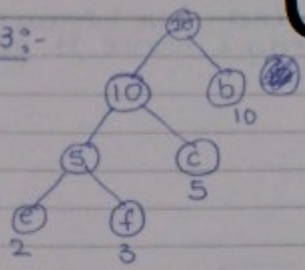
Step 2:- Now 2 minimums are c=5 & 5 from summation of e & f.



a=50	b=10
d=30	10

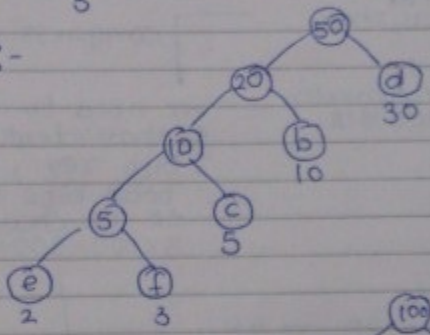
Cosmos

Step3:-



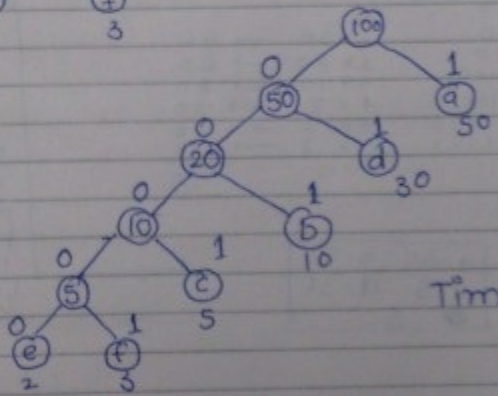
a=50
d=30
20

Step4:-



a=50
50

Step5:-



← empty

Time Complexity :-
$O(n \log_2 n)$

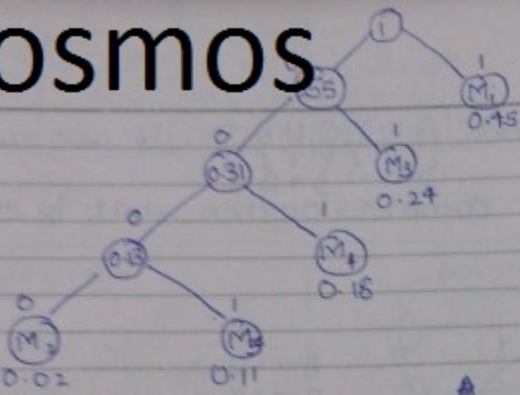
Huffman Coded Tree

a → 1
b → 001
c → 0001
d → 01
e → 00000
f → 00001

→ 185 bits
Average no. of bits = $\frac{185}{100} = 1.85$

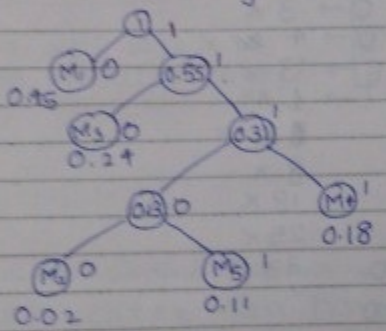
Q. Message M contains
M = (M1, M2, M3, M4, M5)
0.45, 0.02, 0.24, 0.18, 0.11

Cosmos



↑ My answer

$M_1 \rightarrow 1$ → 1 × 0.45 = 0.4
$M_2 \rightarrow 0000$ → 4 × 0.02 = 0.0
$M_3 \rightarrow 01$ → 2 × 0.24 = 0.4
$M_4 \rightarrow 001$ → 3 × 0.18 = 0.5
$M_5 \rightarrow 0001$ → 4 × 0.11 = 0.



↑ Sir's answer

$M_1 \rightarrow 0$
$M_2 \rightarrow 1100$
$M_3 \rightarrow 10$
$M_4 \rightarrow 111$
$M_5 \rightarrow 1101$

★ Put smaller no. on left hand side & bigger no. on right hand side.

Total no. of bits:
$0.45 \times 1 + 0.02 \times 4 + 2 \times 0.24 + 3 \times 0.18 +$
$0.11 \times 4 = 1.99$ bits
Avg. no. of bits = 1.99 bits.

In Huffman Coding
★ One message will contain 1 bit (but not always)
★ There are two messages with same no of bits.

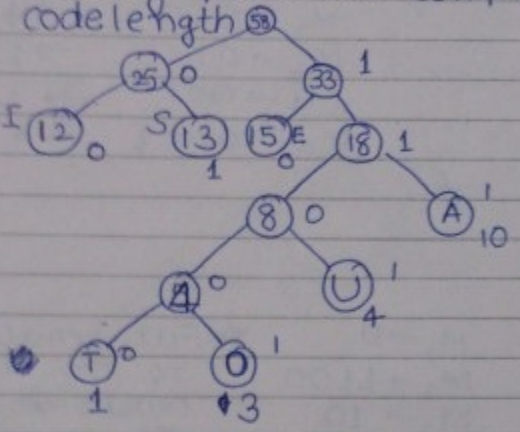
Q Decode the mes encoded message :-
10110011011111001100101111101

Steps: See the tree & start from root.
1 → go right from root
10 → go left from 0.55 & we reach leaf,
M_3 is detected.
1 → go right from root

Cosmos

A file contains characters A, E, I, O, U, S, T, if we use Huffman coding for data compression, then what is the avg. codelength

A	E	I	O	U	S	T
10	15	12	3	4	13	1



A → 10 ×
E → 15 ×
I → 12 ×
O → 3 ×
U → 4 ×
S → 13 ×
T → 1 ×
4 ×
8 ×
18 ×
25 ×
33 ×

A → 111	→ 3 × 10	= 30	2
E → 10	→ 2 × 15	= 30	
I → 00	→ 2 × 12	= 24	
O → 11001	→ 3 × 5	= 15	
U → 1101	→ 4 × 4	= 16	
S → 01	→ 2 × 13	= 26	
T → 11000	→ 5 × 1	= 5	
		146	→ total bits

Avg. codelength = $\frac{146}{58}$

Q. If $f(n) = O(g(n))$
T/F ?

a. $\log(f(n)) = O(\log(g(n)))$ ✓

b. $2^{f(n)} \equiv O(2^{g(n)})$ ✗ → take $f(n) = 2^{2n}$

c. $f(n) \equiv O(f(n/2))$ ✗ → take $f(n) = \frac{1}{n}$

d. $f(n) = O(f(n))^2$ ✗ → take $f(n) = n$

$f(n) = O(g(n))$

$n^2 < n^3$

$n^2 = O(n^3)$

~~$2\log n = \Theta(3\log n)$~~

(applying log on both sides)

$n! = O(n^n)$

$\log n! = O(n \log n)$

$f\ n^3 = \Omega(n^2)$

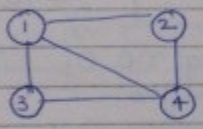
$3\log n = \Theta(2\log n)$

★ Sometimes they becomes equal on application of log on both sides.

Minimum Cost Spanning Tree

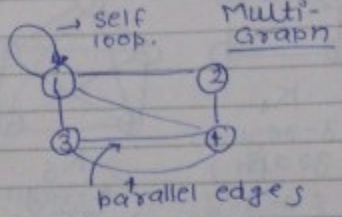
$G(V, E)$

Simple Graph



① No self loop.
② No parallel edges

Multi-Graph



① Self loops or parallel edges

Cosmos

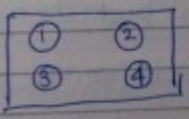
★ In a simple graph with 'n' vertices, then at max. each vertex can have $(n-1)$ degree.

★ In multigraph with 'n' vertices, then at max. each vertex can have infinite degree.

Types Of Simple Graphs :-

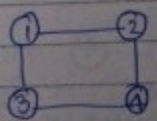
★ All null graphs are regular graphs but all regular graphs are not null graphs.

1. Null Graphs



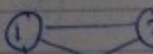
if the degree of every vertex is zero, then it is null graph.

2. Regular Graph → no. of edges = $\frac{n \times m}{2}$ (n → no. of vertices, m → degree of each vertex)



2-regular graph

if the degree of every vertex is same, then it is regular graph.

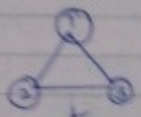


Cosmos

All complete graphs are regular but all regular graphs are not complete

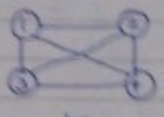
Complete Graph:-

If each vertex is connected to all other vertices, then it is called complete graph.



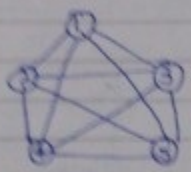
K_3 (2-regular graph)

$\frac{3 \times 2}{2} = 3$



K_4 (3-regular graph)

$\frac{4 \times 3}{2} = 6$



K_5 (4-regular graph)

$\frac{5 \times 4}{2} = 10$

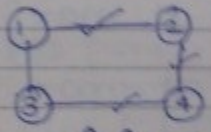
Maximal Graph

degree of each edge

$\frac{(n-1) \times n}{2}$ → total no. of vertices

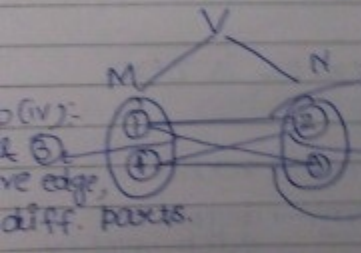
→ total no. of edges.

Bipartite Graph



V = {1, 2, 3, 4}

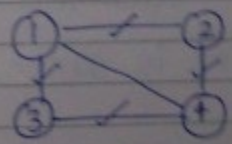
divide the vertices into two parts M & N & all the edges must be b/w M & N only & not within M only or N only.

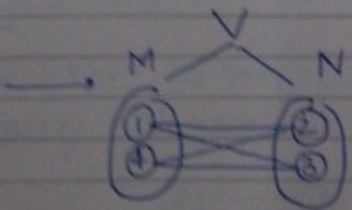


Step (i) :- ① & ② have edge ∴ put them in M & N respec.

Step (ii) :- ② & ④ edge, ∴ different parts

Step (iii) :- ③ & ④ edge, ∴ " "

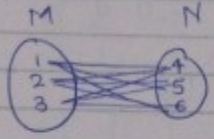




But ① — ④ not possible, ∴ it is not bipartite graph.

★ Connected Graph :-
In a graph each there is a path b/w every pair of vertices. Min. no. of edges = V-1 [V:-no. of vertices.]

Complete Bipartite



each element in M every element in M is adjacent to N, so it is complete bipartite graph.

$G(V, E)$

Cosmos

$|E| \leq \frac{V(V-1)}{2}$ [in simple graph] but possible for multigraphs too

★ $|E| = O(V^2)$

apply log on both sides:-

$\log E = O(2 \log V)$

$\log E = O(\log V)$

$\therefore E \log E \Rightarrow E \log V$ (from this)

No. of graps

n=2 ① ② ① ②

∴ no. of graphs = 2

n=3 ①

② ③

no. of graphs = $2^{n(n-1)/2}$ (simple)

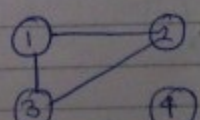
★ no. of graphs with 'n' vertices = $2^{\frac{n(n-1)}{2}}$

Spanning Tree :-

A connected subgraph H of a given Graph $G(V, E)$ is said to be spanning tree iff

① It should contain all vertices of G.

② Edge H should contain (V-1) edges if G contains V-vertices.

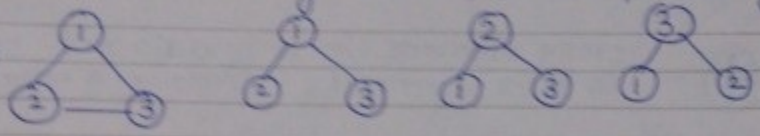
e.g. 

→ it contains all vetices & contains (V-1) edges, so it should be spanning tree,

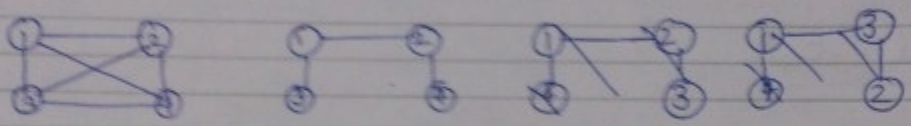
Cosmos

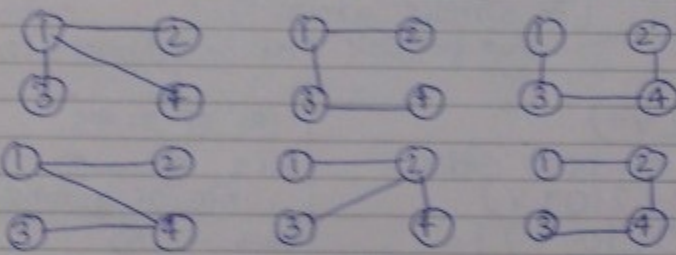
Q. Find no. of spanning trees for the following graph :-

(i)



(ii)

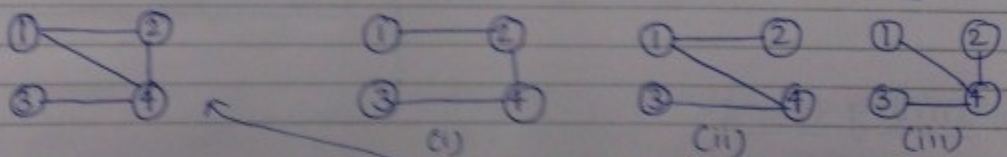




→ total = 16

☆ The no. of spanning trees in a 'n'-vertex complete graph = n^{n-2}.

Q. Find no. of spanning trees :- [when graph is not complete.]



(i) (ii) (iii)

Kirchaff-Theorem

$O(n^2)$

① Make Adjacency Matrix

Graph.

$$M = \begin{array}{c} \\ 1 \\ 2 \\ 3 \\ 4 \end{array} \begin{array}{c} \begin{array}{cccc} 1 & 2 & 3 & 4 \end{array} \\ \begin{bmatrix} 0 & 1 & 0 & 1 \\ 1 & 0 & 0 & 1 \\ 0 & 0 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix} \end{array}$$

② (i) replace all non-diagonal one 1's by -1.

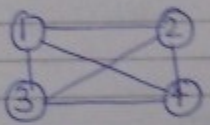
Cosmos

$$M = \begin{array}{c} \\ 1 \\ 2 \\ 3 \\ 4 \end{array} \begin{array}{c} \begin{array}{cccc} 1 & 2 & 3 & 4 \end{array} \\ \begin{bmatrix} 2 & -1 & 0 & -1 \\ -1 & 2 & 0 & -1 \\ 0 & 0 & 1 & -1 \\ -1 & -1 & -1 & 5 \end{bmatrix} \end{array}$$

③ Co-factors of any element will give no of spanni trees.

$$C_{11} = \begin{vmatrix} 2 & 0 & -1 \\ 0 & 1 & -1 \\ -1 & -1 & 3 \end{vmatrix} = 2(3-1) - 1(1)$$
$$= 2 \times 2 - 1 = \boxed{3}$$

Q Find no. of spanning trees for following graph:-



Step (i):- Adjacency Matrix

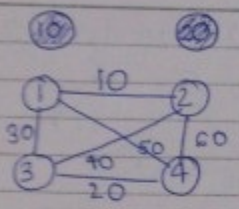
$$M = \begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \\ 1 & 1 & 1 & 0 \end{bmatrix}$$

Step (ii):-

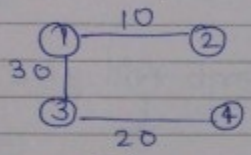
$$M = \begin{bmatrix} 3 & -1 & -1 & -1 \\ -1 & 3 & -1 & -1 \\ -1 & -1 & 3 & -1 \\ -1 & -1 & -1 & 3 \end{bmatrix}$$

$$C_{11} = \begin{vmatrix} 3 & -1 & -1 \\ -1 & 3 & -1 \\ -1 & -1 & 3 \end{vmatrix} = 3(9-1) + 1(-3-1) - 1(1+3)$$
$$= 24 - 4 - 4$$
$$= 16$$

Minimum Cost Spanning Tree
① Covering all vertices of the graph with min. cost of the tree edges.



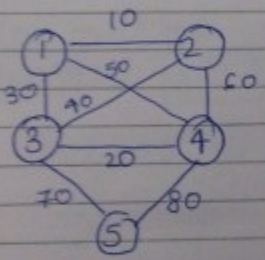
• for the graph 16 spanning trees are possible, in all of them M is the min. cost spanning tree.



M.

*To find out Min. Cost Spanning tree, we have 2 algorithms:-
i) Krushkal
ii) Prim's

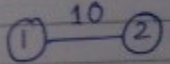
Krushkal :-



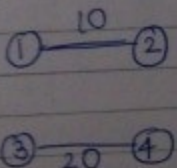
Step(i) :- Take 1st min.

edge weight	starting pt.	ending pt.
* 10	1	2
20	3	4
30	1	3
40	3	2
50	1	4
60	2	4
70	3	5
80	4	5

Now, make a min heap for all the edges, so no. of edges = E, ∴ time complexity = O(E)
Now take 1st min. (Delete it from heap).
Deletion from heap :- $O(\log_2 n)$ but $n = E$ ∴ $O(\log_2 E)$

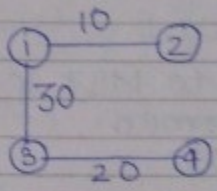


Step (ii) :- Take 2nd next min.

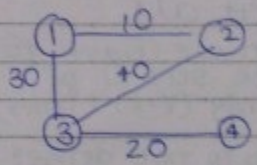


• In best case we will get min. cost spanning tree when 1st (V-1) edges do not make any cycle.
∴ (V-1) times $\log_2 E$ + O(E) = (V-1)$\log_2 E$ + O(E)
= O(V log E)

• take next min. :- (this also takes $\log_2 E$)

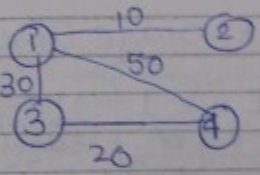


• take next min. :- (this also takes $\log_2 E$)



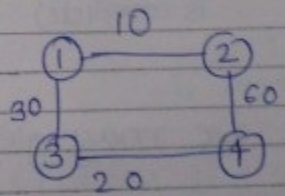
→ cycle → so delete this edge.

• take next min. :- (this also takes $\log_2 E$ from min. heap)



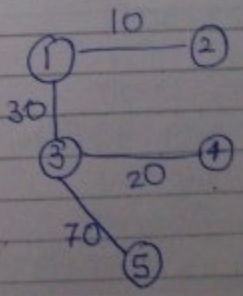
→ cycle, so delete this edge

• take next min. :-



→ cycle, so delete this edge

• take next min. :-



→ no cycle, & all vertices are covered, so this is min. cost spanning tree.

Cosmos

Algorithm:-

① Create Min. heap. $\longrightarrow O(E) \log_2 E$
② Delete one by one min. & add to MST if no cycle add to MST if no-cycle formed until (V-1) edges are added to MST.

Best Case:- When we get MST first (V-1) times without getting a cycle at all.
(i.e. When we have to execute step ② of algo only (V-1) times)

$E + (V-1)\log E$
$= E + V \log E$
but $\log E = O(\log V)$
$= E + V \log V$
$= \boxed{O(E + V \log V)}$ → Best Case & Average Case

When graph is not dense:- $\Omega(V \log V)$
When graph is dense:- $O(V^2)$

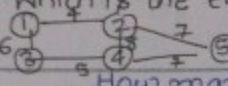
Worst Case:-
(When we have to execute step ② of algo E times.)
$E + E \log E$
$= E + E \log V$
$= \boxed{O(E \log V)}$

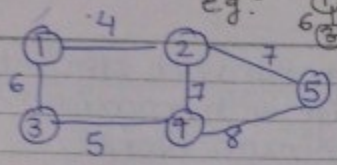
in best case:- $E = V-1$ → when graph contains (V-1) edges only.
in worst case:- $E = V(V-1) = O(V^2)$ (when graph is complete)

★ In Kruskal's algorithm, in middle we may get disconnected graph (Forest), but in case of Prim's algo, we always get connected graph.

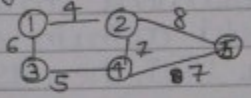
Prim's

Consider the following graph:-

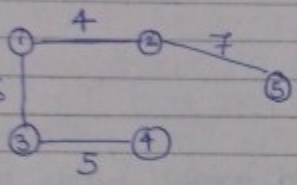
* if there are more than one MST for a given connected graph, then all the graph contain atleast one edge weight which is the edge weight of more than one edges.
e.g.  has 2 MST's, so it must contain atleast 2 edges with same weight.



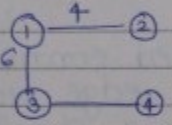
How many diff. min. cost spanning trees are possible?

* but converse is not true, i.e. if graph contains edges with same weight, then it might not have more than one MST,
e.g.  has only one MST.

Ans. Using Krushkal :-

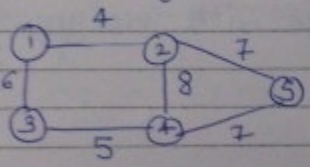


Only 1 at this stage :-

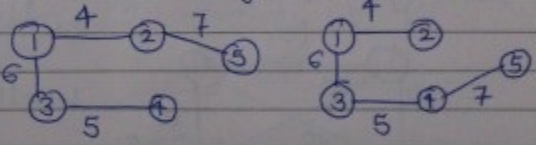


next min. is 7 which is b/w ② & ④ and ② & ⑤, but ② & ④ makes a cycle, so we can't use it, we can only use edge b/w ② & ⑤, & we stop because we have 4 edges now.

but if the graph is :-



∴ min spanning trees are :-



1

If there are

Note :- ① For the given graph more than 1 MST may be possible.

② For the given graph, if there are more than one MST, then atleast one of the edge weight will be repeated. (but converse is not true.) ← implies that *

✓ ③ If the given graph is connected & no edge weight is repeated, then exactly one MST is possible.

④ There will be no MST if graph is disconnected

★★ If edge weight is not given, then we assume it to be 1.

Q. Let G be an undirected connected graph with distinct edge weights. Let emax be the edge with maximum weight & emin be the edge with minimum weight, then check following statements are T/F?

T (a) Every MST should contain emin.

T (b) If emax is in MST, then its removal must disconnect G.

F (c) No MST contain emax.

T (d) G has unique MST.

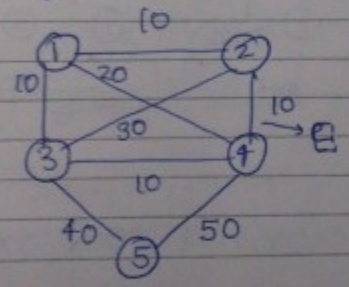
Q. Let G be an undirected connected graph with N vertices. If 'W' is the min. edge weight among all edge weights in G & E be a specific edge with weight 'W'. Then

(a) E may be there in MST.

(b) If E is not in MST, then all edges are having weight 'W' in that cycle.

(c) Every MST should contain an edge with weight 'W'.

(d) Every MST should contain E.



W = 10

edge weight is not given, assume it to be 1.

An undirected graph G has 'n' nodes, the adjacency matrix is given by n×n square matrix. Whose

(i) Diagonal elements are zero's.
(ii) Non-Diagonal elements are one's } → connected graph.

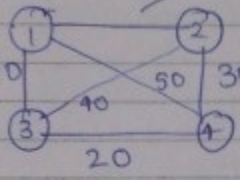
then check :-

(a) G has no MST.

(b) G has multiple MST's with diff. costs. :- multiple MST's with same cost may be possible.

(c) G has unique MST of cost n-1.

→ if there is no edge weight, we assume it to be 1.



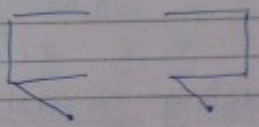
→ MST
↓
31

Cosmos

Q. Let T & T' be 2-spanning trees of a connected graph G. Suppose that an edge e_1 is in T but not in T' & edge e' in T' but not in T, then which is T after performing following op^n on T & T'. (remaining edges are assumed to be same)

(i) $(T - \{e\}) \cup \{e'\}$

(ii) $(T' - \{e'\}) \cup \{e\}$



(a) Both T & T' are ST.

(b) T → ST but not T'.

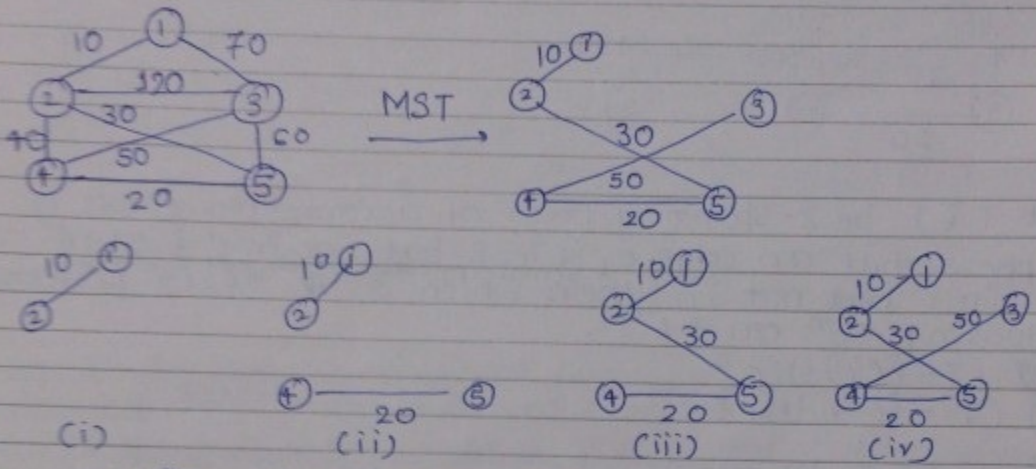
(c) T' → ST but not T.

(d) both are not ST.

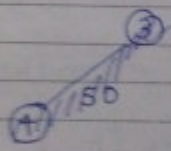
My answer:- removing e from T & adding e' makes it T' & removing e' from T' & adding e makes it T.

☆ if the remaining edges are not same:-

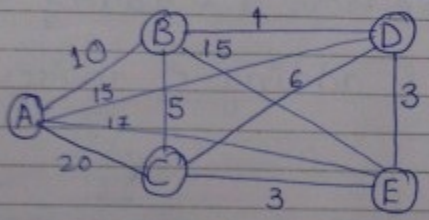
Prim's Algorithm



Using Kruskal algo



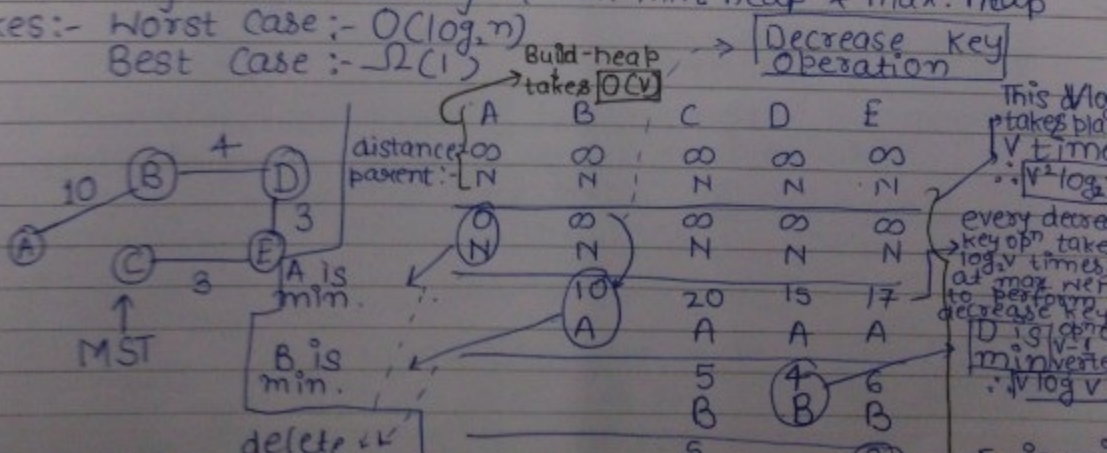
Apply Prim's algorithm :-



★

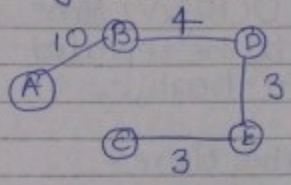
Decrease key & Increase key opn in min heap & max. heap takes:- Worst case :- $O(\log_2 n)$

Best case :- $\Omega(1)$

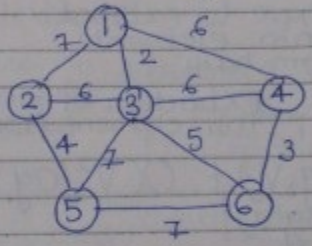


Now,

Using Kruskal's



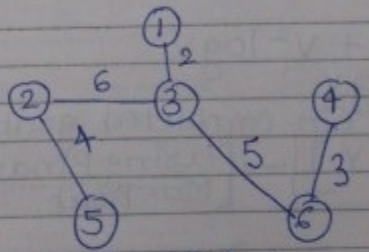
Apply Prim's algo on the following graph:-



distance:- ∞ ∞ ∞ ∞ ∞ ∞

parent:- N

0	~~7~~	6	2	~~6~~ 3	~~7~~ 4	5
1	2	3	4	5	6
~~N~~	~~1~~	~~1~~	~~1~~	~~1~~	~~1~~
X	1	1	1	8	3
	3	X	6	2	X
	~~6~~		X		
	X				



using Prim's.

Note:-

Step (i) Build-heap of all the vertices :- takes O(V) time.

(ii) Take min. out of the min-heap :- takes O(1) time, rearranging heap will takes :- $O(\log_2 V)$

(iii) Now, change the distance value using decrease key opⁿ :- one decrease key opⁿ on one vertex takes $\log_2 V$ time.

(iv) Now, at max. in one stage all the vertices will have to undergo decrease key opⁿ, so time taken :- $O(V \log_2 V)$.

(v) Now, we take out next min. from heap & have to rearrange the heap, it takes :- $O(\log_2 V)$ time.

(vi) Now, the decrease key op^n have to be repeated till only one vertex is left in min-heap, ∴ the total no. of stages = V.

∴ total decrease key op^n takes time :-

$V \times V\log_2 V$

- V times decrease key op^n
- time complexity of decrease key op^n in one stage.

We have to extract 'V' min. elements from heap

extraction of next min. element from heap

My answer :-

∴ Time Complexity = $O(V \cdot \log_2 V) + O(V^2 \log_2 V) + O(V$

$= \boxed{O(V^2 \log V)}$

Sir's Answer :-

Time Complexity = $V + V\log V + V^2 \log V$

$= (V + V^2)\log V$

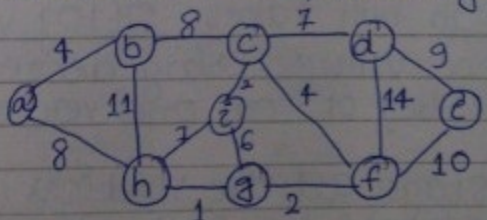
But $V^2 = E$ (in completed graph)

$\boxed{O[(V+E)\log V]}$ → [using Binary Min. heap.]

∴ Time Complexity :- $O[E + V\log V]$ [using Fibonacci min heap.]

:- $O(V+E)$ [using Binomial min heap.]

Consider the following graph



Cosmos

Which one of the following sequence of edges of MST is not true using Prim's algo when the algo. started from vertex a.

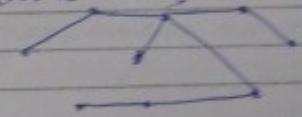
(a) (a,b) (b,c) (c,i) (c,f) (f,g) (g,h) (c,d) (d,e)
(b) (a,b) (b,c) (f,c) (c,i) (f,g) (g,h) (c,d) (d,e)
(c) (a,b) (a,h) (h,g) (g,f) (f,c) (c,i) (c,d) (d,e)
(d) (a,b) (b,c) (h,g) (f,g) (c,i) (c,f) (c,d) (d,e)

(d) is wrong because it creates unconnected graph.

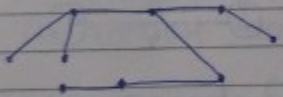
		4	8	7		4	6	7 8 2
distance	0	∞	∞	∞	∞	∞	∞	∞ ∞
		a	b	c	d	e	g	h i
parents	N	N	N	N	N	N	N	N N
		a	b	c	c	i	a i	c

(b) is wrong because using Prim's after (a,b) & (b,c) → (c,i) should come but (f,c) is coming.

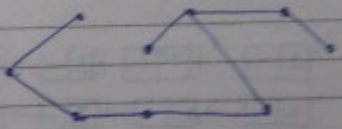
(a)



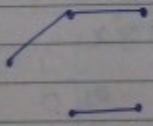
(b)



(c)



(d)



← This is wrong because it created undi graph.

(a) & (c) are correct using Prim's.

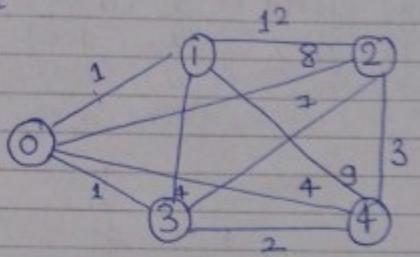
Q. Consider a complete undirected graph with vertex set V = {0, 1, 2, 3, 4}
Entry Wij in matrix W below is the weight of t edge {i, j}

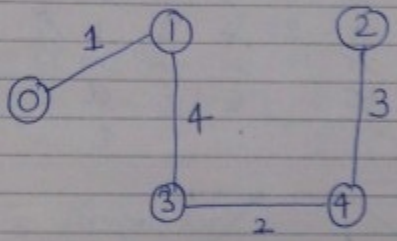
W =	0	1	2	3	4
0	0	1	8	1	4
1	1	0	12	4	9
2	8	12	0	7	3
3	1	4	7	0	2
4	4	9	3	2	0

What is the cost of MST for the above graph such that vertex 0 is a leaf node in that spanning tree?

Cosmos

(a) 8
(b) 9
(c) 10 :- My answer
(d) 11

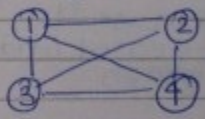


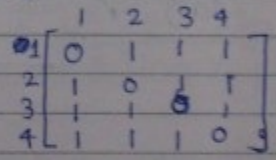


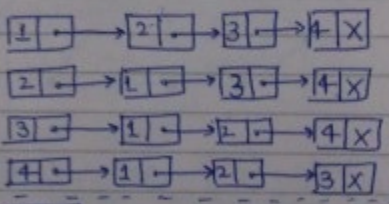
(using Kruskal's)

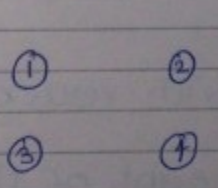
Graph Representations

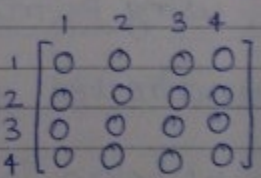
① Adjacency Matrix

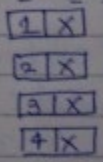






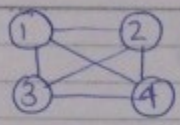


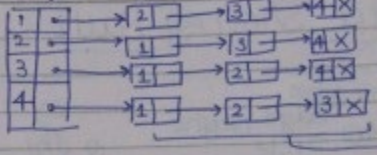




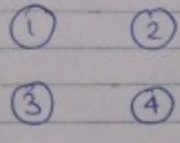
- represent graph Adjacency matrix will take $O(V^2)$ always.
- find out the degree of a vertex, it will take $\theta(V)$.
- To find out that the graph is connected, it will take.

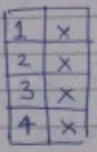
Adjacency List:-





Space Complexity:-
$V + 2E = \boxed{O(V+E)}$
or
$\boxed{\Theta(V+E)}$
Whichever is big, it will be the answer.





In this case :-
* To find out the degree of vertex :-
  $\Omega(1)$ → When no adjacent neighbours.
  $O(V)$ → When graph is connected.
  (Sparse graph)
* When less no. of edges :- Adjacency List
* When more no. of edges :- Adjacency Matrix
  (dense graph)

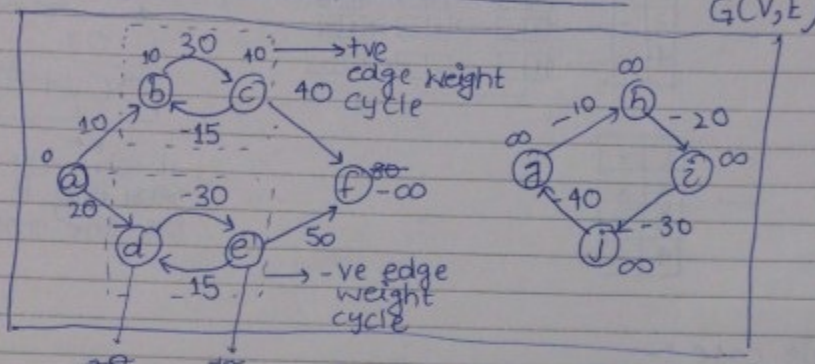
Single Source Shortest Path

Date

21.07.12

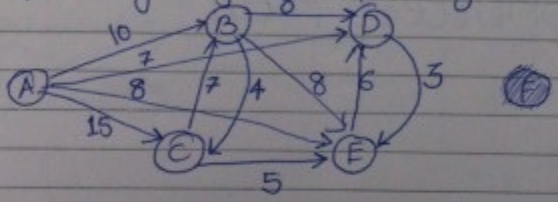
Cosmos

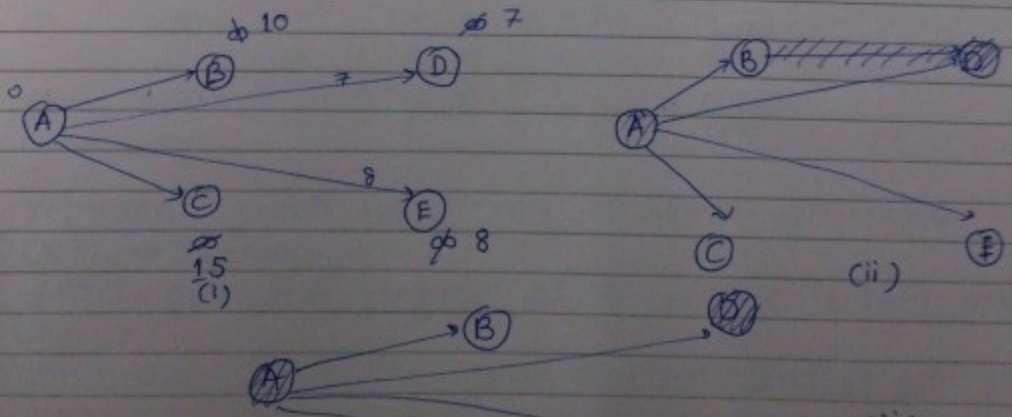
Single Source Shortest Path



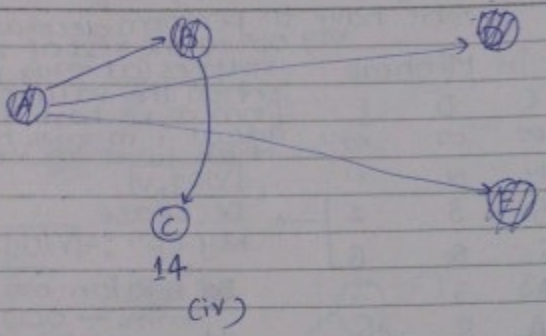
- MC(A,B) = ∞, if there is no path b/w A & B
- MC(A,B) = -∞, B is participant of negative edge weight cycle or dependent on -ve edge weight cycle
  - d, e → participants
  - f → dependent

Q. Find shortest path from vertex 'a' for the following graph using dijkstra's algo?

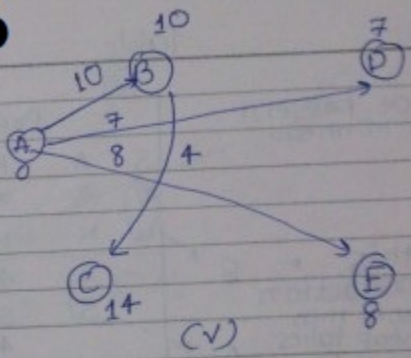




Cosmos

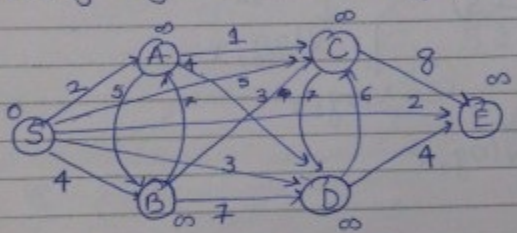


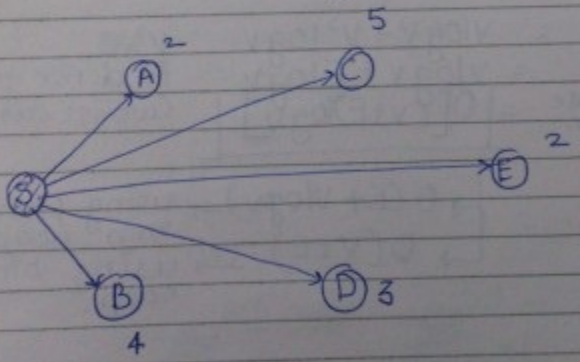
(iv)

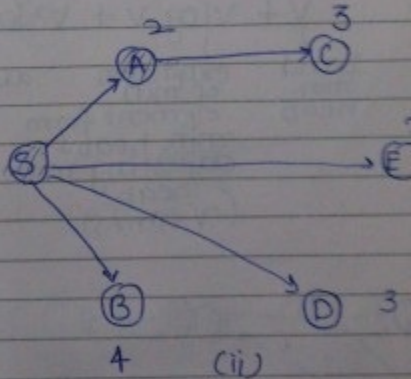


(v)

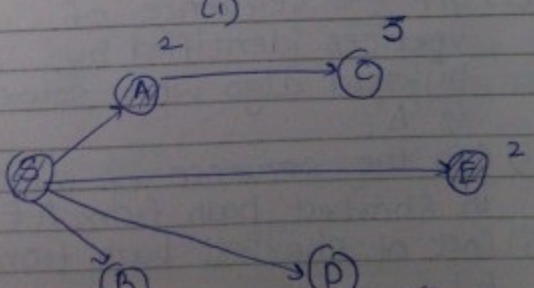
Q. Apply Dijkstra starting from vertex 'S'.



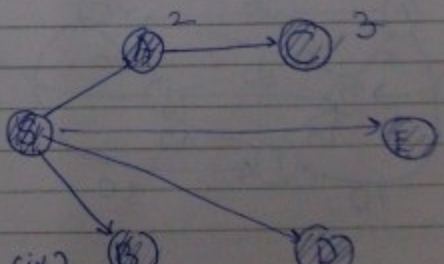




(i)



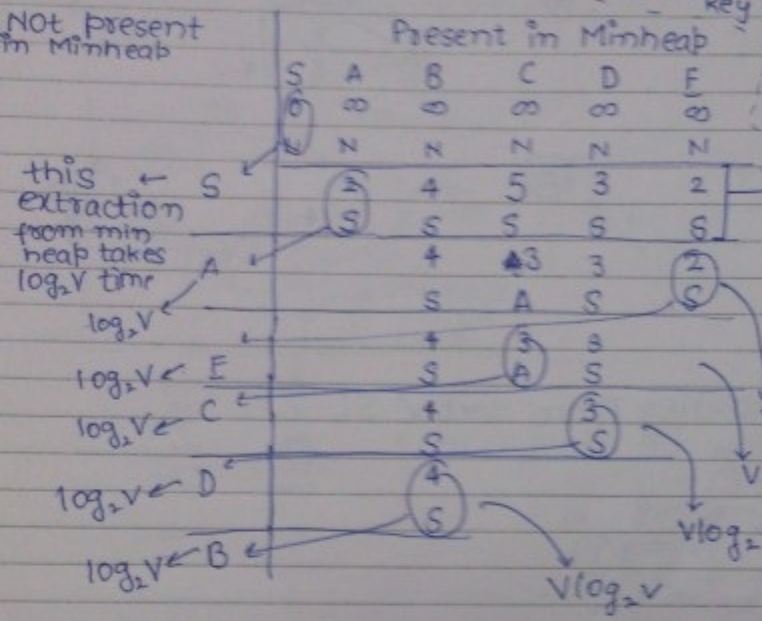
(ii)



Cosmos

Explanation

Table



Decrease Key opn :-

- Delete the min. from min. heap → O(1)
- Adjust the root :- $\log_2 V$
- We have to perform decrease key opn to the rest of the vertices (at max.)
- When the value changes from ∞, we have to adjust it in min. heap, which takes $\log_2 V$ time.

$V \log_2 V$

Decrease key opn :- $V \log V$

Deletion on min. → O(1)

at max. we have to perform decrease opn. (V-1) times.

(V) $\log_2 V$.

Time Complexity:-

$V + V\log_2 V + V^2 \log_2 V = V \log V + V^2 \log V$

$= V \log V + E \log V$

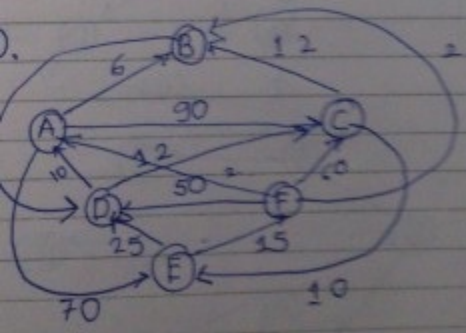
$= O[(V+E) \log V]$

$V^2 = E$ in dense graphs (worst case)

build min. heap — extraction of min. element from min. heap & adjusting min heap (V times) — all decrease key opn

↳ $O(E + V \log V)$ → using fibonacci min heap

↳ $O(V+E)$ → using binary heap.

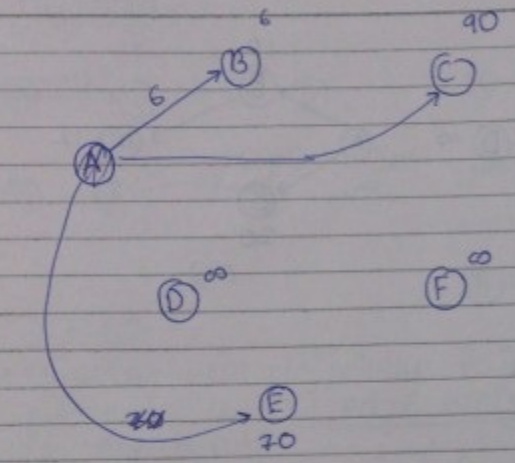


(i) O/p the sequence of vertices identified by Dijkstra algo when source is 'A'.

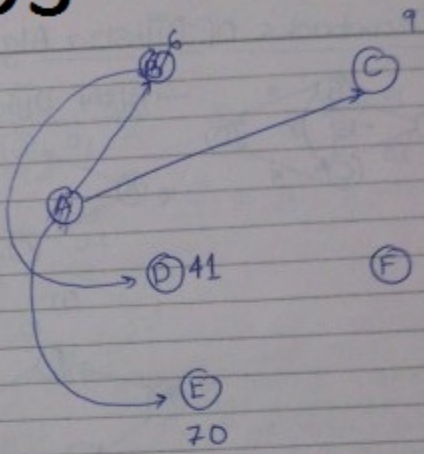
(ii) O/p the sequence of vertices in shortest path from A-E.

(iii) Cost of shortest path from A-E?

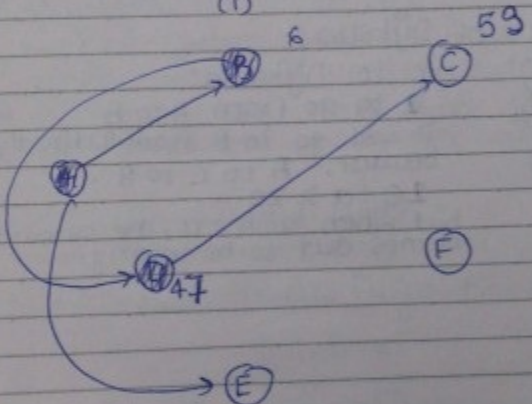
Cosmos



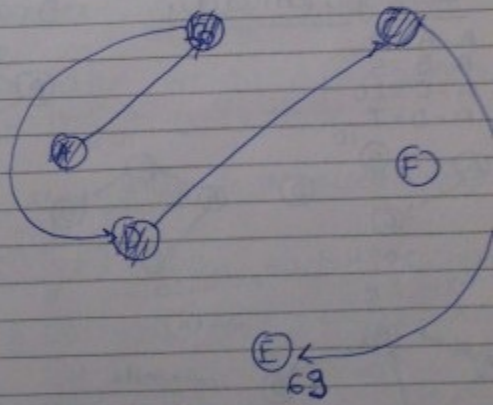
(i)



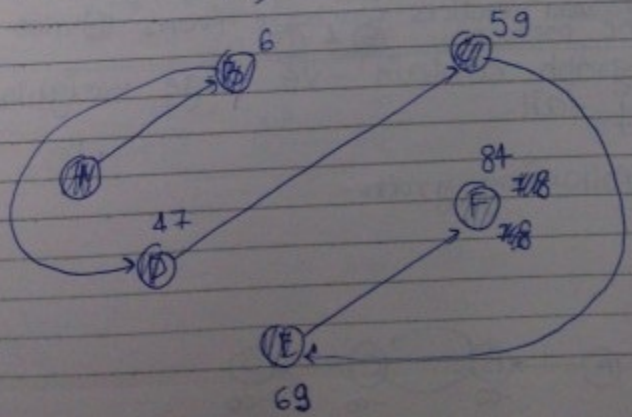
(ii)



(iv)



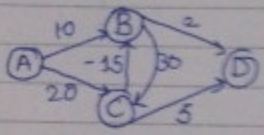
(v)



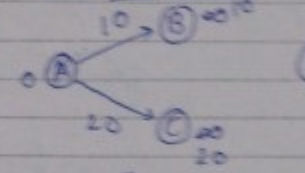
Ans. (i) A-B-D-C-E-F
(ii) A-B-D-C-E
(iii) ~~60~~ 69

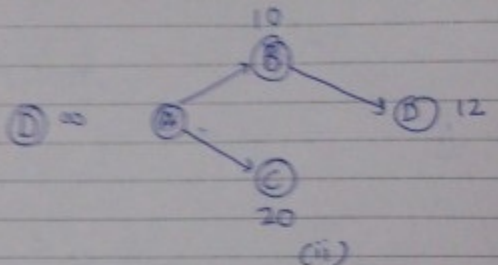
Cosmos

Drawbacks of Dijkstra Algo

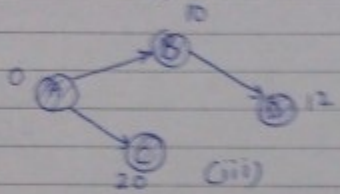


→ Apply Dijkstra:-



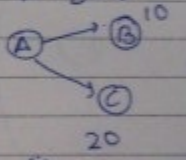


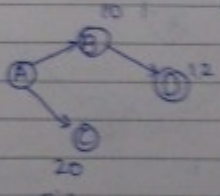
(i)

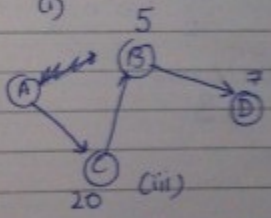


→ Using Manually

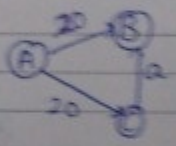
A - A = 0
A - B = 5
A - C = 20
A - D = 7







Drawback of Dijkstra:-

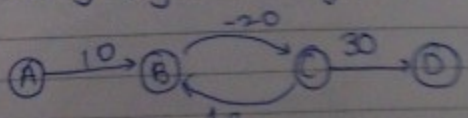


acc to Dijkstra:-
2 to go from A to B
It will go to B from A directly because:- A to C to B :-
20 + a > 20
but when a is -ve, the answer comes out to be wrong.

→ Whi
→ Which vertices gives incorrect result using Dijkstra:- to go to B:- A to C to B = 5
→ then ~~vertices~~ vertices reachable from B also faces the problem.

★ Note:- If the given graph contain -ve edge weights then Dijkstra algo may fail.

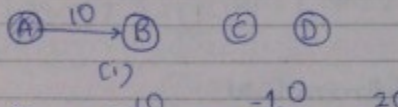
Q. Apply Dijkstra Algo for following graph:-



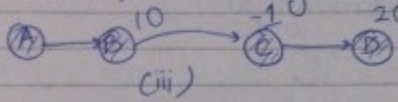
using manually:-

Cosmos

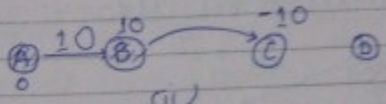
Dijkstra :-



(i)

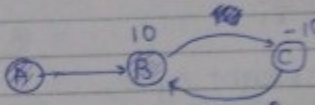


(iii)



(ii)

How many vertices gives wrong answer :-



15

but we can't take this edge because B is out of the min. heap. ∴ B is affected, & hence we can go to C via B, ∴ C is also affected & we can go to D via C, ∴ D is also affected.

Note :-

★ Dijkstra will defmitely fail when there is a -ve edge weight cycle in the graph.

★ Dijkstra will work fine for -ve edge weight when we eliminate a vertex after using the vertex 2 times.

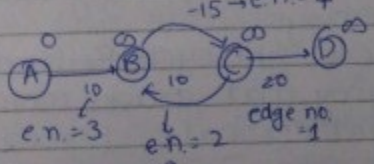
★ Dijkstra will work fine for -ve edge weight cycle when we eliminate a vertex after using it (V-1) times.

Bellman-Ford Algo :-

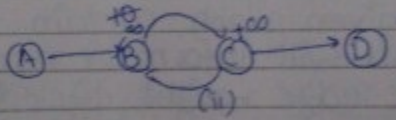
the algo is repeated (V-1) times.

Note :- ① Bellman Algo takes O(VE) time. → each time all the edges are calculat using bellman ford

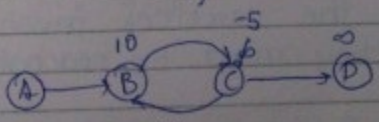
Q.



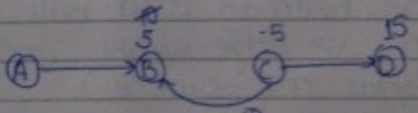
(i)



(ii)



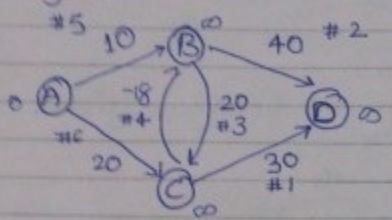
(iii)



(iv)

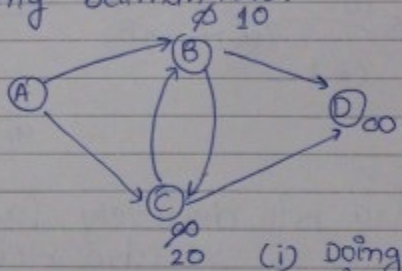
Cosmos

2. Apply Bellman-Ford algo :-



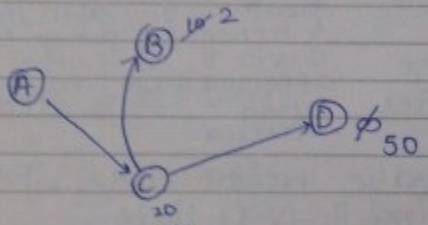
V=4 → we have to do (V-1)=3 times.

→ using Bellman Ford :-

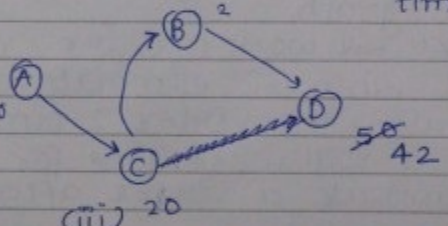


(i) Doing 1st time

-ve (no -ve edge weight cycle, only -ve edge weights, we have to do it only 2 times ~~& not~~ (V-1) times.)



(ii) doing 2nd time

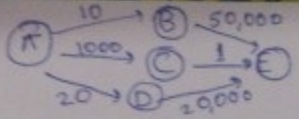


(iii) doing 3rd time

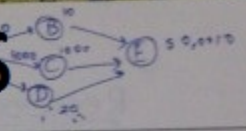
★ after performing V-1 times, do it one more time to check, if the values change then there is -ve edge weight cycle, otherwise not.

★★ When there is
- If given graph contain both +ve & -ve edge weight, bellman ford will give correct result always.
- If given graph contain -ve edge weight cycle, Bellman ford will inform (or report) saying that graph contain -ve edge weight cycle & we can't compute shortest path.
- The Bellman Ford will also give the vertices involved in -ve edge weight cycle. (but they must be reachable from the source.)

Dynamic Programming



Cosmos



Greedy :- ① Sometimes give wrong answer. ② Takes less time. ③ One decision at a time.

Dynamic : ① Always give correct answer.
② Take More time because it cover all possibilities.
③ Sequence of decisions.

- In greedy to choose shortest distance from A to E, greedy will choose path A to B (shortest path from A), & then choose B to E, which gives wrong answer.
- In dynamic programming, it will search every possibility & hence take more time.

★★ Greedy guess the answer which may or may not be correct.

Applications Of Dynamic Programming (Recursive)

① Fibbonacci Series
② Longest Common Subsequence
③ Multistage Graph
④ Matrix Chain Multiplication
⑤ Travelling Sales Person
⑥ 0/1 Knapsack
⑦ All pair shortest path
⑧ Sum of subset problem.

★ In divide & conquer → height of tree :- $\log n$
★ In Dynamic Programming → height of tree :- n

when we will consid all possibl combinatio

Fibonacci - Series :-

n	0	1	2	3	4	5	6	7	8	9	10
fibn	0	1	1	2	3	5	8	13	21	34	55

Recurrence Reln:

$$T(n) = \begin{cases} 0, & \text{if } n=0 \\ 1, & \text{if } n=1 \\ T(n-1)+T(n-2), & \text{otherwise } (n>1) \end{cases}$$

T(n) :- Fib(n)

```
fib(n)
{ if (n==0)
      return 0;
  if (n==1)
      return 1;
  else
      return (fib(n-1)+fib(n-2));
}
```

// or

```
a = fib(n-1);
b = fib(n-2);
c = a+b;
return c;
```

```
fib(6)
{
  a=fib(5);
  b=fib(4);
  return(a+b);
}
```

```
fib(5)
{
  a=fib(4);
  b=fib(3);
  return(a+b);
}
```

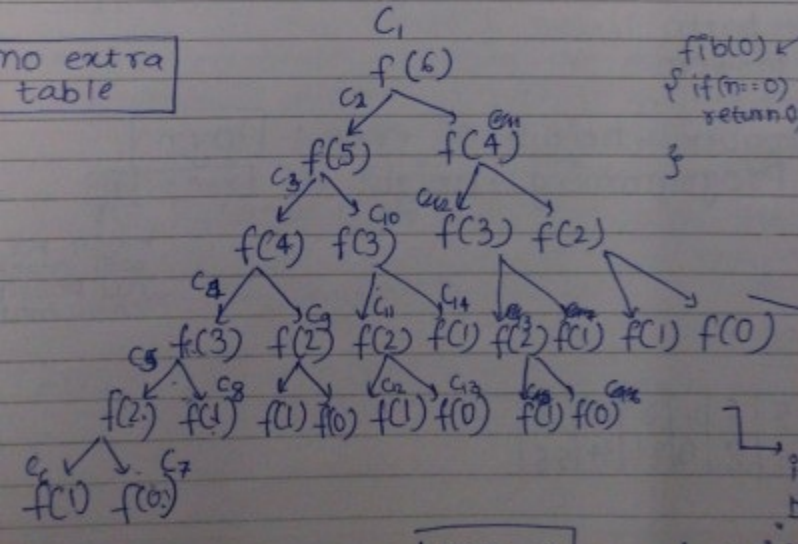
```
fib(4)
{
  a=fib(3);
  b=fib(2);
  return(a+b);
}
```

```
fib(2)
{
  a=fib(1);
  b=fib(0);
  return(a+b);
}
```

```
fib(0)
{ if(n==0)
    return 0;
}
```

```
fib(1)
{ if (n==1)
    return 1;
}
```

no extra table

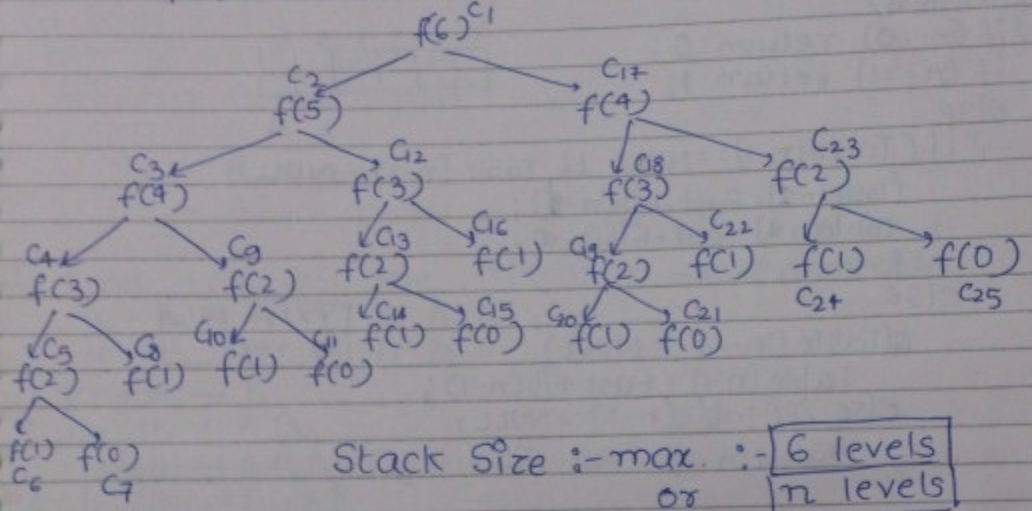


→ n level binary tree.

↳ if we assume that it is complete complete binary tree,

Cosmos

f(6) → f(5) → f(4) → f(3) → f(2) → f(1)



Stack Size :- max. :- 6 levels
or n levels

∴ in fibonacci series fn. :- each fn. requires 6 Bytes.
∴ Space complexity = max. stack size × 6 Bytes + 2 (i/p size)
= n × 6 Bytes + 2
= $\boxed{\Theta(n)}$

Drawback Of Recursive Soln. :-
- Note :- In the above recursive tree many subproblems are repeating (Overlapping subproblems).
- In dynamic programming, we will solve every subproblem only once.
- Dynamic Programming time complexity for fibonacci series $\boxed{O(n)}$ because it will call only distinct fn calls. As soon as a particular fn. is completed, that value is stored in a table so that no need to complete the function again.

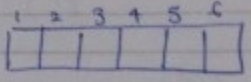
Fibonacci Series using Dynamic Programming :-
Time Complexity :- $\boxed{O(n)}$
Space Complexity :- 2 + 6n + n = $\boxed{\Theta(n)}$

f(3) = f(1) + f(2)

Cosmos

Fast Fibonacci Series :-

```
Fast-Fib(n)
{if (n==0) return 0;
 if (n==1) return 1;
 else
 { if (Table[n-1]==NULL && Table[n-2]==NULL)
      {Table[n-1] = Fast-Fib(n-1);
       Table[n-2] = Fast-Fib(n-2);
      }
   else
      {if (Table[n-1]==NULL)
          Table[n-1] = Fast-Fib(n-1);
       else if (Table[n-2]==NULL)
          Table[n-2] = Fast-Fib(n-2);
      }
   Table[n] = Table[n-1] + Table[n-2];
   return(Table[n]);
 }
}
```



extra table is req.

Max. level of the tree :- n

LCS :-

Subsequence :- of a given sequence is just the given sequence only in which 0 or (more) symbols are left out.

1 2 3 4 5 6

e.g. S = (A, B, B, A, B, B) → sequence

S1 = (A, A, B, B) → subsequence
(1, 4, 5, 6)

S2 = (A, A) → subsequence
(1, 4)

S3 = (B, B, B, B) → "
(2, 3, 5, 6)

S4 = (B, B, A, A) → not the subsequence

S5 = (A, B, B, B, A) → " "

Common Subsequence:-

z is a common subsequence of x & y if z is subsequence to both x & y.

e.g. x = A B B A B B
y = B A A B A A
CS_1 = A B A
CS_2 = B A B
CS_3 = A B

Q. Find LCS for the following subsequences:-
x = A B C B D A B
y = B D C A B A

1 length → A
2 length → A B
3 length → A B A
4 length → B D A B
5 length

Q. x: B B B A A A
y: B B A A A A

(decrementing 2nd sequence when it doesn't matches.)

LCS(6,6) = 1 + LCS(5,5)
= 1 + LCS(4,4)
= 1 + LCS(3,3)
= 0 + LCS(3,2)
= 1 + LCS(2,1)
= 1 + LCS(1,0)
= 0

Q. x: B B A A B A
y: B B B A B A

(decrementing 1st sequence when it doesn't matches.)

LCS(6,6) = 1 + LCS(5,5)
= 1 + LCS(4,4)
= 1 + LCS(3,3)
= 0 + LCS(2,3)
= 1 + LCS(1,2)

Cosmos

* In such case, when we have the choice of decrementing either of the sequence,
  - greedy says decrement either of them, but not both (this sometimes give correct result & sometimes incorrect.
  - dynamic says do both options, this always give correct result.

LCS(i,j) = Length of longest common subsequence possible with the 2 sequences x & y, where sequence x contains 'i' symbols & y contain 'j' symbols

```
LCS(i,j)
{ if (i==0 || j==0)
    { printf("No common subsequence");
    }
  else if (i
```

$$LCS = \begin{cases} 0 ; \text{if } i=0 \text{ or } j=0 \\ 1+LCS(i-1,j-1) \quad \text{if } x[i]=y[j] \\ \max(LCS(i,j-1), LCS(i-1,j)) \quad , \text{otherwise} \end{cases}$$

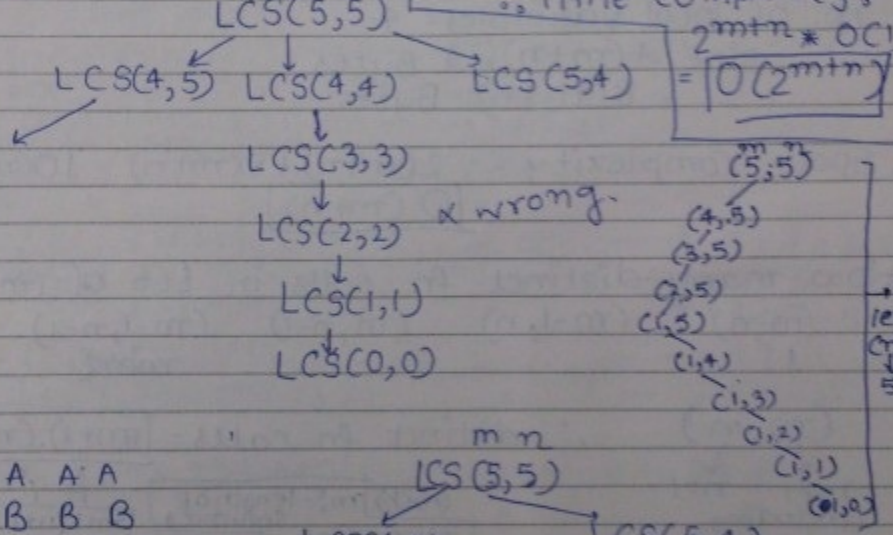
```
LCS (i,j)
  { if (i==0 || j==0)
      {printf(" 0 common sequence");
       return 0;
      }
    else if (x[i]==y[j])
      { a= LCS(i-1,j-1);
        return (a+1);
      }
    else { if((LCS(i,j-1) > LCS(i-1,j))
```

Cosmos

If we dont use dynamic programming then many fn. calls are repeated, at max. there are (m+n) levels, so for complete binary tree, the max. no. of nodes are 2^{m+n}, let's suppose each fn call takes O(1) time. ∴ complexity [illegible]

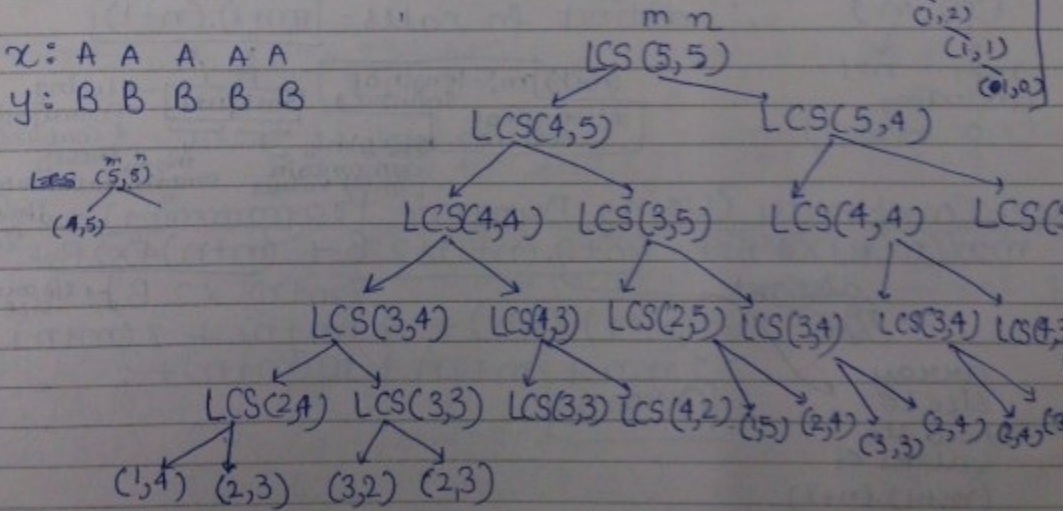
```
a = LCS(i-1, j);
b = LCS(i, j-1);
if (a > b)
    return a;
else
    return b;
}
}
```

★★ In the recursive tree
If we draw recursive tree for LCS of (m,n), in the worst case we will get (m+n) level complete binary tree. So total no. of nodes in (m+n) level complete binary tree is $2^{m+n}-1$
∴ $2^{m+n}-1$ no. of function calls.
We assume each fn. call takes O(1) time,
∴ Time complexity :-
$2^{m+n} * O(1) = \boxed{O(2^{m+n})}$



✗ wrong.

x: A A A A A
y: B B B B B



★ The height of the tree (at max.) = $\boxed{m+n}$

Cosmos

In above recursive tree many fn. calls are repeated, so we have to go to dynamic programming, which will solve every subproblem only once.

Space complexity:-

2 i/p arrays are req. of size m & n.

∴ i/p size = (m+n) × 2 Bytes (each integer/char of 2 B).

Stack size (max.) = m+n

no. of local variables = 4

∴ size = 4(m+n) × 2 Bytes

= 8(m+n) Bytes.

∴ Space complexity = 2(m+n) + 8(m+n) = 10(m+n)

= O(m+n)

How many distinct fn. calls in LCS of (m,n).

(m,n) (m-1,n) (m,n-1) (m-1,n-1)

1

(m,n)

↓ ↘

m+1 n+1

(including 0 as well)

∴ distinct fn. calls = (m+1).(n+1)

m:- length of sequence 1

n:- length of sequence 2.

m+1 n+1 → using permutations & combinations theory

this can contain (m+1) values

this can contain (n+1) values

Space Complexity (Using Dynamic Programming):- distinct fn. calls

max(m,n) × 4 B + (m+1).(n+1) × 2 B + (m+n) 4 × 2 B + (m+n) × 2 B → i/p array size

= 2(m+1).(n+1) + 8(m+n) + 2(m+n)

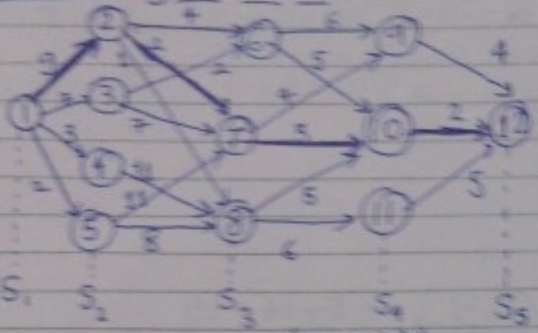
= 2mn + 2m + 2n + 10(m+n) + 2

= O(mn)

distinct fn. calls → array size to hold the values of (m+1).(n+1) size

Cosmos

Multistage Graph



→ edge weight

using Greedy :-

1 - 5 - 8 - 10 - 12 ; - 17

MSG(1,1) — from 1st stage, from 1st vertex

MSG(3,8) — from 3rd stage, from 8th vertex

→ to final dest^{n}. which is vertex 12

$$MSG(1,1) = \begin{cases} C(1,2) + MSG(2,2) \\ C(1,3) + MSG(2,3) \\ C(1,4) + MSG(2,4) \\ C(1,5) + MSG(2,5) \end{cases} \begin{matrix} \longrightarrow 4+7=11 \\ \longrightarrow 2+5=7 \checkmark \\ \longrightarrow 3+18=21 \\ \longrightarrow 2+15=17 \end{matrix}$$

min of the 4.

$$MSG(2,2) = \begin{cases} C(2,6) + MSG(3,6) \\ C(2,7) + MSG(3,7) \\ C(2,8) + MSG(3,8) \end{cases} \begin{matrix} \longrightarrow 4+7=11 \\ \longrightarrow 2+5=7 \checkmark \\ \longrightarrow 1+7=8 \end{matrix}$$

min. of 3

$$MSG(2,3) = \begin{cases} C(3,6) + MSG(3,6) \\ C(3,7) + MSG(3,7) \end{cases} \begin{matrix} \longrightarrow 2+7=9 \checkmark \\ \longrightarrow 7+5=12 \end{matrix}$$

min. of 2

$$MSG(2,4) = C(4,8) + MSG(3,8) \longrightarrow 11+7=18$$

$$MSG(2,5) = \begin{cases} C(5,7) + MSG(3,7) \\ C(5,8) + MSG(3,8) \end{cases} \begin{matrix} \longrightarrow 11+5=16 \\ \longrightarrow 8+7=15 \checkmark \end{matrix}$$

min. of 2

$$MSG(3,6) = \begin{cases} C(6,9) + MSG(4,9) \\ C(6,10) + MSG(4,10) \end{cases} \begin{matrix} \longrightarrow 6+4=10 \\ \longrightarrow 5+2=7 \checkmark \end{matrix}$$

min. 2

$$MSG(4,9) = C(9,12) = 4$$

$$MSG(4,10) = C(10,12) = 2$$

$$MSG(3,7) = \begin{cases} C(7,9) + MSG(4,9) \\ C(7,10) + MSG(4,10) \end{cases} \begin{matrix} \longrightarrow 4+4=8 \\ \longrightarrow 3+2=5 \checkmark \end{matrix}$$

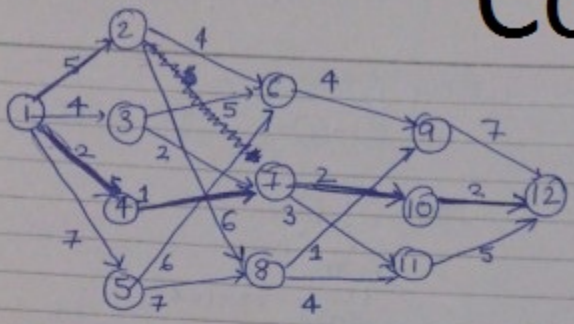
min.

$$MSG(3,8) = \begin{cases} C(8,10) + MSG(4,10) \\ C(8,11) + MSG(4,11) \end{cases} \begin{matrix} \longrightarrow 5+2=7 \checkmark \\ \longrightarrow 6+5=11 \end{matrix}$$

min.

$$MSG(4,11) = C(11,12) = 5$$

Cosmos



$$MSG(1,1) = \begin{cases} \min \begin{pmatrix} c(1,2)+MSG(2,2) \\ c(1,3)+MSG(2,3) \\ c(1,4)+MSG(2,4) \\ c(1,5)+MSG(2,5) \end{pmatrix} \end{cases} \begin{matrix} \rightarrow 5+14=19 \\ \rightarrow 4+6=10 \\ \rightarrow 2+5=7 \checkmark \\ \rightarrow 7+15=22 \end{matrix}$$

$$MSG(2,2) = \begin{cases} \min \begin{pmatrix} c(2,6)+MSG(3,6) \\ c(2,8)+MSG(3,8) \end{pmatrix} \end{cases} \begin{matrix} \rightarrow 4+11=15 \\ \rightarrow 6+8=14 \checkmark \end{matrix}$$

$MSG(3,6) = \{ c(6,9) + MSG(4,9) \rightarrow 4+7=11$

$MSG(4,9) = c(9,12) = 7$

$$MSG(3,8) = \begin{cases} \min \begin{pmatrix} c(8,9)+MSG(4,9) \\ c(8,11)+MSG(4,11) \end{pmatrix} \end{cases} \begin{matrix} \rightarrow 1+7=8 \checkmark \\ \rightarrow 4+5=9 \end{matrix}$$

$MSG(4,11) = c(11,12) = 5$

$$MSG(2,3) = \begin{cases} \min \begin{pmatrix} c(3,6)+MSG(3,6) \\ c(3,7)+MSG(3,7) \end{pmatrix} \end{cases} \begin{matrix} \rightarrow 5+11=16 \\ \rightarrow 2+4=6 \checkmark \end{matrix}$$

$$MSG(3,7) = \begin{cases} \min \begin{pmatrix} c(7,10)+MSG(4,10) \\ c(7,11)+MSG(4,11) \end{pmatrix} \end{cases} \begin{matrix} \rightarrow 2+2=4 \checkmark \\ \rightarrow 3+5=8 \end{matrix}$$

$MSG(4,10) = c(10,12) = 2$

$MSG(2,4) = c(4,7) + MSG(3,7) = 1+4=5$

$$MSG(2,5) = \text{etc} \begin{cases} \min \begin{pmatrix} c(5,6)+MSG(3,6) \\ c(5,8)+MSG(3,8) \end{pmatrix} \end{cases} \begin{matrix} \rightarrow 6+11=17 \\ \rightarrow 7+8=15 \checkmark \end{matrix}$$

nd Method :-

from ⑥ :- 11 (6-9-12) ✓

from ⑦ :- 4 (7-10-12) ✓
8 (7-11-12)

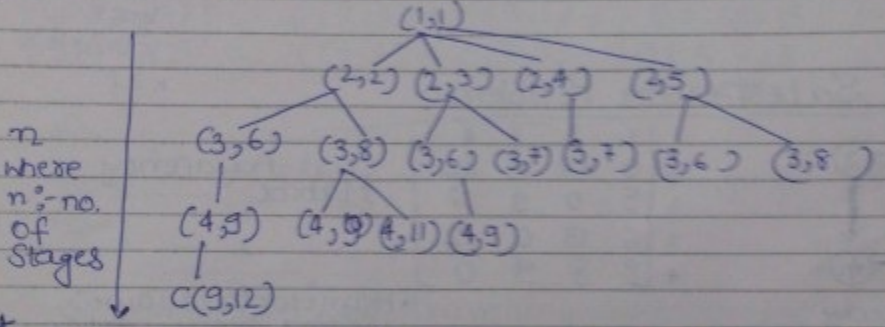
from ⑧ :- 8 (8-9-12) ✓
9 (8-11-12)

om ② :- 15 (2-6-min. distance from ⑥)
6+8=14 (2-8-min. distance from ⑧) ✓

☆ Dynamic programming & brute force becomes same when all problems are distinct.

from ③ :- 5+11=16 (3-6-min dis from ⑥)
2+4=6 (3-7-min " " ⑦) ✓
from ④ :- 6+4=10 (4-7-" " " ⑦)

Cosmos



☆ Multistage graph will generate in the worst case 'n' level complete binary tree, n:- no of stages, so total time complexity is $O(2^n)$ → [no. of nodes in complete binary tree $= 2^n - 1$]
But in multistage graph problem many subproblems are repeated, so we perform dynamic programming

☆ MSG (G(V,E), N) ⟶ V distinct function call

☆ All function calls combined take O(E) times.
∴ $O(V+E)$

Space Complexity :- (V+E) + n + V
(V+E): i/p in form of adjacency list
n: stack-size (no. of stages)
V: table size (distinct fn. calls)
= V+E+V, + V = 3V+E
max stack size
= $O(V+E)$

Cosmos

$MSG(S_i, V_j)$ = the min. cost req. from stage S_i & vertex V_j to destination.

$$MSG(S_i,V_j) = \begin{cases} (C(V_j,k) + MSG(S_{i+1},k)) \rightarrow min. \\ \forall k \in S_{i+1} \wedge (V_j,k) \in E \\ \quad \downarrow \text{edge should be these.} \\ C(V_j,D) \quad \text{if } S_i == (F-1) \end{cases}$$

D → Destⁿ

F → Final stage

$c(V_j,k)$

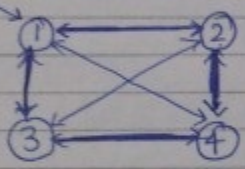
$V_j \in F$

$c(V_j,k) \in E$

$k \in F$

Date
22.07.12

Travelling Salesperson Problem:-



	1	2	3	4
1	0	10	15	20
2	5	0	9	10
3	6	13	0	12
4	8	8	9	0

Cost-Adjacency Matrix

Acc. to greedy:-

1 – 2 – 3 – 4 – 1

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- Hamiltonian Tour:- Visiting every vertex exactly once.
- Euler Tour:- Visiting every edge exactly once.

Using Dynamic Programming:-

$$TSP\{1,\{2,3,4\}\} = \min \begin{cases} c(1,2) + TSP\{2,\{3,4\}\} \rightarrow 10+25=35 \checkmark \\ c(1,3) + TSP\{3,\{2,4\}\} \rightarrow 15+25=40 \\ c(1,4) + TSP\{4,\{2,3\}\} \rightarrow 20+23=43 \end{cases}$$

(1: starting from 1; can go to either 2 or 3 or 4)

$$TSP\{2,\{3,4\}\} = \min \begin{cases} c(2,3) + TSP\{3,\{4\}\} \rightarrow 9+20=29 \\ c(2,4) + TSP\{4,\{3\}\} \rightarrow 10+15=25 \checkmark \end{cases}$$

$$TSP\{3,\{4\}\} = c(3,4) + c(4,1) = 12+8=20$$

($c(4,1)$: go back to 1 from 4)

$$TSP\{4,\{3\}\} = c(4,3) + c(3,1) = 9+6=15$$

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$TSP\{3,\{2,4\}\} = \min \begin{cases} c(3,2) + TSP\{2,\{4\}\} \rightarrow 13+18=31 \\ c(3,4) + TSP\{4,\{2\}\} \rightarrow 12+13=25 \checkmark \end{cases}$

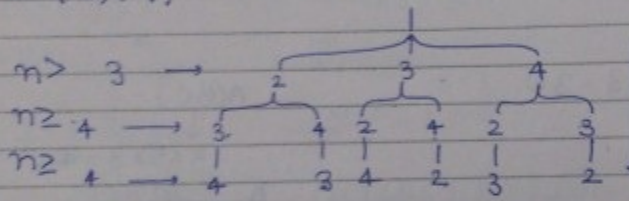
$TSP\{2,\{4\}\} = c(2,4) + c(4,1) = 10+8=18$

$TSP\{4,\{2\}\} = c(4,2) + c(2,1) = 8+5=13$

$TSP\{4,\{2,3\}\} = \min \begin{cases} c(4,2) + TSP\{2,\{3\}\} \rightarrow 8+15=23 \\ c(4,3) + TSP\{3,\{2\}\} \rightarrow 9+18=27 \end{cases}$

$TSP\{2,\{3\}\} = c(2,3) + c(3,1) = 9+6=15$

$TSP\{3,\{2\}\} = c(3,2) + c(2,1) = 13+5=18$



at each level we have (at max.) n nod & there are n levels. n+n+.....+n(n ti

- Time complexity of Travelling Salesperson w/o dynamic programming :- $\Omega(2^n)$ [Brute force method]
- Using Dynamic Programming :- $\Omega(2^n)$ [No overlapping subproblem.]

number of nodes in the above tree

$n \times (n-1) \times (n-2) \times \ldots \times 1 \leq n \times n \times \ldots \times n$ (n times)

$= n^n$

$= O(n^n)$

★ Travelling Salesperson is one of the NP complete proble because they can't be solved in polynomial time.

Space Complexity = V^2 (adjacency matrix) + $V+E$ (graph) + V (stack size) + 2^V (distinct function calls)

$= O(2^V)$ $\Omega(2^V)$

TSP(A,R):- min. cost required to go from vertex A to all other remaining vertices(R) exactly once & coming back to source S

Recurrence Reln.:-

$$TSP(A,R) = \min_{K \in R} \{ C(A,K) + TSP(K,R') \} \quad \text{where } \forall K \in R \text{ \& } R' = R - \{K\}$$

$$= C(A,S) \quad \text{if } R = \emptyset.$$

Source is S.

Matrix Chain Multiplication

e.g. $A_{2\times10}$ $B_{10\times5}$

$C = AB$

2×5

∴ Total no. of multiplications $= 2\times5\times10 = 100$

$C(k,S) + C(A,k)$

Q. $A_{2\times3}$ $B_{3\times5}$ $C_{5\times3}$

$(AB)C \rightleftharpoons (AB) \rightarrow 2\times5\times3 = 30$

$(AB)_{2\times5}\ C_{5\times3}$

$2\times3\times5 = 30$

$(AB)C \rightarrow \boxed{60}$ multiplications

$A(BC)$

$3\times5\times3 = 45$

$A_{2\times3}\ (BC)_{3\times3}$

$2\times3\times3 = 18$

$A(BC) \rightarrow \boxed{63}$ multiplications

No. of multiplications of $m\times p$ & $p\times n$ matrices = $\boxed{m\times n\times p}$

ABCDE

(A)(BCDE) → this will have 5 because BCDE can be multiplied in 5 ways

(AB)(CDE) → 2

(ABC)(DE) → 2

(ABCD)(E) → 5

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$A_{1\times5}\ B_{5\times2}\ C_{2\times1}$

$(AB)_{1\times2}\ C_{2\times1} \rightarrow 1\times1\times2 = 2$

$1\times2\times5 = 10$

$10 + 2 = \boxed{12}$ ✓

$A(BC)_{5\times1}$

$5\times2\times1 = 10$

$A_{1\times5}\ (BC)_{5\times1}$

$1\times1\times5 = 5$

Cosmos

Q. $A_{2\times5}\ B_{5\times4}\ C_{4\times2}\ D_{2\times3}$

(A)(BCD)

(A)(BC)(D)

$5\times4\times2=40$

$A_{2\times5}(BC)_{5\times2}\ D_{2\times3}$

((A)(BC))D

$2\times5\times2=20$

$(ABC)_{2\times2}\ D_{2\times3}$

$2\times2\times3=12$

((A)(BC))(D) → 72

(A)((B)(CD))

$4\times2\times3=24$

$5\times4\times3=60$

$2\times5\times3=30$

(A)((BC)(D))

$5\times2\times3=30$

$(A)(BCD)_{5\times3}$

$2\times5\times3=30$

(A)((BC)(D)) → 100

$5\times4\times3$ → 114

(AB)(CD) → 88

$2\times5\times4=40$

24

$(AB)_{2\times4}(CD)_{4\times3}$

$2\times4\times3=24$

(ABC)(D) → 40+16+12 = 68

((AB)C)(D)

40

$(AB)_{2\times4}\ C_{4\times2}$

$2\times4\times2=16$

$(ABC)_{2\times2}\ D_{2\times3}$

$2\times2\times3=12$

Q. $A_{1\times5}\ B_{5\times1}\ C_{1\times2}\ D_{2\times5}$

n-levels [n:- no. of matrices]

ABCD

(A)(BCD)

(A)((BC)D)

$5\times1\times2=10$

$(BC)_{5\times1}\ D_{2\times5}$

$5\times2\times5=50$

(A)(B(CD)) → 60

$1\times2\times5=10$

$B_{5\times1}(CD)_{1\times5}$

$1\times5\times5=25$

(AB)(CD) → 20

$1\times5\times1=5$

10

$(AB)_{1\times1}(CD)_{1\times5}$

$1\times5=5$

(ABC)(D) 17

((AB)C)(D)

5

$(AB)_{1\times1}\ C_{1\times2}$

$1\times2=2$

$(ABC)_{1\times2}(D)_{2\times5}$

(ABC)D 30

$A_{1\times5}(BC)_{5\times1}$

10

$(ABC)_{1\times1}$

each comparison takes O(1) time, as there are (n-1) comparisons, ∴ O(n) time

• MCM(1,4):- min. no. of multiplications req. to multiply 1 to 4 matrices.

cost

$$MCM(1,4) = \min \begin{cases} MCM(1,1) + MCM(2,4) + 25 & \text{combined multiplication of } (1,1) \& (2,4) \text{ e.g. } (A) \& (BCD) \\ MCM(1,2) + MCM(3,4) + 5 & \text{combined multiplication of } (1,2) \& (3,4) \text{ e.g. } (AB) \& (CD) \\ MCM(1,3) + MCM(4,4) + 10 & \end{cases}$$

(n-1) comparisons ∴ O(n)

$$MCM(2,4) = \min \begin{cases} MCM(2,2) + MCM(3,4) + 25 & \text{combined multiplication of } (2,2) \& (3,4) \text{ e.g. } (B) \& (C,D) \\ MCM(2,3) + MCM(4,4) + 50 & \end{cases}$$

(n-2) comparisons ∴ O(n)

$$MCM(3,4) = \min \{ MCM(3,3) + MCM(4,4) + 0 \}$$ combined multiplication cost of (3,3) & (4,4) e.g. (C) & (D)

$$MCM(2,3) = \min \{ MCM(2,2) + MCM(3,3) + 10 \}$$

$$MCM(1,2) = \min \{ MCM(1,1) + MCM(2,2) + 5 \}$$

$$MCM(1,3) = \min \begin{cases} MCM(1,1) + MCM(2,3) + 10 & \text{combined multiplication cost of } (1,1) \& (2,3) \text{ e.g. } (A) \& (BC) \\ MCM(1,2) + MCM(3,3) + 2 & \end{cases}$$

(n-2) comparisons ∴ O(n)

* Without Dynamic Programming n-level binary tree will generate 2^n

In the above tree, many subproblems are repeated, so we will perform Dynamic Programming.

MCM(1,n) contain how many distinct subproblem?

Starting with 1: (1,4) (1,3) (1,2) (1,1)

Starting with 2: (2,2) (2,3) (2,4)

Starting with 3: (3,3) (3,4)

Starting with 4: (4,4)

∴ MCM(1,4) contains 10 distinct fn calls.

	1	2	3	4
1	✓	✓	✓	✓
2	×	✓	✓	✓
3	×	×	✓	✓
4	×	×	×	✓

distinct.

For MCM(1,n) contains $\frac{n^2}{2}$ function calls.

half of the n^2 matrix

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Time Complexity : $\frac{n^2}{2}$ x time complexity of each fn. call

$= \frac{n^2}{2} \times O(n)$

$= \boxed{O(n^3)}$

Space Complexity :- $4(n) + n + \frac{n^2}{2}$

- $4(n)$: to store size of the matrix (4 Bytes for each matrix); no. of matrices
- n: Stack size
- $\frac{n^2}{2}$: table-size (no. of distinct fn. calls)

$= O(n^2)$

Recurrence Relⁿ :-

$MCM(i,j) =$ { min. no. of multiplications req. to multiply i to j matrices.

$MCM(i,j) = \begin{cases} \min(MCM(i,k) + MCM(k+1,j)) \ \forall\ k \text{ from } i \text{ to } (j-1) \\ 0 \quad \text{if } i == j. \end{cases}$

↑ my answer

Correct answer $= \begin{cases} \min\left[MCM(i,k) + MCM(k+1,j) + \text{no. of multiplications req. to multiply the matrices } (i,k) \text{ and } (k+1,j)\right] \\ i \le k < j \\ 0, \text{ if } i == j \end{cases}$

0/1 Knapsack Problem

objects	ob_1	ob_2	ob_3	$n=3$
profits	5	3	4	$M=5$
weights	3	2	1	

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Manually:-

case 1:- ~~ob_1~~ & ob_3 ∴ profit = 9
case 2:- ob_1 & ob_2 ∴ profit = 8
case 3:- ob_2 & ob_3 ∴ profit = 7

Using Greedy:-

$ob_1 \rightarrow \frac{5}{3} = 1.66$
$ob_2 \rightarrow \frac{3}{2} = 1.5$
$ob_3 \rightarrow \frac{4}{1} = 4$

take ob_3 first, capacity left = 4
take ob_1 2nd capacity left = 1.
but we can't take fraction, ∴ we can't take fraction of ob_2.

Q n = 3 M = 15

Objects	ob_1	ob_2	ob_3	ob_4	ob_5
profits	50	29	33		
weight	10	6	7		
profit:weight	5	4.83	4.7		

Using Greedy:-
take ob_1 1st, ∴ capacity left = 5
we can't take any more.
∴ profit = 50

Manually:-
ob_2 & ob_3 → profit = 62
ob_1 → " = 50

* Note:- Greedy algo will not give optimal soln. for 0/1 Knapsack, so we use dynamic programming which will cover all possible combination.

n = 5 M = 10

bjects	ob_1	ob_2	ob_3	ob_4	ob_5
rofits	7	2	1	6	12
eights	3	1	2	4	6

e.g. 01KP(5,10) → 5 objects to choose from (from 1 to 5)
→ capacity of knapsack

01KP(4,10)
4 objects to choose from (1 to 4) excluded 5.

1KP(n,m) = max. profit we will get in 0/1 Knapsack proble
here n :- no. of objects & m is the capacity of knapsack.
↓ among which we can select
↓ current capacity.

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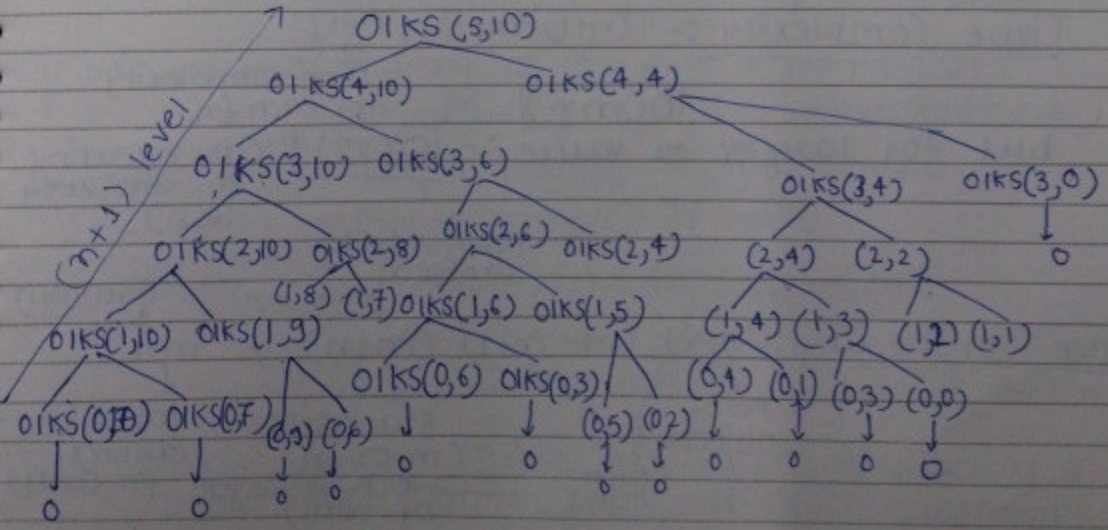
$$01KS(n,m) = \begin{cases} 0 & \text{if } (n) \text{ or } (m) == 0 \\ \max\{ 01KP(n-i, m-wi) + bi \ \forall i \\ \quad 01KP(n-i, m) \\ 01KS(n-1, m) & \text{if } w_n > M \\ \max.\{ 01KS(n-1, m-w_n) + P_n & \text{if } w_n \le M \\ \quad 01KS(n-1, m) \end{cases}$$

```
01KS(n,m){
   if (n==0 or m==0)
     {return 0;
     }
   else if (Wn>M)
     {01KS(n-1,m);
     }
   else
     {a = 01KS(n-1,m-Wn)+Pn ;
      b = 01KS(n-1,m);
      return (max (a,b));
     }
}
```



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0/1 Knapsack of (n,m) will generate n-level (approx.) almost complete binary trees.
∴ 2^n function calls. & every fn. call takes O(1) time [as we have to make only 1 comparison in max. fn.]
∴ W/O using Dynamic :- $O(2^n)$

☆ In above rec. tree some fn. calls are repeating, so we perform dynamic programming, which calls distinct fn. calls exactly once.

Q 01 Knapsack contains how many distinct fn. calls
01KS(5,10)
- 5 can decrease in 6 ways (5,4,3,2,1,0)
- the weight can decrease acc. to the weights of the object, but maximum it can have values from 10 to 0, ∴ 11

∴ 6 × 11 = 66 distinct functions (max.)
∴ 01KS(n, m) ∴ distinct fn. calls = (n+1)(m+1)
(n+1) (m+1)

∴ Time Complexity :- (n+1)(m+1) × O(1)
(O(1) : time complexity of each fn)
:- $O(mn)$
but for larger n value :- $O(2^n)$ (as repeating fn. are very less.)
($O(2^n)$: NP Complete problem)

Space Complexity :- n × k + (max.(m,n)) (m+n) + (m+1)(n+1) mn
- n × k : n = no. of objects, k = size of each object
- (m+n) : stack size (m or n which is larger of both)
- mn : table size (distinct fn. calls)

Cosmos

Sum Of Subsets Problem :-

$S = \{70, 20, 90, 50, 25, 15, 40, 30\}$ (indices 1 to 8)

$M = 200$

$S_1 = (90, 90, 50, 40)$

$S_2 = (70, 90, 40)$

$S_3 = (70, 90, 25, 15)$

$S_4 = (70, 20, 25, 15, 40, 30)$

Q. To find a subset whose sum is 200.

SOS(8, 200)

- 8 : Set contains 8 elements
- 200 : the subset sum is 200 (or we can say that the remaining value to make it 200).

e.g. • SOS(8, 200) means set is over but no element is choosen.

• SOS(5, 0) means 5 elements are left but we get the subset whose sum is 200.

• SOS(4, 100) means 4 elements are left & we need 100 more to make the subset sum to be 200.

$$SOS(n,m) = \begin{cases} SOS(n-1,m) & \text{if } m > W_n \\ & \text{if } m < W_n \text{ but } m \neq 0 \\ \begin{pmatrix} SOS(n-1, m-W_n), S = S \cup W_n \\ SOS(n-1, m) \end{pmatrix} & \\ S \text{ (return subset)} & \text{if } m = 0. \end{cases}$$

- initial value of $S = \phi$
- initial value of n :- no. of elements in the subset.
- " " " m :- M (e.g. M = 200 in above example.)

$$SOS(n,m) = \begin{cases} -1 \text{ (error message)} & \text{if } n == 0 \text{ \& } m \neq 0 \\ S \text{ (return subset)} & \text{if } m == 0 \\ SOS(n-1, m) & \text{if } w_n > m \\ \begin{pmatrix} SOS(n-1, m-w_n), S = S \cup W_n \\ SOS(n-1, m) \end{pmatrix} & \text{if } W_n \leq \end{cases}$$

SOS(n, m) :- Sum of subsets of (n, m) is finding a subset from given set of 'n' elements whose sum is 'm'.

Cosmos

★ The rec. rel^n makes n-level binary tree which have 2^n nodes.

∴ Time Complexity :- $2^n \times O(1)$
↳ time complexity of each fn.

$:- O(2^n)$.

★ Using Dynamic Programming:-

SOS (n, m)

this can decrease in (n+1) ways
e.g. n=8
∴ it can be (8,7,6,...,1,0)

this can decrease into (m+1) ways [maximum]
[e.g. in above example m=200, ∴ it can be (200,199,198,....,2,1,0) but many of these values are not possible, but m can take max. (m+1) values.

∴ Time Complexity :- $O(mn)$
When n is large :- $O(2^n)$

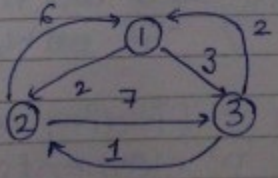
0/1 Knapsack problem can be polynomially reducible to Sum of Subsets problem, so both are NP Complete.

Space Complexity :- i/p size + Stack size (n-level binary tree) + distinct fn. calls.

$(n+1) + n + (m+1)(n+1)$

$= O(mn)$

All pairs Shortest Path



① +ve edge weights
|V|* Dijkstra algo ⇒ V.(V+E)logV

② -ve edge weight & -ve edge weight cycle :-
|V|* Bellman Ford
|V|* VE

Cosmos

$A^k(i,j)$ = min. cost required to go from vertex i to vertex j with the intermediate vertices $k \in \{0,1,2,3,\ldots,k\}$

A^0	1	2	3
1	0	2	3
2	6	0	7
3	2	1	0

A^1	1	2	3
1	0	2	3
2	6	0	7
3	2	1	0

A^2	1	2	3
1	0	2	3
2	6	0	7
3	2	1	0

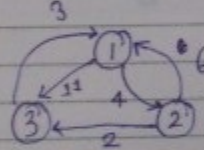
A^3	1	2	3
1	0	2	3
2	6	0	7
3	2	1	0

$A^1(2,3) = \begin{cases} A^0(2,3) \\ A^0(2,1) + A^0(1,3) \end{cases}$ (via 1)

this means path from 2 to 3 via 1

$A^2(2,3)$ → this means path from 2 to 3 via 1 & 2.

Q.



A^0	1	2	3
1	0	4	11
2	6	0	2
3	3	∞	0

A^1	1	2	3
1	0	4	11
2	6	0	2
3	3	7	0

A^2	1	2	3
1	0	4	6
2	6	0	2
3	3	7	0

A^3	1	2	3
1	0	4	6
2	5	0	2
3	3	7	0

- $A^2(2,3) = \min\{A^1(2,3), A^1(2,2) + A^1(2,3)\}$

via 0,1,2 from A^2. / from $A^2(2,3)$

- $A^3(1,2) = \min\{A^2(1,2), A^2(1,3) + A^2(3,2)\}$

via 0,1,2,3 from A^3. / from $A^2(1,2)$

[via 0 means direct path

- $A^0(i,j) \rightarrow V^2$ function calls
- $A^1(i,j) \rightarrow V^2$ function calls
- $A^2(i,j) \rightarrow$ " " "
- $A^3(i,j) \rightarrow$ " " "

all these are distinct function calls.

$\therefore V^2 \times V$ → V no. of times

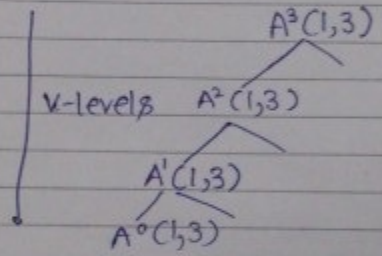
distinct fn

Cosmos

```
Allpairsshortestpath(A°,n)
{ for(k=1 to n)
    { for(i=1 to n)
        { for(j=1 to n)
            {
              A^k(i,j) = min.{A^(k-1)(i,j), A^(k-1)(i,k) + A^(k-1)(k,j)};
            }
        }
    }
```

Time Complexity:-

$O(V^3)$



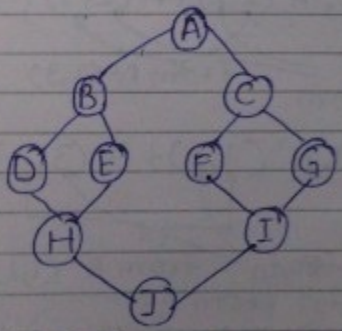
Space Complexity:-

$V^2 + V + V^2 = O(V^2)$

- V^2 → i/p matrix A°
- V → Stack size (V-levels)
- V^2 → table size (next table like A° & A^1, A^1 & A^2, A^2 & A^3)

Graph Traversal

1. Breadth First Traversal (BFT)
2. Depth First Traversal (DFT)



BFT

A - B - C - D - E - F - G - H - I - J

DFT

A - B - D - H - I - E - C - F - G - J

```
BFT(v)
{ visited(v) = 1;
  add(v,Q);
  while(Q is not empty)
```

→ while loop is executing V times

```
{x = delete(Q); printf(x);
 for all w adj x
  {if (w not visited)
    {visited(w) = 1;
     add(w, Q);
    }
  }
}
```

→ at max. each vertex is adjacent to each other vertex.

Note1:- In order to implement BFT, we are using queue as the data structure.

Time Complexity :- $V(1+V)$

- V → 'V' times while loop
- 1 → each vertex is connected to only one vertex
- V → each vertex is connected to each other vertex

$= V + V^2$

$= V + E \rightarrow \boxed{O(V+E)}$ or $\boxed{\Theta(V+E)}$

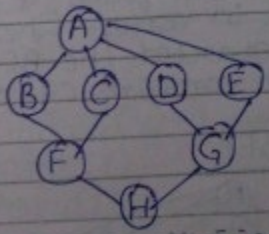
↑ whichever is greater.

- Best Case :- $\boxed{\Omega(V)}$
- Worst Case :- $\boxed{O(E)}$

★ BFT is also known as Level Order Traversal.

Q. Consider following graph :-

Give 4 diff. BFT's



① A-B-C-D-E-F-G-H

② A-E-B-C-D-G-F-H

③ A-C-B-D-E-F-G-H

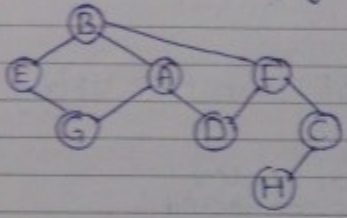
④ A-D-E-B-C-G-F-

Cosmos

Starting from C:-
C-A-F-B-D-E-H-G

Starting from H:-
H-F-G-B-C-D-E-A

Q. Consider the following Graph:-

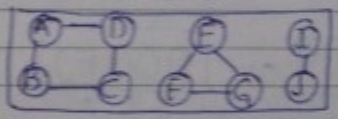


Give 4 diff. BFT's starting from H.

1. H-C-F-D-B-A-E-G
2. H-C-F-B-D-A-E-G
3. H-C-F-B-D-E-A-G
4. H-[illegible] only 3.

Applications of BFT:-

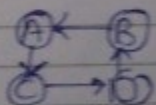
G(V,E)

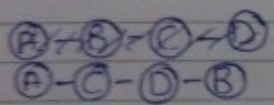


Starting from A:-
A-B-D-C I-J
E-F-G

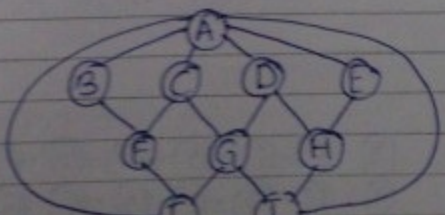
1. Using BFT, we can check whether the graph is connected or not.
2. To obtain the no. of connected components, we use BFT.

[connected component:- maximal subgraph which is connected.] → O(V+E)

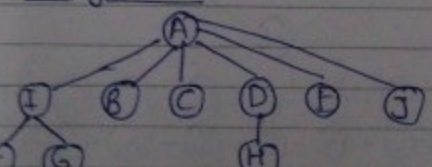




In a directed graph, if we come back to vertex where we have started the BFT, then graph contains cycle.



Using BFT

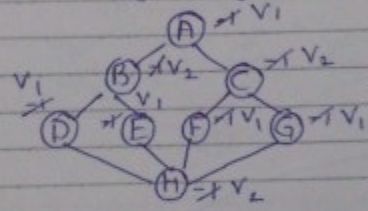


★★Any new problem is given in GATE about Graph, most of the times it is BFT & DFT.

In the given unweighted graph (all edge weights are same), to find out shortest path from the given source, we use BFT algorithm.

★★If given graph is unweighted, then
Dijkstra ⟶ O[(V+E)log V]
Bellman-Ford → O(VE)
BFT ⟶ O(V+E) ✓
BFT takes shortest time.

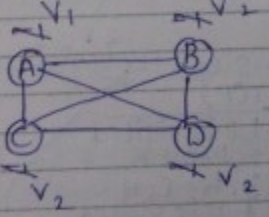
⑤ To check given graph is Bipartite or not, we use BFT algorithm



Step 1:- Visit a graph & mark '-1' to all vertices :- O(V+E)
Step 2:- Place parent in V_1 & children in V_2

V_1	V_2
A	B
D	C
E	H
F	
G	

→ Bipartite



V_1	V_2
A	B
	C
	D

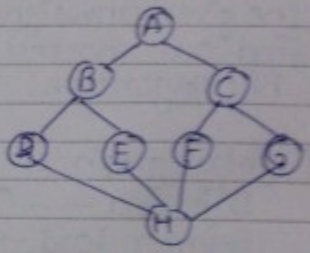
→ Not a bipartite. (as C & D are adjacent to B & are in V_2).

Depth First Traversal (DFT):-

★ DFT will take stack as a data structure.

```
DFT(V)
{1. visited(V) = 1;
 2. print(V);
 3. for all w adj to V
    {4. if(w is not visited)
        5.  DFT(w);
    }
```

O/p:- A B D H E F C G



DFT(A)
1. A is visited
2. Print A
3. W = B, C
4. B is not visited — C is now visited
5. DFT(B)
   1. B is visited
   2. Print B
   3. W = A, D, E
   4. A is visited — E is visited now
      D is not visited
   5. DFT(D)
      1. D is visited
      2. Print D
      3. W = B, H
      4. B is visited
         H is not visited
      5. DFT(H)
         1. H is visited
         2. Print H
         3. W = D, E, F, G
         4. D is visited — F is not visited — G is visited now
            E is not visited
         5. DFT(E)
            1. E is visited
            2. Print E
            3. W = B, H
            4. Both B & H visited
         5. DFT(F)
            1. F is visited
            2. Print F
            3. W = C, H
            4. H is visited
               C is not visited
            5. DFT(C)
               1. C is visited
               2. Print C
               3. W = A, F, G
               4. A & F is visited
                  G is not visited
               5. DFT(G)
                  1. G is visited
                  2. Print G
                  3. W = C, H
                  4. Both visited

★ No. of function calls = no. of vertices in = $\boxed{V}$.

★ In worst case, each function call 'for loop' is repeated $(V-1)$ times in case of complete graph.

★ In best case, in each fn. call 'for loop' is repeated only once.

Cosmos

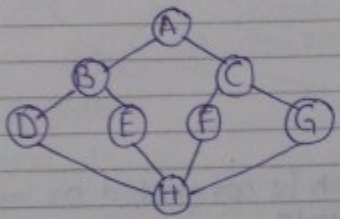
Note 1:- In order to implement DFT, we uses stack data structure.

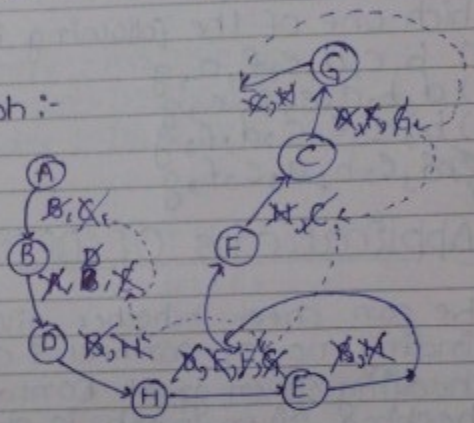
★Time complexity:- Best case

$O(V) \; V.(1+(V-1)) = \boxed{\theta(V+E)}$ as $\boxed{V^2 \geq E}$

V fn. calls; worst case

Q. Consider the following graph:-



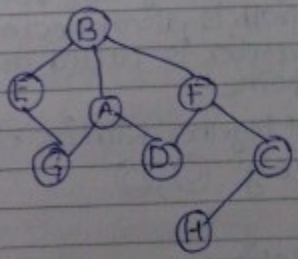


H-E-B-A-C-G-F-D

D-H-F-C-G-A-B-E

Q. Consider the following graph:-

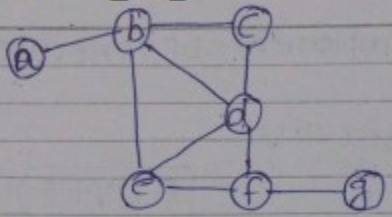
Give 4 diff. DFT's starting from H.



1. H-C-F-D-A-B-E
2. H-C-F-B-E-G-D
3. H-C-F-D-A-G-E
4. H-C-F-B-A-D-G
5. H-C-F-B-A-G-D

Q. Consider the following graph:-

Cosmos

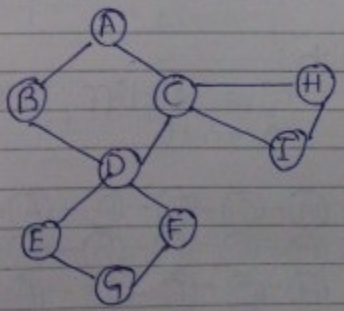


Which one of the following is correct DFT?

(a) b, c, d, e, f, a, g
(b) d, b, a, e, f, c, g
(c) b, a, e, c, d, f, g
(d) d, e, b, a, c, f, g

Applications Of DFT

1. We can check whether given graph is connected or not.
2. Finding no. of connected components.
3. Checking given graph contain cycle or not.
4. Checking given graph is strongly connected or not.
5. Finding no. of articulation points or not.
6. Finding no. of bridges.



Articulation pts.:-

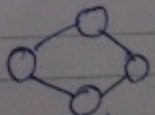
In the given connected graph, by deleting any vertex* if the graph is disconnected, then the vertex is called articulation point.

* together with its edges

e.g. C, D

Bridges:-

In the given connected graph, by deleting any edge, if the graph is disconnected, then edge is called bridge.



Date

★In order to convert recursive procedure to iterative procedure (or non-recursive procedure) requires stack data structure.

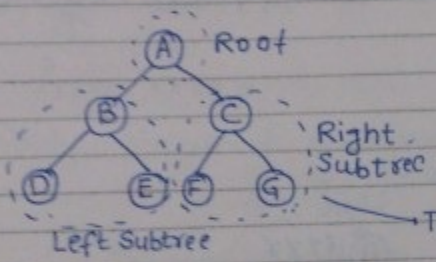
28.07.12

Tree Traversal

① Inorder (LST Root RST)
② Preorder (Root LST RST)
③ Postorder (LST RST Root)

Time taken for tree traversal: O(V+E) (from graphs); no. of edges in tree are (V-1)

∴ O(V+V-1) ⇒ O(V)



preorder :- ABDECFG
postorder :- DEBFGCA
inorder :- DBEAFCG

There are 'V' function calls.

```
Preorder(t)
{ print(t->data);
  Preorder(t->
  if(t!=NULL)
   { printf(t->data);
     Preorder(t->leftchild);
     Preorder(t->rightchild);
   }
}
```

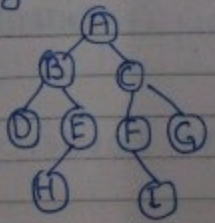
```
inorder(t)
{ if(t!=NULL)
   { inorder(t->leftchild);
     printf(t->data);
     inorder(t->rightchild);
   }
}
```

★Note:- If the binary tree contain V nodes, inorder, preorder & postorder traversal takes θ(V) time.

```
postorder(t)
{ if(t!=NULL)
   { postorder(t->leftchild);
     postorder(t->rightchild);
     printf(t->data);
   }
}
```


Q1. Consider the following C program.

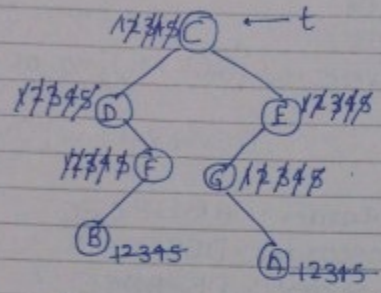
```
→ Aorder(t)
  { if(t!=NULL)
     { printf(t->data);
       Aorder(t->rightchild);
       printf(t->data);
       Aorder(t->leftchild);
```



O/p:-
AEGGGCFIIIFFC
BEEHHHEBDDD

Cosmos

Q. Apply the above code on the following binary tree:-

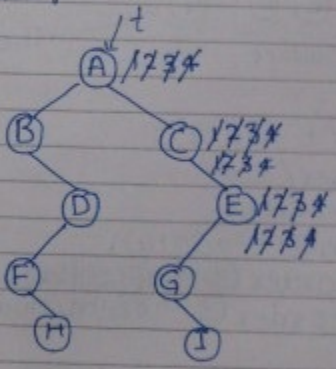


O/p:-

CEEGAAAGGECDFFBBBFDDC

Q. Consider following C program:-

```
B(order)
{if(t!=NULL)
  {1.Border(t->right);
   2.printf(t->data);
   3.Border(t->right);
   4.printf(t->data);
  }
}
```



O/p:-

EECEECAEECEECA

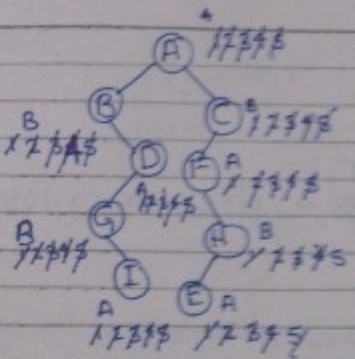
Q. Consider following C program:-

```
A(ord Aorder(t)
  { if (t!=NULL)
    {1.printf(t->data);
     2.Border(t->right);
     3.printf(t->data);
     4.Border(t->left);
     5.printf(t->data);
    }
  }
```

```
Border(t)
{if(t!=NULL)
  {1.printf(t->data);
   2.Aorder(t->left);
   3.printf(t->data);
   4.Aorder(t->right);
   5.printf(t->data);
  }
}
```

Start from A(order Aorder(t)

Cosmos

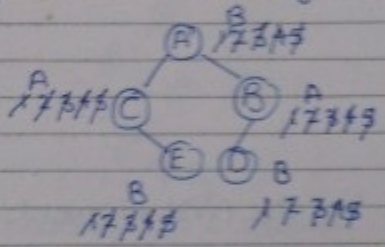


O/p:-

ABFI

ACFHEEEHHFFCCABBDDGGIIIGDBA

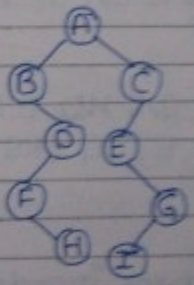
Q. Apply above C program starting from Border



O/p:-

ACEEECCABBDDDBA

Q. Consider the following binary tree:

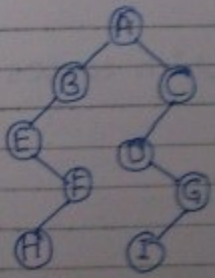


preorder :- ABDFHCEGI

inorder :- BFHDAEIGC

postorder :- HFDBIGECA

Q. Consider the following binary tree:-



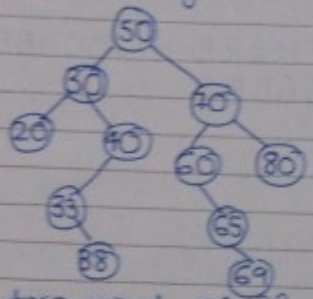
preorder :- ABEFHCDGI

inorder :- EHFBADIGC

postorder :- HFEBIGDCA

Cosmos

Q. Consider following BST :-



Inorder :- 20,30,33,38,40,50,60,65,69,70,80

★ Inorder travesal of Binary Search tree will always give ascending order.

Q.

pre :- A, B, D, F, H, C, E, G, I
in :- B, F, H, D, A, E, I, G, C

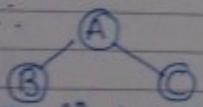
Step 1 :- identify the root, it will be the first element of preorder.

(A)

Step 2 :- now, check A in inorder, left to A is LST of A & right to A is the RST of A.

∴ LST :- B F H D
RST :- E I G C

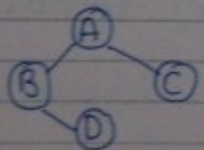
from preorder B comes first from D F H B (LST)
& from preorder C comes first from E I C G (RST).



step3 :- in order to find LST & RST of B check inorder see left of B in inorder

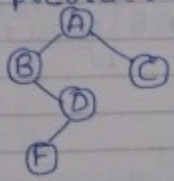
LST :- NULL
RST :- F, H, D preorder

from inorder, we check from F, H, D which comes 1st.
D comes 1st.

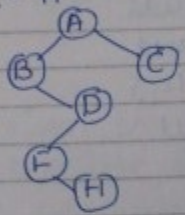


Cosmos

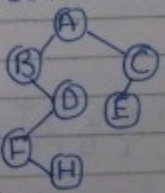
Step4 :- Checking LST & RST of D (from inorder)
LST :- H, F
RST :- Null
from preorder F comes 1st from H, F.



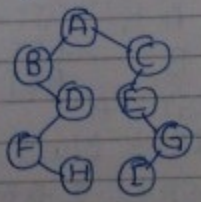
Step5 :- Checking LST & RST of F (from inorder)
LST :- Null
RST :- H



Step6 :- Checking LST & RST of C (from inorder)
LST :- E, I, G
RST :- Null
Now E comes 1st from E, I, G from preorder



Step7 :- Similarly :-

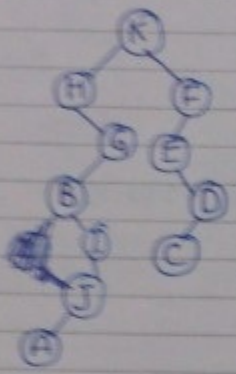


Note :-
If postorder & inorder are given, then procedu is similar, the only diff is that we have to chec the last element from postorder (in the amon the elements of a subtree).

Cosmos

Time complexity $O(n^2)$ for constructing the tree from pre/post & inorder because for every element, we have to search all n elements & there are total n elements, $n \times n$

Q. pre :- K, H, G, B, I, J, A, F, E, D, C
in :- H, B, I, A, J, G, K, E, C, D, F



K :- LST :- H B I A G
RST :- E C D F

F :- LST :- E C D
RST :- Null

E :- LST :- Null
RST :- C D

D :- LST :- C
RST :- Null

H :- LST :- Null
RST :- B I A J G

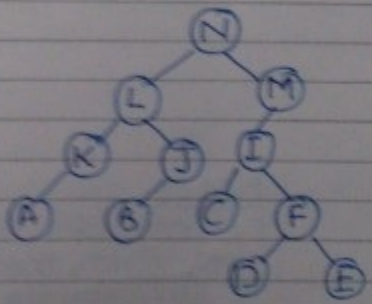
G :- LST :- B I A J
RST :- Null

B :- LST :- Null
RST :- I, A, J

I :- LST :- Null
RST :- A J

J :- LST :- A
RST :- Null

post :- L, A, M, K, G, D, B, J, I, F
in :- L, K, M, A, E, D, C, F, H, J, I, B
post :- A, K, B, I, L, C, D, E, F, I, M, N
in :- A, K, L, B, J, N, C, I, D, F, E, M



N :- LST :- A K L B J
RST :- C I D F F M

L :-
LST :- A K
RST :- B J

J :-
LST :- B
RST :- Null

K :- LST :- A
RST :- Null

M :- LST :- C I D F E
RST :- Null

I :- LST :- C
RST :- D F E

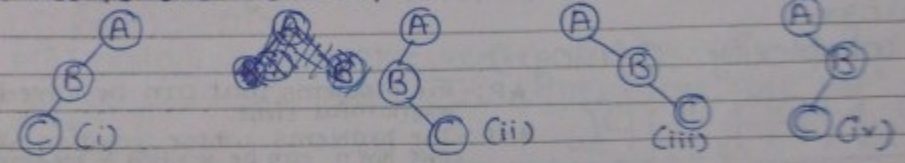
F :- LST :- D
RST :- E

C :- LST :- Null
RST :- Null

Cosmos

Note:- In order to construct unique Binary Tree fro the given (preorder & inorder) or (inorder & postorder) requires $O(n^2)$ [worst case & avera case].

Q. pre:- A, B, C
post:- C, B, A
what is the inorder?



Note:- For the given preorder & postorder, unique binary tree is not possible, ∵ inorder is not available, ∴ inorder is necessary to construct unique binary tree.

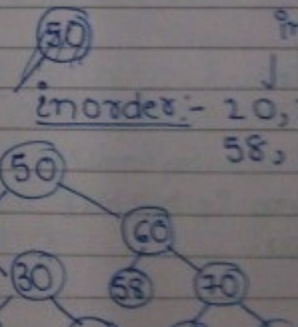
Time Complexity:- There is no algo to construct unique Binary Tree for the given preorder & postorder other th Brute Force, so it will take $O(2^n)$ times.

★ If preorder & postorder is given & we are asked to find inorder & 4 options

Q. A BST is traversed in preorder & values are printed i the following order:-
preorder:- 50, 20, 30, 25, 40, 60, 58, 70, 65, 80
what is the postorder:-

50 → root
50:- LST :- 20, 30, 25, 40
RST :- 60, 58, 70, 65, 80

postorder:- 25, 40, 30, 20, 58, 65, 80, 70, 60, 50



★NPC ~~problems~~ doesn't apply directly to optimization problems, but rather applies to decision problems, in which the answer is simply 'yes' or 'no' (or, '0' or '1').

(Note:- In order to construct unique B.T. from given (inorder & postorder) or (inorder & preorder) requires O(nlogn) time, if inorder is already sorted.

① → To sort inorder:- nlogn
② → To search every element in inorder, we can perform Binary Search:- logn
∴ 4 every element (n) will perform binary search,
∴ nlogn
∴ time complexity:- O(nlogn)

P, NP, NPH, NPC

★P:- The problems that can be solved in polynomial time.
★NP:- The problems whose solution when once given can be verified in polynomial time

Problems
- Decidable (algo is available)
  - Decision problems (answer is yes or no)
    - P-Class (taking less time to answer yes or no.) [time to solve the problem takes max. polynomial time.] $\log_2 n$ falls into this category, though it is not polynomial, but $\log_2 n <$ polynomial
    - NP-Class (taking more time to answer yes or no.)
  - Optimization problems (answer is not yes or no.)
- Undecidable (no algorithms)

★Every P problem is NP, but not vice-versa.

★A problem 'A' is in P-Class if there exist a deterministic polynomial time algo to solve that problem.

★ P-Class problems:-
1. Binary Search :- $O(\log_2 n)$
2. Linear Search :- $O(n)$
3. Job sequencing with deadline :- $O(n^2)$
4. Quicksort :- $O(n^2)$
5. LCS (m,n) :- $O(mn)$
6. MST- prim's :- $O((V+E)\log V)$
7. ✗ Travelling Salesperson problem :- $O(2^n)$:- X → so it is not a p-class.
8. ✗ 0/1 knapsack :- $O(2^n)$ → so it is not in p-class.
9. ✗ SOS (m,n) :- $O(2^n)$:- so it is not in p-class.
10. matrix multiplication (strassen) :- $O(n^{2.81})$

• NP-Class:- A problem 'A' is in NP-class if there exist a non-deterministic polynomial time algo to solve that problem.
→ Linear Search → O(1) in case of NP, so without
→ Binary Search → O(1)
→ Job sequencing with deadline → O(n)

★ ★ if we can solve a problem w/o guess in polynomial time, then we can solve it with guess in polynomial time

in this slot, guess the best slot

∴ n-job slots

in this slot, guess the best job

∴ $O(n)$

④ Quicksort:-

take 1st pivot & guess its position → $O(1)$
similary other $(n-1)$ elements ∴ $nO(1)$ = $O(n)$

⑤ LCS(n,n):- $O(n)$
⑥ MST-prim's :- take one vertex & pick the smallest edge (guess) which will take $O(1)$ time, there are 'V' vertices, ∴ $O(V)$
⑦ Travelling Salesperson Problem:- $O(V)$
⑧ 0/1 Knapsack problem:-
for every object we have two choices
(_ _ _ _ _ _ _)
n objects (we can either take it or leave it.)

∴ guessing to take it or leave it, ∴ $O(1)$ time.
n-objects :- $O(n)$
⑨ SOS(n,m) :- $O(n)$
⑩ Strassen Matrix Multiplication:- $O(n^2)$

• Non-determinism :- Takes max polynomial time, but only guesses & no proofs.

★ Every P-class problem is NP class problem, because P-class (with proof) are taking polynomial time then when done without proof it will take less than or equal polynomial time (like that in NP class).
★ ∴ every P-problem is also a NP-problem, but every NP-problem need not be a P-problem (e.g. Travelling Salesperson is NP but not P.)

~~P ⊆ NP~~ P ⊂ NP acc. to the above statement.

Cosmos

★If we solve the Travelling Salesperson problem, then 4 other 2^n complex problems, then

$P = NP$

So, by now

$\boxed{P \subseteq NP}$

we can't say $P = NP$ & $P \subset NP$ alone.

$\therefore \boxed{P \subseteq NP}$

have to be proved, area of open research.

- $P \subset NP$ → it is true when someone proofs that Travelling Salesperson problem can't be solved in less than or equal to polynomial time.
- $P = NP$ → it is true when someone proofs that Travelling Salesperson problem can be solved in polynomial time.

★P-Class :- Any problem is solved in polynomial time without any guess.

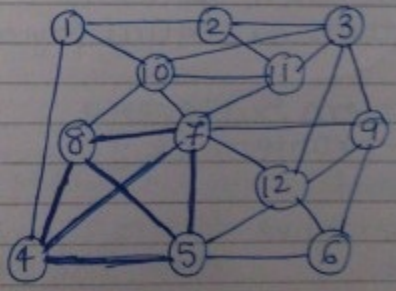
★ NP-Class :- Any problem is solved in polynomial time with guess.

★An algo without guess gives time complexity of $O(2^n)$, then it is not P, but we can't say anything about whether it is NP or not.

that's why $\boxed{P \subseteq NP}$

i/p :- Graph $G(V, E)$, integer K

Ques. Does G contain K-Clique or not ?



Q. Does the given graph contain subgraph which is k-vertex complete graph

↑

this problem is not P. because it takes $O(2^n)$ time complexity.

Cosmos

*With guess (NP-class) we choose a subgraph with k-vertices in O(1) time & then verify each vertex will take O(V) time [where V:- no. of vertices in original graph.]

★ NP means :- if we are given correct soln. to a probl then we can verify the soln. in max. polynomia time (we dont solve the problem.)

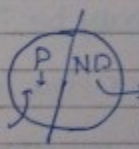
★ P means :- we can solve the problem in polynomial time.

e.g. Linear Search :- Solution O(n)

Verification (whether the element search is what we wanted or not ∴ verification will take O(1) time.

① P-Class :- any problem is solved (for all i/p's) we can say yes or no in polynomial time, then the problem is called as P-class.

② NP-Class :- any problem is verified (correct input only) in polynomial time is called NP-class.

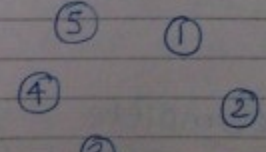


that can be solved in max. polynomial time

that can be verified in max. polynomial time.

★★ Any problem that can be easily solved (means problem can be solved in max. polynomial time) t it can also be verified easily (verification can be done in max. polynom time). ∴ every P probler is a NP problem but no vice-versa.

Polynomial Time reduction :-



• 3-SAT was the 1st NP problem discovered by Stephen Cook. It took 3 years.

• Next 3000 problems were discovered in next 10 years.

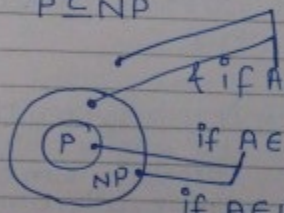
Cosmos

- A is an unknown problem & B is known problem, if A is polynomially reducible to B, then
  if B is NP takes 2^n time,
  then A will take $O(n^c)+O(2^n) = \boxed{O(2^n)}$ c:- constant.
- ★ conversion must be done in polynomial time.
  - if B is NP, then A is NP.
  - if B is P, then A is P.
- ★ If convesion takes more than polynomial time, then A & B dont have similar properties.

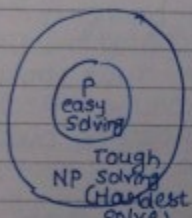
NP-Hard :-

A problem L is said to be NP-Hard if & only if L' is
- polynomially reducible to L for all $L' \in NP$.

! $L' \leq_p L, \forall L' \in NP$



& if A∈NPH, then it can be in only NP or outside NP.
if A∈P (can be solved & verified in max. polynomial time.)
if A∈NP (then it is dilemna, that it can be in both only P & only NP.)

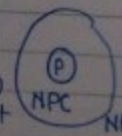


★ To prove our problem is NP-Hard, then all problems in NP should be reducible to our problem in max. polynomial time.
★ all hardest problem in NP is reducible to L, then L is also Hardest problem.

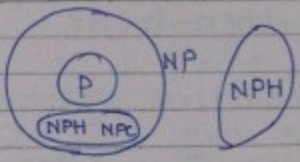
NP-Complete :-

A problem L is said to be NP-Complete
- L is NP-Hard. (or) A is NPC when A & NP
- $L \in NP$.

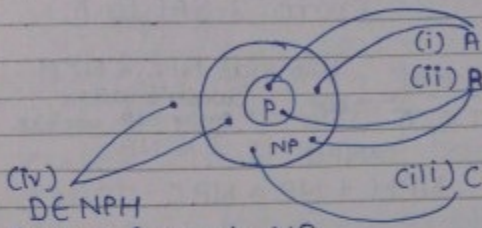
(Hardest problem inside NP (but not



- $A \in NP$ → means A is in NP but may or may not be in P.
- $A \in NPC$ → means $A \in NP$ but not in P.
- $A \in NPH$ → $A \in NPC$ or $A \notin NP$

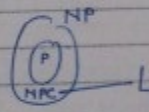


Cosmos



(i) $A \in P$ (In both)
(ii) $B \in NP$ (either in P & NP both or only NP)
(iii) $C \in NPC$ (only in NP, not in P)
(iv) $D \in NPH$ (either in only NP or outside NP).

Note1:- If L is NPC & $L \leq L'$ then L' is NPH. (Hardest)
(L is polynomially reducible to L')
NPC, NPH, undecidable



★ Hardest Language → (L) is reducible to L', then L' is hardest langua

Note2:- If L is NP & $L' \leq_p L$ then L' is NP. (Easier)
(right to left)
unknown on left.
P, NP, decidable

★ If languages are reducible to our easy language, then those languages are easy.

Q1. A, B & C are 3-decisions problems. Let A be a NPC, then which are true/false?

(a) $A \in NP$ (T)
(b) $\forall L \in NP, L \leq_p A$ (T)
(c) If $C \in NP$ & $A \leq_p C$ then C is NPC (T)
(d) If $D \in NP$ and $D \leq_p A$ then D is NPC. (F)

(c) → C is in NP & A is polynomially reducible to our problem C, then our problem is hardest, ∴ C is NPC.
(d) → D is in NP & our problem D is polynomially reducible to NPC problem, but NPC doesn't work from right to left, but D is NP.

Cosmos

(d) → If there is:- if $D \in NP$ & some $B \leq_p D$, then $B \in NP$.

Q. Two people 'A' & 'B' have been asked to show that a problem Π is NPC.
A shows polynomial time reduction from Π to 3-SAT.
B shows " " " from 3-SAT to Π.
then check for T/F?
(a) Π - P (F)
(b) Π - NP (T)
(c) Π - NPH (T)
(d) Π - NPC (T)

$\Pi \leq_p$ 3-SAT → NP
3-SAT $\leq_p \Pi$ → NPH
This side NPC, & NPH doesn't work, ∴ only NP works. ∴ Π is NP
This side both NPH & NPC works. but NPC can be converted into NPH.
∴ NPH + NP = NPC

Q. Let S be a NPC & (Q & R) be two other problems known to be in NP, if $S \leq_p R$ & $Q \leq_p S$, then T/F?
(a) $R \in NPC$ (T)
(b) $R \in NPH$ (F)
(c) $Q \in NPC$ (F)
(d) $Q \in NPH$ (F).

• $S \leq_p R$ → R is NPH but R is in NP
∴ R is NPC.

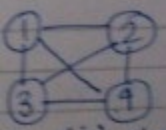
• $Q \leq_p S$ → if our problem is polynomially reducible to Harder problem,
∴ Q is NP.

Euler Graph: (If both euler path & euler cycle are possible).
Path:- sequence of zero or more edges. (not disconnected)

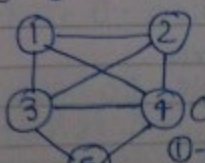
Euler Path:- It is a path which covers every edge of the given graph exactly once.

Euler Cycle (circuit):- It is a closed Euler Path (starting vertex is also the ending vertex).

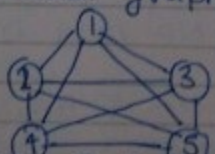
Check the following graphs are euler graphs or not.



(i) No euler path, & as no euler



(ii) 1-2-5-4-1-2-3-4-2
∴ euler path but no euler



(iii) 1-2-4-5-1-2-5
Euler Path & Euler

*If we have two vertices with odd degree & rest others with even degree, then (i) if we start from either of the odd degree vertices we will end up completing a euler circuit on the other odd degree vertex (we can't get a euler circuit). (ii) & if we start with an even degree vertex, we can't even get a euler path.

Note 1:- In the given graph, every vertex degree is even, then to that graph both euler path & euler circuit is possible, so given graph is euler graph.

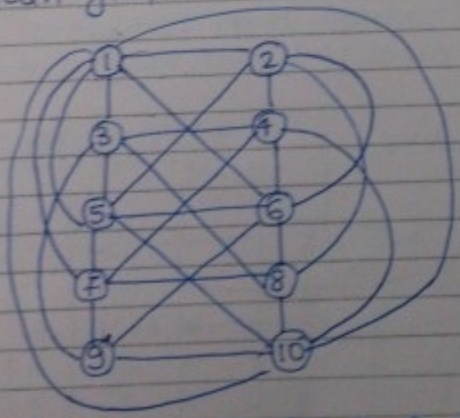
* Checking given graph is Euler Graph or not, there is $O(V^2)$ time complexity algo, so it is a P-class problem. (Checking). [checking degree of one vertex using adjacency matrix/list will take $O(V)$ time & we have to repeat this V times, ∴ complexity is $O(V^2)$.] or $O(E)$ (for dense graphs)

Note 2:- If the given graph contain exactly 2 vertices with odd degree then to that graph euler path is possible but euler circuit is not possible, ∴ given graph is not euler graph.

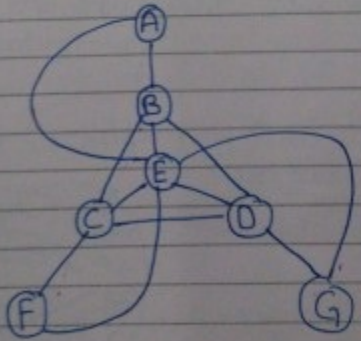
Note 3:- If the given graph contain more than 2 vertices with odd degree, euler path & euler circuit both are not possible.

Hamiltonian Graph

- Hamiltonian path:- It is a path in which every vertex of a given graph is covered exactly once.
- Hamiltonian cycle:- It is a closed Hamiltonian path (starting & ending vertex is same)
- Graph containing both Hamiltonian cycle & path is a Hamiltonian graph.



1-4-5-7-9-10-8-6-4-2-1
Hamiltonian path & cycle possible.
∴ Hamiltonian Graph.



A-B-C-F-E-D-G
- Hamiltonian path
- but no hamiltonian cycle.

★ There is no algo to check given graph is Hamiltonian graph or not other than Brute force, so it is one of the NPC.

3-SAT (3 variable satisfiable problem)

$(x_1 \vee x_2 \vee x_5) \wedge (x_6' \vee x_{10}' \vee x_2') \wedge (x_{1000} \vee x_8 \vee x_9) \wedge (x_{50} \vee x_{99} \vee x_{26})$

Note :- 3-SAT is NPC, but 2-SAT problem is P.

P-Class	NP-Class
① Shortest path	① Longest path.
② Euler	② Hamiltonian
③ 2-SAT	③ 3-SAT
④ Edge-cover	④ Vertex-cover.

Cosmos